

INTERNATIONAL STANDARD

ISO
8326

First edition
1987-08-15



INTERNATIONAL ORGANIZATION FOR STANDARDIZATION
ORGANISATION INTERNATIONALE DE NORMALISATION
МЕЖДУНАРОДНАЯ ОРГАНИЗАЦИЯ ПО СТАНДАРТИЗАЦИИ

Information processing systems — Open Systems Interconnection — Basic connection oriented session service definition

*Systèmes de traitement de l'information — Interconnexion de systèmes ouverts — Service de
session en mode connexion*

Reference number
ISO 8326:1987 (E)

Foreword

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International Standard ISO 8326 was prepared by Technical Committee ISO/TC 97, *Information processing systems*.

Users should note that all International Standards undergo revision from time to time and that any reference made herein to any other International Standard implies its latest edition, unless otherwise stated.

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Information processing systems — Open Systems Interconnection — Basic connection oriented session service definition

0 Introduction

This International Standard is one of a set of International Standards produced to facilitate the interconnection of computer systems.

This International Standard is related to other International Standards in the set as defined by the Reference Model for Open Systems Interconnection (ISO 7498). The Reference Model subdivides the area of standardization for interconnection into a series of layers of specification, each of manageable size.

The purpose of this International Standard is to define the service provided to the Presentation Layer at the boundary between the Session and Presentation Layers of the Reference Model. The session service is provided by the session protocol making use of the services available from the Transport Layer. This International Standard also defines the session service characteristics which the presentation protocol may exploit. The relationship between the International Standards for the session service, session protocol, transport service, and the presentation protocol is illustrated in figure 1.

It is recognized that, with respect to session Quality of Service, (described in clause 10), work is still in progress to provide an integrated treatment of QOS across all of the layers of the OSI Reference Model and to ensure that the individual treatments in each layer service satisfy overall QOS objectives in a consistent manner. As a consequence, an addendum may be added to this International Standard at a later time which reflects further QOS developments and integration.

1 Scope and field of application

This International Standard defines in an abstract way the externally visible service provided by the OSI Session Layer in terms of

- a) the primitive actions and events of the service;
- b) the parameter data associated with each primitive action and event;
- c) the relationship between, and the valid sequence of these actions and events.

The service defined in this International Standard is that which is provided by the OSI session protocol (in conjunction with the transport service) and which may be used by the OSI presentation protocol.

This International Standard does not specify individual implementations or products, nor does it constrain the implementation of entities and interfaces within a computer system. There is, therefore, no conformance to this International Standard.

2 References

ISO 7498, *Information processing systems — Open Systems Interconnection — Basic Reference Model*.

ISO 7498/Add.3, *Information processing systems — Open Systems Interconnection — Basic Reference Model — Addendum 3: Name including addressing*.¹⁾

ISO 8072, *Information processing systems — Open Systems Interconnection — Transport service definitions*.

ISO 8327, *Information processing systems — Open Systems Interconnection — Basic connection oriented session protocol specification*.

ISO/TR 8509, *Information processing systems — Open Systems Interconnection — Service conventions*.

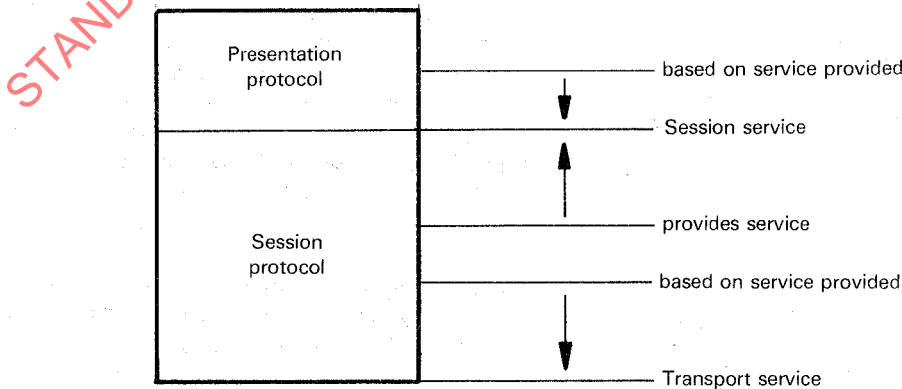


Figure 1 — Relationship of this International Standard to other OSI standards

1) At present at stage of draft; publication anticipated in due course.

Section one: General

3 Definitions

NOTE — The definitions contained in this clause make use of abbreviations defined in clause 4.

3.1 Reference Model definitions

This International Standard is based on the concepts developed in ISO 7498, and makes use of the following terms defined in it:

- a) expedited-session-service-data-unit;
- b) session-connection;
- c) Session Layer;
- d) session-service;
- e) session-service-access-point;
- f) session-service-data-unit;
- g) Transport Layer;
- h) duplex;
- i) half-duplex.

3.2 Service convention definitions

This International Standard also makes use of the following terms defined in ISO/TR 8509, as they apply to the Session Layer:

- a) service-user;
- b) service-provider;
- c) primitive;
- d) request;
- e) indication;
- f) response;
- g) confirm.

3.3 Session-service definitions

For the purpose of this International Standard, the following definitions also apply.

3.3.1 calling SS-user: An SS-user that initiates a session connection establishment request.

3.3.2 called SS-user: An SS-user with whom a calling SS-user wishes to establish a session connection.

NOTE — Calling SS-users and called SS-users are defined with respect to a single connection. An SS-user can be both a calling and called SS-user simultaneously.

3.3.3 sending SS-user: An SS-user that acts as a source of data during the data transfer phase of a session connection.

3.3.4 receiving SS-user: An SS-user that acts as a sink of data during the data phase of a session connection.

NOTE — An SS-user can be both a sending and a receiving SS-user simultaneously.

3.3.5 requestor; requesting SS-user: An SS-user that initiates a particular action.

3.3.6 acceptor; accepting SS-user: An SS-user that accepts a particular action.

3.3.7 token: An attribute of a session connection which is dynamically assigned to one SS-user at a time to permit certain services to be invoked.

3.3.8 conditional (parameter): A parameter whose presence in a request or response depends on conditions defined in the text of this International Standard; and whose presence in an indication or confirm is mandatory if that parameter was present in the preceding session service primitive, or absent if that parameter was absent in the preceding session service primitive.

3.3.9 proposed parameter: The value for a parameter proposed by an SS-user in an S-CONNECT request or an S-CONNECT response that it wishes to use on the session connection.

3.3.10 selected parameter: The value for a parameter that has been chosen for use on the session connection.

4 Symbols and abbreviations

4.1 Abbreviations

SS	: session-service
SSAP	: session-service-access-point
SSDU	: session-service-data-unit
NSSDU	: normal-data-session-service-data-unit
TSSDU	: typed-data-session-service-data-unit
XSSDU	: expedited-session-service-data-unit
QOS	: quality of service

4.2 Service variables

V(A)	See 11.4.1.1
V(M)	See 11.4.1.2
V(R)	See 11.4.1.3
Vsc	See 11.4.1.4

5 Conventions

This International Standard uses the descriptive conventions defined in ISO/TR 8509 except that, where indicated in this International Standard, parameter values associated with a service primitive may be passed in a direction opposite to the direction of the service primitive.

6 Model of the session service

This International Standard uses the abstract model for a layer service defined in ISO/TR 8509. The model defines the interactions between the SS-user and the SS-provider which take place at the two SSAPs. Information is passed between an SS-user and the SS-provider by service primitives, which may convey parameters.

7 Overview of the session service

7.1 General overview

The session service provides the means for organized and synchronized exchange of data between cooperating SS-users. It provides its users with means to

- establish a connection with another SS-user, exchange data with that user in a synchronized manner, and release the connection in an orderly manner;
- negotiate for the use of tokens to exchange data, synchronize and release the connection, and to arrange for data exchange to be half-duplex or duplex;
- establish synchronization points within the dialogue and, in the event of errors, resume the dialogue from an agreed synchronization point;
- interrupt a dialogue and resume it later at a prearranged point.

7.2 Token concept

A token is an attribute of a session connection which is dynamically assigned to one SS-user at a time to permit certain services to be invoked. It is the right to exclusive use of the service.

Four tokens are defined:

- the data token;
- the release token;
- the synchronize-minor token;
- the major/activity token.

A token is always in one of the following states:

- available, in which case it is always
 - assigned to one SS-user, who then has the exclusive right to use the associated service (provided that no other restrictions apply); and
 - not assigned to the other SS-user, who does not have the right to use the service but may acquire it later;
- or
- not available to either SS-user, in which case neither SS-user has the exclusive use of the associated service. The service then becomes inherently available to both SS-users (data transfer and release), or otherwise unavailable to both SS-users (synchronization and activities).

Restrictions related to the availability and assignment of tokens are defined in 11.2.

7.3 Synchronization and dialogue unit concepts

SS-users may insert synchronization points into the data they are transmitting. Each synchronization point is identified by a serial number maintained by the SS-provider (see 11.4).

Any semantics which SS-users may give to their synchronization points are transparent to the SS-provider.

There are two types of synchronization points:

- minor synchronization points;
- major synchronization points.

Major synchronization points are used to structure the exchange of data into a series of dialogue units. The characteristic of a dialogue unit is that all communication within it is completely separated from all communication before and after it. A major synchronization point indicates the end of one dialogue unit and the beginning of the next. Each major synchronization point is confirmed explicitly.

Minor synchronization points are used to structure the exchange of data within a dialogue unit. Figure 2 illustrates how a dialogue unit is structured through the use of minor synchronization points. Each minor synchronization point may or may not be confirmed explicitly.

7.4 Activity concept

The activity concept allows SS-users to distinguish between different logical pieces of work called activities. Each activity consists of one or more dialogue units. Only one activity is allowed on a session connection at a time, but there may be several consecutive activities during a session connection. An activity may also span more than one session connection.

An activity can be interrupted and then resumed on the same or on a subsequent session connection. This can be considered as a form of resynchronization.

Figure 3 shows how an activity may be structured into dialogue units through the use of major synchronization points. In addition, the SS-users may transfer data outside an activity.

7.5 Resynchronization

Resynchronization may be initiated by either SS-user. It sets the session connection to a defined state, and therefore includes reassignment of tokens and setting the synchronization point serial number to a new value. Resynchronization purges all undelivered data.

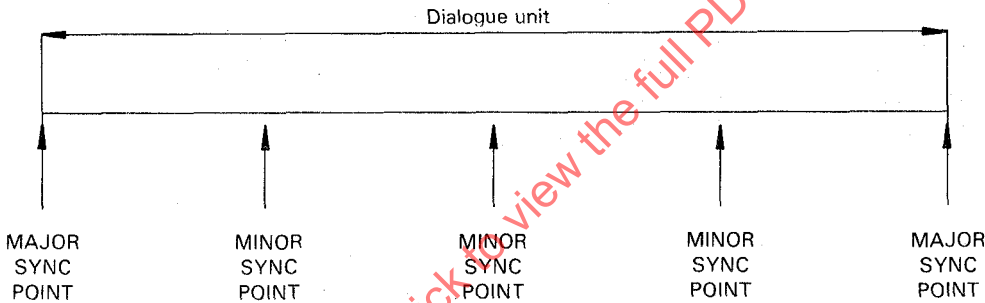


Figure 2 — Example of a structured dialogue unit

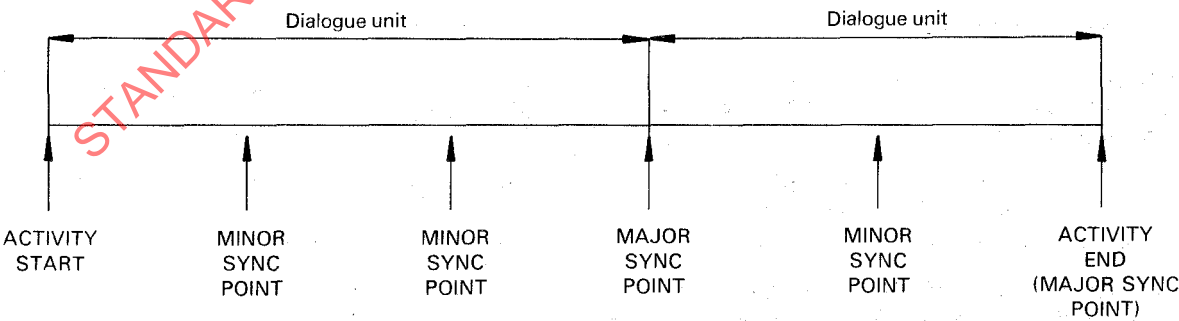


Figure 3 — Example of a structured activity

Three options are provided:

- a) abandon option which is used to set the synchronization point serial number to an unused value;
- b) restart option which is used to set the synchronization point serial number to any used value which is greater than the synchronization point serial number which identifies the last acknowledged major synchronization point;
- c) set option which is used to set the synchronization point serial number to any value chosen by the SS-user.

7.6 Negotiation

Negotiation takes place between both SS-users during the session connection establishment phase according to the following rules.

7.6.1 Negotiation of functional units

The kernel functional unit (see clause 9) is always used. Each SS-user proposes the use or non-use of each of the other functional units. A functional unit is selected only if both SS-users propose use of the functional unit and it is supported by the SS-provider. Specific negotiation rules are given in 12.1.2.

7.6.2 Negotiation of initial token settings

When the calling SS-user proposes use of a functional unit that requires a token, it also proposes the initial token settings:

- a) calling SS-user side;
- b) called SS-user side;
- c) called SS-user choice.

If the use of the functional unit is selected, the token is set to

- d) the side proposed by the called SS-user, if "called SS-user choice" is proposed by the calling SS-user; or
- e) in all other cases, the side proposed by the calling SS-user.

7.6.3 Negotiation of initial synchronization point serial number

When a calling SS-user proposes any of the major synchronize, minor synchronize or resynchronize functional units, but does not propose the activity management functional unit, it also proposes a value for the initial synchronization point serial number.

The calling SS-user may also propose a value for the initial synchronization point serial number even if the activity management functional unit is proposed provided that any of the minor synchronize, major synchronize or resynchronize functional units are also proposed. If the called SS-user selects use of any of the minor synchronize, major synchronize or resynchronize functional units, but does not select use of the activity management functional unit, it returns a value for the initial synchronization point serial number which may or may not be the same as the value proposed by the calling SS-user. The value returned by the called SS-user is used as the initial synchronization point serial number for the session connection.

In all other combinations of functional units, no initial synchronization point serial number is proposed.

8 Phases and services of the session service

The session service comprises three phases. The purpose of each phase, and a short description of the associated services is given in this clause. The services and the primitives by which they are invoked are defined in clauses 12, 13 and 14.

8.1 Session connection establishment phase

The session connection establishment phase is concerned with establishing a connection between two SS-users. It has one service associated with it:

the session connection service (see 12.1) is used to set up a session connection and to negotiate tokens and parameters to be used for the connection.

8.2 Data transfer phase

The data transfer phase is concerned with the exchange of data between the two SS-users connected in the session connection establishment phase.

There are four services associated with data transfer:

- a) the normal data transfer service (see 13.1) allows the transfer of normal data SSDUs (NSSDUs) over a session connection. Its use is controlled by the data token if the half-duplex functional unit has been selected;
- b) the expedited data transfer service (see 13.2) allows the transfer of expedited SSDUs (XSSDUs) over a session connection free from the token and flow control constraints of the normal data transfer service, typed data transfer service and capability data exchange service;
- c) the typed data transfer service (see 13.3) is used to transfer typed data SSDUs (TSSDUs) independent of the availability and assignment of the data token;
- d) the capability data exchange service (see 13.4) is used to exchange a limited amount of confirmed SS-user data while not within an activity.

There are three services concerned with token management:

- e) the give tokens service (see 13.5) allows an SS-user to surrender one or more specific tokens to the other SS-user;
- f) the please tokens service (see 13.6) allows an SS-user to request the other SS-user to transfer one or more specific tokens to it;
- g) the give control service (see 13.7) allows an SS-user to surrender all available tokens to the other SS-user.

There are three services associated with synchronization and resynchronization:

- h) the minor synchronization point service (see 13.8) allows the SS-user to separate the flow of NSSDUs and TSSDUs transmitted before the service was invoked from the subsequent flow of NSSDUs and TSSDUs. Its use is controlled by the synchronize-minor token;

- i) the major synchronization point service (see 13.9) allows the SS-user to confine the flow of sequentially transmitted NSSDUs, TSSDUs and XSSDUs in each direction within a dialogue unit. Its use is controlled by the major/activity token;
- j) the resynchronize service (see 13.10) is used to set the session connection to a previous or to a new synchronization point and to reassign the available tokens. This service may cause loss of NSSDUs, TSSDUs and XSSDUs.

There are two services for reporting errors or unanticipated situations:

- k) the provider-initiated exception reporting service (see 13.11) (P-exception reporting service) permits SS-users to be notified of exception conditions or SS-provider protocol errors. This service may cause loss of NSSDUs, TSSDUs and XSSDUs;
- l) the user-initiated exception reporting service (see 13.12) (U-exception reporting service) is used by the SS-user to report an exception condition when the data token is available but not assigned to the SS-user. This service may cause loss of NSSDUs, TSSDUs and XSSDUs.

There are five services associated with activities:

- m) the activity start service (see 13.13) is used to indicate that a new activity is entered. Its use is controlled by the major/activity token;
- n) the activity resume service (see 13.14) is used to indicate that a previously interrupted activity is re-entered. Its use is controlled by the major/activity token;
- o) the activity interrupt service (see 13.15) allows an activity to be abnormally terminated with the implication that the work so far achieved is not to be discarded and may be resumed later. Its use is controlled by the major/activity token. This service may cause loss of NSSDUs, TSSDUs and XSSDUs;
- p) the activity discard service (see 13.16) allows an activity to be abnormally terminated with the implication that the work so far achieved is to be discarded, and not resumed. Its use is controlled by the major/activity token. This service may cause loss of NSSDUs, TSSDUs and XSSDUs;
- q) the activity end service (see 13.17) is used to end an activity (and set a major synchronization point). Its use is controlled by the major/activity token;

Using the activity services may lead to a state where no activity is in progress on the session connection. When activity services are employed, but no activity is in progress, only the activity start, activity resume, token management, capability data, typed data, normal data, expedited data, abort and release services may be invoked by the SS-users.

8.3 Session connection release phase

The session connection release phase is concerned with releasing a previously established session connection. It has three services associated with it:

- a) the orderly release service (see 14.1) provides a means of achieving the orderly release of a session connection;
- b) the user-initiated abort service (see 14.2) (U-abort service) is used to initiate the release of a session connection in a way that will terminate any outstanding service request. This service may cause loss of NSSDUs, TSSDUs and XSSDUs;
- c) the provider-initiated abort service (see 14.3) (P-abort service) is used by the SS-provider to indicate the release of the session connection for internal reasons. This service may cause loss of NSSDUs, TSSDUs and XSSDUs. Any outstanding service request is terminated.

9 Functional units and subsets

9.1 Functional units

Functional units are logical groupings of related services defined by this International Standard for the purpose of

- a) negotiation of SS-user requirements during the session connection establishment phase;
- b) reference by other International Standards.

Table 1 specifies the association of tokens and functional units. When a functional unit implies the availability of a token, services concerned with the management of that token are provided in order to be able to request and transfer the available tokens.

The services associated with each functional unit are specified in table 2.

Table 1 — Functional units using tokens

Functional unit	Token
Negotiated release	Release token
Half-duplex	Data token
Minor synchronize	Synchronize-minor token
Major synchronize	Major/activity token
Activity management	Major/activity token

Table 2 — Services associated with each functional unit

Functional unit	Service(s)	Reference
Kernel (non-negotiable)	Session connection	12.1
	Normal data transfer	13.1
	Orderly release	14.1
	U-Abort	14.2
	P-Abort	14.3
Negotiated release	Orderly release	14.1
	Give tokens	13.5
	Please tokens	13.6
Half-duplex	Give tokens	13.5
	Please tokens	13.6
Duplex	No additional service	
Expedited data	Expedited data transfer	13.2
Typed data	Typed data transfer	13.3
Capability data exchange	Capability data exchange	13.4
Minor synchronize	Minor synchronization point	13.8
	Give tokens	13.5
	Please tokens	13.6
Major synchronize	Major synchronization point	13.9
	Give tokens	13.5
	Please tokens	13.6
Resynchronize	Resynchronize	13.10
Exceptions	Provider exception reporting	13.11
	User exception reporting	13.12
Activity management	Activity start	13.13
	Activity resume	13.14
	Activity interrupt	13.15
	Activity discard	13.16
	Activity end	13.17
	Give tokens	13.5
	Please tokens	13.6
	Give control	13.7

9.1.1 Kernel functional unit

The kernel functional unit supports the basic session services required to establish a session connection, transfer normal data and release the session connection.

9.1.2 Negotiated release functional unit

The negotiated release functional unit supports the negotiated orderly release service. The release token is available when this functional unit has been selected.

9.1.3 Half-duplex functional unit

The half-duplex functional unit supports the half-duplex service. The data token is available when this functional unit has been selected. It is not possible to select both this functional unit and the duplex functional unit for use on the same session connection.

9.1.4 Duplex functional unit

The duplex functional unit supports the duplex service. It is not possible to select both this functional unit and the half-duplex functional unit for use on the same session connection.

9.1.5 Expedited data functional unit

The expedited data functional unit supports the session expedited data transfer service.

9.1.6 Typed data functional unit

The typed data functional unit supports the typed data transfer service.

9.1.7 Capability data exchange functional unit

The capability data exchange functional unit supports the capability data exchange service. This functional unit can only be selected when the activity management functional unit has been selected.

9.1.8 Minor synchronize functional unit

The minor synchronize functional unit supports the minor synchronization point service. The synchronize-minor token is available when this functional unit has been selected.

9.1.9 Major synchronize functional unit

The major synchronize functional unit supports the major synchronization point service. The major/activity token is available when this functional unit has been selected.

9.1.10 Resynchronize functional unit

The resynchronize functional unit supports the resynchronize service.

9.1.11 Exceptions functional unit

The exceptions functional unit supports the user exception and provider exception reporting services.

This functional unit can only be selected when the half-duplex functional unit has been selected.

9.1.12 Activity management functional unit

The activity management functional unit supports the activity management services and the give control service. The major/activity token is available when this functional unit has been selected.

9.2 Subsets

A subset is defined as a combination of the kernel functional unit together with any other set of functional units provided that

- a) if the capability data functional unit is included in the subset, then the activity management functional unit is also included in the subset; and
- b) if the exceptions functional unit is included in the subset, then the half-duplex functional unit is also included in the subset.

NOTE — This International Standard contains no requirements for the registration of subsets. Users of this International Standard may define subsets to meet their session service needs. Other International Standards may identify subsets that comply with the above definition.

10 Quality of session service

The term "Quality of Service" (QOS) refers to certain characteristics of a session connection as observed between the session connection endpoints. QOS describes aspects of a session connection which are attributable solely to the SS-provider; such aspects are independent of SS-user behaviour (which is beyond the control of the SS-provider). SS-user behaviour does not impact the QOS provided.

Once a session connection is established, the SS-users at the two ends have the same knowledge and understanding of what the QOS over the session connection is.

10.1 Determination of QOS

QOS is described in terms of QOS parameters.

The definition of the QOS parameters associated with the session service is given in 10.3. These definitions provide both SS-users and the SS-provider with a common understanding of the QOS characteristics.

Two types of session service QOS parameters are identified:

- a) those which are negotiated during the session connection establishment phase:
 - 1) session connection protection (see 10.3.9);
 - 2) session connection priority (see 10.3.10);
 - 3) residual error rate (see 10.3.5);
 - 4) throughput, for each direction of transfer (see 10.3.3);
 - 5) transit delay, for each direction of transfer (see 10.3.4);
 - 6) optimized dialogue transfer (see 10.3.13); and
 - 7) extended control (see 10.3.12);
- b) those which are not negotiated during the session connection establishment phase but whose values are selected and/or known by other methods (for example, *a priori* knowledge and agreement, or by means of management functions) not defined in this International Standard:
 - 1) session connection establishment delay (see 10.3.1);
 - 2) session connection establishment failure probability (see 10.3.2);
 - 3) transfer failure probability (see 10.3.6);
 - 4) session connection release delay (see 10.3.7);
 - 5) session connection release failure probability (see 10.3.8);
 - 6) session connection resilience (see 10.3.11).

The negotiation procedures for parameters listed in a) are defined in 10.2. Once the session connection is established, the selected QOS parameters are not re-negotiated during the lifetime of the session connection. The SS-user should be aware that changes in QOS during a session connection are not signalled in the session service.

10.2 Session connection QOS negotiation procedures

QOS negotiation is described in terms of parameters which can be conveyed by the S-CONNECT primitives during the session connection establishment phase (see clause 12). For the parameters which are negotiated during the session connection establishment phase [see 10.1a)], the parameter values and negotiation rules are defined as follows:

- a) In the S-CONNECT request primitive, the calling SS-user can specify:
 - 1) for session connection protection, session connection priority, extended control, and optimized dialogue transfer, a single parameter value which is the "desired" QOS; for extended control and optimized dialogue transfer, one of the two values "feature desired" or "feature not desired" is conveyed;

NOTE — If the calling SS-user proposes use of the expedited data functional unit, the extended control parameter has the value "feature desired".

- 2) for residual error rate, and for each direction of throughput and transit delay, two parameter values which are the "desired" QOS and the "lowest acceptable" QOS to which the calling SS-user will agree.

b) In the S-CONNECT indication primitive, for each of the negotiated parameters, an "available" value is conveyed which is specified as follows:

1) for session connection protection, if the SS-provider agrees to provide a QOS value equivalent to the "desired" value specified in the S-CONNECT request, then the SS-provider specifies that value as "available"; if the SS-provider does not agree to provide the "desired" QOS requested, the SS-provider refuses to establish the session connection by issuing the S-CONNECT (reject) confirm primitive to the calling SS-user;

2) for session connection priority, the SS-provider specifies the QOS value it is willing to provide (a value which is equal to or better than the "desired" value specified in the S-CONNECT request) as "available";

3) for the residual error rate and each direction of throughput and transit delay, if the SS-provider agrees to provide a value of QOS which is equal to or better than the "lowest acceptable" QOS value specified in the S-CONNECT request, then the SS-provider specifies the value as "available"; if the SS-provider does not agree to provide this QOS, then the SS-provider refuses to establish the session connection by issuing the S-CONNECT (reject) confirm primitive to the calling SS-user;

4) for extended control and optimized dialogue transfer, if the "desired" value in the S-CONNECT request primitive is "feature not desired" then "feature not desired" is specified as "available"; if the "desired" value is "feature desired" and the SS-provider agrees to provide the feature on the session connection, then "feature desired" is specified as "available"; otherwise if the SS-provider does not agree to provide the feature, "feature not desired" is specified as "available".

c) In the S-CONNECT response primitive, for each of the negotiated parameters, an "agreed" value is conveyed which is specified as follows:

1) for optimized dialogue transfer, if the "available" value in the S-CONNECT indication primitive is "feature not desired" and the called SS-user agrees not to have the feature provided on the session connection, then "feature not desired" is specified as "agreed"; otherwise the SS-user may reject establishment of the session connection; if the "available" value in the indication primitive is "feature desired" and the SS-user agrees to have the feature provided, then "feature desired" is specified as "agreed"; otherwise, if the SS-user does not agree to provision of the feature, the value "feature not desired" is specified as "agreed";

2) for each of the other parameters, if the called SS-user agrees to the QOS value specified as "available" in the S-CONNECT indication primitive, then the identical value is specified as "agreed"; if the SS-user does not agree to the "available" value, the SS-user may reject establishment of the session connection.

d) In the S-CONNECT confirm primitive, for each of the negotiated parameters, an "agreed" value is conveyed which is identical to the "agreed" value conveyed in the S-CONNECT response.

10.3 Definition of QOS parameters

QOS parameters can be classified as

- parameters which express session service performance parameters, as shown in table 3;
- parameters which express other session service characteristics, as shown in table 4.

These session service QOS parameters are defined in this sub-clause.

Table 3 — Classification of performance QOS parameters

Phase	Performance criterion	
	Speed	Accuracy/reliability
Session connection establishment	Session connection establishment delay	Session connection establishment failure probability (misconnection/session connection refusal)
Data transfer	Throughput Transit delay	Residual error rate (corruption, duplication/loss) Session connection resilience Transfer failure probability
Session connection release	Session connection release delay	Session connection release failure probability

Table 4 — Parameters specifying other session service features

Extended control Session connection protection Session connection priority Optimized dialogue transfer

10.3.1 Session connection establishment delay

Session connection establishment delay is the maximum acceptable delay between an S-CONNECT request and the corresponding S-CONNECT confirm primitive.

NOTE — This delay includes SS-user dependent components.

10.3.2 Session connection establishment failure probability

Session connection establishment failure probability is the ratio of total session connection establishment failures to total session connection establishment attempts in a measurement sample.

Session connection establishment failure is defined to occur when a requested session connection is not established within the specified maximum acceptable session connection establishment delay as a result of misconnection, session connection refusal, or excessive delay on the part of the SS-provider. Session connection establishment attempts which fail as a result of error, session connection refusal, or excessive delay on the part of an SS-user are excluded in calculating session connection establishment failure probability.

10.3.3 Throughput

Throughput is defined for each direction of transfer, in terms of a sequence of at least two SSDUs successfully transferred by an S-DATA request/S-DATA indication or S-TYPED-DATA request/S-TYPED-DATA indication sequence of primitives. Given such a sequence of n SSDUs, where $n \geq 2$, the throughput is defined to be the smaller of

- the number of SS-user data octets contained in the last $n-1$ SSDUs divided by the time between the first and last S-DATA or S-TYPED-DATA request in the sequence, and
- the number of SS-user data octets contained in the last $n-1$ SSDUs divided by the time between the first and last S-DATA or S-TYPED-DATA indications in the sequence.

Successful transfer of the octets in a transmitted SSDU is defined to occur when the bits are delivered to the intended receiving SS-user without error, in the proper sequence, prior to release of the session connection by the receiving SS-user.

Throughput is only meaningful for a sequence of complete SSDUs and each specification is based on a previously stated average SSDU size.

Throughput is specified separately for each direction of transfer on a session connection. In each direction, a specification of throughput will consist of a "maximum throughput" value and an "average throughput" value. The "maximum throughput" value represents the maximum rate at which the SS-provider can continuously accept and deliver SSDUs, in the absence of sending SS-user input delays or flow control applied by the receiving SS-user. Thus, the sequence of SSDUs in the calculation above are defined to be presented continuously at the maximum rate. The "average throughput" value represents the expected transfer rate on a session connection including the effects of expected user-attributable delays (for example non-continuous SSDU input, receiving SS-user flow control). Thus, the sequence of SSDUs in the calculation above are defined to be presented at a rate which includes components representing "average" user delays.

It is possible for either the input or the output of a sequence of SSDUs to be excessively delayed by the SS-users. Such occurrences are excluded in calculating "average throughput" values.

For each direction of transfer, and for each of the "maximum throughput" and "average throughput" specifications, the throughput QOS for a particular session connection is negotiated between the SS-users and the SS-provider (see 10.2).

Throughput on a session connection relates only to the transfer of normal data and typed data over the session connection. There is no specification of the throughput for data which is transferred in association with the issue of any other session service primitives (for example S-CONNECT, S-CAPABILITY-DATA, etc.).

10.3.4 Transit delay

Transit delay is the elapsed time between the completion of any session service request primitive and the corresponding session service indication primitive occurring during the data transfer phase of a session connection. Elapsed time values are calculated only on service primitive pairs which are successfully completed.

Successful completion of a service primitive pair is defined to occur when the issue of the request primitive by one SS-user results in the issue of the corresponding indication primitive to the peer user (including any SS-user data associated with the primitive) which is without error, and in a proper sequence with respect to other primitives, prior to release of the session connection by the receiving SS-user.

In duplex and half-duplex session connections, transit delay is specified independently for each direction of transfer. In general, each transit delay specification defines both the average value and the maximum value expected for a session connection. Each specification of transit delay assumes a previously stated average size for SS-user data included in the service primitive pair.

An attempt to measure the transit delay for an individual service primitive pair may be greatly influenced if the receiving SS-user exercises flow control. Such occurrences are excluded in calculating both average and maximum transit delay values.

10.3.5 Residual error rate

Residual error rate is the ratio of total incorrect, lost, and duplicate units of SS-user data to the total units of SS-user data transferred across the session service boundary in association with any SS-primitive issued in the data transfer phase of a session connection during a measurement period. The relationship between these quantities for a particular SS-user pair is defined in figure 4.

10.3.6 Transfer failure probability

Transfer failure probability is the ratio of total transfer failures to total transfer samples observed during a performance measurement.

A transfer sample is a discrete observation of SS-provider performance in handling service requests made by the SS-user. A transfer sample begins with the initiation of session service requests during the data transfer phase and continues until the outcome of a given number of service requests have been determined. These service requests may include the transfer of SS-user data or other service requests (such as S-ACTIVITY-START request, S-TOKEN-PLEASE request, etc.) made by the SS-user. A transfer sample will normally correspond to the duration of an individual session connection.

A transfer failure is a transfer sample in which the observed performance is worse than a specified minimum acceptable level. Transfer failures are identified by comparing the measured values for applicable supported performance parameters with specified transfer failure thresholds. The three supported performance parameters which may apply are throughput, transit delay, and residual error rate.

In systems where session service QOS is reliably monitored by the SS-provider, transfer failure probability can be estimated by the probability of an S-P-ABORT or an S-P-EXCEPTION-REPORT during a transfer sample.

10.3.7 Session connection release delay

Session connection release delay is the maximum acceptable delay between an SS-user initiated S-U-ABORT request and the successful release of a particular session connection. Session connection release delay is normally specified independently for each SS-user.

Issue of an S-U-ABORT request by either SS-user marks the beginning of the session connection release delay for both users. Successful release for one SS-user is defined to occur when that SS-user is first able to initiate a new session connection. Successful release is signalled to the SS-user not initiating the S-U-ABORT request by an S-U-ABORT indication. The SS-user initiating the S-U-ABORT request will normally receive a similar signal of local significance.

10.3.8 Session connection release failure probability

Session connection release failure probability is the ratio of total SS-user initiated abort requests resulting in session connection release failure to total SS-user initiated abort requests included in a measurement sample. Session connection release failure probability is normally specified independently for each SS-user.

Session connection release failure is defined to occur, for a particular SS-user, if that SS-user is not successfully released (as defined in 10.3.7) within the specified maximum session connection release delay as a result of error or excessive delay on the part of the SS-provider. Session connection release attempts which fail as a result of error, release refusal, or excessive delay on the part of an SS-user are excluded in calculating session connection release failure probability.

10.3.9 Session connection protection

Session connection protection is the extent to which an SS-provider attempts to prevent unauthorized monitoring or manipulation of SS-user originated information. Session connection protection is specified qualitatively by selecting one of the following session connection protection options:

- a) no protection features;
- b) protection against passive monitoring;
- c) protection against modification, replay, addition or deletion;
- d) both b) and c).

10.3.10 Priority

The specification of priority is concerned with the relationship between session connections. This parameter specifies the relative importance of a session connection with respect to

- a) the order in which session connections are to have their QOS degraded, if necessary; and
- b) the order in which session connections are to be broken to recover resources, if necessary.

This parameter only has meaning in the context of some management entity or structure able to judge relative importance. The number of priority levels is limited.

10.3.11 Session connection resilience

Session connection resilience parameters specify the probability of

- a) an SS-provider initiated non-orderly release of a session connection (i.e. issue of an S-P-ABORT indication); and
- b) an SS-provider exception report (i.e. issue of an S-P-EXCEPTION-REPORT indication) during a specified time interval on an established session connection.

10.3.12 Extended control parameter

The extended control parameter allows the SS-users to make use of the resynchronize, abort, activity interrupt and activity discard services when normal flow is congested.

NOTE — When the expedited data functional unit has been selected the extended control QOS is always provided to the SS-users.

10.3.13 Optimized dialogue transfer

The optimized dialogue transfer QOS parameter permits the concatenated transfer of certain session service requests. How this concatenation of service requests is achieved is a local implementation matter.

NOTE — This QOS parameter invokes the SS-provider extended concatenation protocol option.

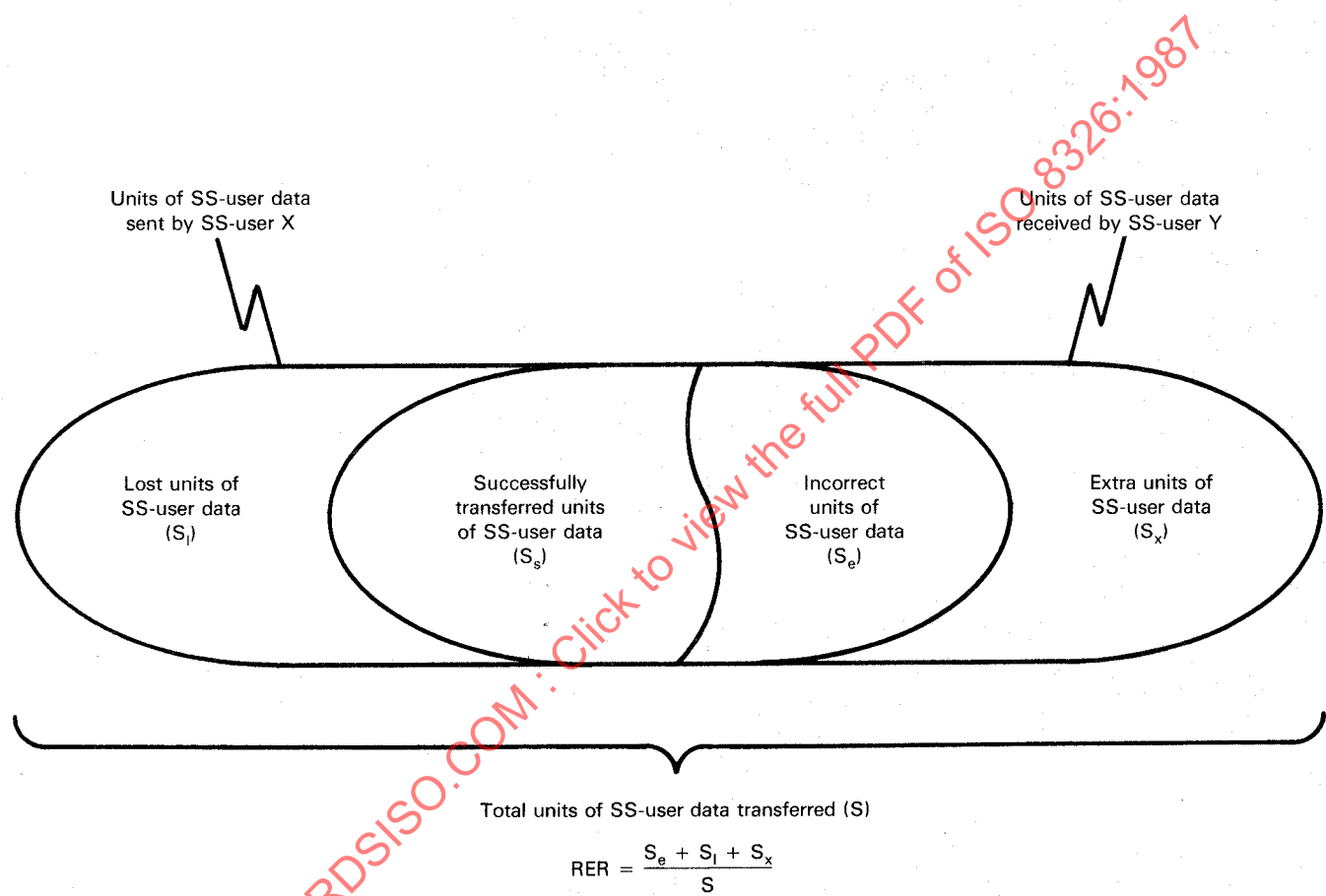


Figure 4 — Components of residual error rate

Section two : Session service primitives

11 Introduction to session service primitives

11.1 Summary of primitives

Each of the services constituting the session service is achieved by invoking a sequence of session service primitives. Tables 5, 6 and 7 summarize the primitives and their parameters occurring in each phase of the session service. The parameters are defined in clauses 12, 13 and 14.

11.2 Token restrictions on sending primitives

Table 8 defines the conditions under which those service primitives requiring tokens may be issued.

11.3 Sequencing of primitives

All SS-user requests and responses are delivered by the SS-provider in the order in which they are submitted by the SS-user, except for the following:

- a) S-EXPEDITED-DATA;
- b) S-RESYNCHRONIZE;
- c) S-ACTIVITY-INTERRUPT;
- d) S-ACTIVITY-DISCARD;
- e) S-U-ABORT,

which may be delivered earlier than previously submitted primitives, but not later than subsequently submitted primitives.

Table 5 — Session connection establishment phase primitives

Service	Primitives	Parameters
Session Connection	S-CONNECT request	Session Connection Identifier, Calling/Called/Responding Session Addresses, Result, QOS, Session Requirements, Synchronization Point Serial Number, Initial Assignment of Tokens, SS-user data
	S-CONNECT indication	
	S-CONNECT response	
	S-CONNECT confirm	

Table 6 — Data transfer phase primitives

Service	Primitives	Parameters
Normal data transfer	S-DATA request S-DATA indication	SS-user data
Expedited data transfer	S-EXPEDITED-DATA request S-EXPEDITED-DATA indication	SS-user data
Typed data transfer	S-TYPED-DATA request S-TYPED-DATA indication	SS-user data
Capability data exchange	S-CAPABILITY-DATA request S-CAPABILITY-DATA indication S-CAPABILITY-DATA response S-CAPABILITY-DATA confirm	SS-user data
Give tokens	S-TOKEN-GIVE request S-TOKEN-GIVE indication	Tokens
Please tokens	S-TOKEN-PLEASE request S-TOKEN-PLEASE indication	Tokens, SS-user data
Give control	S-CONTROL-GIVE request S-CONTROL-GIVE indication	No parameters
Minor synchro- nization point	S-SYNC-MINOR request S-SYNC-MINOR indication S-SYNC-MINOR response S-SYNC-MINOR confirm	Type, Synchronization Point Serial Number, SS-user data
Major synchro- nization point	S-SYNC-MAJOR request S-SYNC-MAJOR indication S-SYNC-MAJOR response S-SYNC-MAJOR confirm	Synchronization Point Serial Number, SS-user data
Resynchro- nize	S-RESYNCHRONIZE request S-RESYNCHRONIZE indication S-RESYNCHRONIZE response S-RESYNCHRONIZE confirm	Resynchronize Type, Synchronization Point Serial Number, Tokens, SS-user data
P-exception report	S-P-EXCEPTION-REPORT indication	Reason
U-exception reporting	S-U-EXCEPTION-REPORT request S-U-EXCEPTION-REPORT indication	Reason, SS-user data
Activity start	S-ACTIVITY-START request S-ACTIVITY-START indication	Activity Identifier, SS-user data
Activity resume	S-ACTIVITY-RESUME request S-ACTIVITY-RESUME indication	Activity Identifier, Old Activity Identifier, Synchronization Point Serial Number, Old Session Connection Identifier, SS-user data
Activity interrupt	S-ACTIVITY-INTERRUPT request S-ACTIVITY-INTERRUPT indication S-ACTIVITY-INTERRUPT response S-ACTIVITY-INTERRUPT confirm	Reason
Activity discard	S-ACTIVITY-DISCARD request S-ACTIVITY-DISCARD indication S-ACTIVITY-DISCARD response S-ACTIVITY-DISCARD confirm	Reason
Activity end	S-ACTIVITY-END request S-ACTIVITY-END indication S-ACTIVITY-END response S-ACTIVITY-END confirm	Synchronization Point Serial Number, SS-user data

Table 7 — Session connection release phase primitives

Service	Primitives	Parameters
Orderly release	S-RELEASE request	Result, SS-user data
	S-RELEASE indication	
	S-RELEASE response	
	S-RELEASE confirm	
U-abort	S-U-ABORT request	SS-user data
	S-U-ABORT indication	
P-abort	S-P-ABORT indication	Reason

Table 8 — Token restrictions on service primitives

Service primitives	data token	synchronize-minor token	major/activity token	release token
S-RELEASE request	2	2	2	2
S-RELEASE response (negative)	nr	nr	nr	0
S-DATA request (Half-duplex)	1	nr	nr	nr
S-DATA request (Duplex)	3	nr	nr	nr
S-CAPABILITY-DATA request	2	2	1	nr
S-TOKEN-GIVE request (data token)	1	nr	nr	nr
S-TOKEN-GIVE request (synchronize-minor token)	nr	1	nr	nr
S-TOKEN-GIVE request (major/activity token)	nr	nr	1	nr
S-TOKEN-GIVE request (release token)	nr	nr	nr	1
S-TOKEN-PLEASE request (data token)	0	nr	nr	nr
S-TOKEN-PLEASE request (synchronize-minor token)	nr	0	nr	nr
S-TOKEN-PLEASE request (major/activity token)	nr	nr	0	nr
S-TOKEN-PLEASE request (release token)	nr	nr	nr	0
S-CONTROL-GIVE request	2	2	1	2
S-SYNC-MINOR request	2	1	nr	nr
S-SYNC-MAJOR request	2	2	1	nr
S-U-EXCEPTION-REPORT request	0	nr	nr	nr
S-ACTIVITY-START request	2	2	1	nr
S-ACTIVITY-RESUME request	2	2	1	nr
S-ACTIVITY-INTERRUPT request	nr	nr	1	nr
S-ACTIVITY-DISCARD request	nr	nr	1	nr
S-ACTIVITY-END request	2	2	1	nr

Key:

- 0: Token available and not assigned to the SS-user who initiated the service primitive
1: Token available and assigned to the SS-user who initiated the service primitive
2: Token not available or token assigned to the SS-user who initiated the service primitive
3: Token not available
nr: No restriction

11.4 Synchronization point serial number management

Certain primitives carry a synchronization point serial number, which is used to identify a synchronization point. Synchronization points are assigned valid synchronization point serial numbers in the range 0 to 999 998 by the SS-provider. It is the responsibility of the SS-user to ensure that the number assigned by the SS-provider in a synchronization point request does not exceed 999 998.

The synchronization point serial number 999 999 is also a valid synchronization point serial number for use by the SS-user, but only in the following services, which require the synchronization point serial number of the next synchronization point:

- a) Session Connection Service;
- b) Resynchronization Service.

The management of synchronization point serial numbers is defined in this International Standard in terms of

- c) operations on abstract local variables $V(M)$, $V(A)$, $V(R)$ and Vsc , managed by the SS-provider, and
- d) primitives issued by the SS-user in order to invoke these operations.

These operations are summarized in table 33 in annex A.

11.4.1 Variables

11.4.1.1 $V(A)$

$V(A)$ is the lowest serial number to which a synchronization point confirmation is expected. No confirmation is expected when $V(A) = V(M)$.

11.4.1.2 $V(M)$

$V(M)$ is the next serial number to be used.

11.4.1.3 $V(R)$

$V(R)$ is the lowest serial number to which resynchronization restart is permitted.

11.4.1.4 Vsc

Vsc is used to determine whether or not the SS-user has the right to send minor synchronization point responses. Vsc has the following values:

- $Vsc = \text{true}$: the SS-user has the right to issue minor synchronization point responses when $V(A)$ is less than $V(M)$;
- $Vsc = \text{false}$: the SS-user does not have the right to issue minor synchronization point responses.

11.4.2 Session connection establishment

When a session connection is established in which at least one of the following functional units has been selected:

- a) minor synchronize functional unit; or
- b) major synchronize functional unit; or
- c) resynchronize functional unit

and the activity management functional unit has not been selected, $V(M)$ and $V(A)$ are set to the initial synchronization point serial number of the response/confirm primitives. $V(R)$ is set to zero. Vsc is set false.

11.4.3 Minor synchronization point

When an S-SYNC-MINOR request is issued, the associated synchronization point serial number, which is indicated to the SS-user, is equal to $V(M)$. $V(R)$ remains unchanged. $V(A)$ is set to $V(M)$ if Vsc is true, otherwise it remains unchanged. $V(M)$ is then incremented by one and Vsc is set to false.

When an S-SYNC-MINOR indication is received, the associated synchronization point serial number, which is indicated to the SS-user, is equal to $V(M)$. $V(R)$ remains unchanged. $V(A)$ is set to $V(M)$ if Vsc is false, otherwise it remains unchanged. $V(M)$ is then incremented by one and Vsc is set to true.

When an S-SYNC-MINOR response is issued, Vsc shall be true and the associated synchronization point serial number, which is supplied by the SS-user, shall be less than $V(M)$ and equal to or greater than $V(A)$. $V(A)$ is set to the serial number plus one. $V(M)$, $V(R)$ and Vsc remain unchanged.

When an S-SYNC-MINOR confirm is received, Vsc is false and the associated synchronization point serial number, which is indicated to the SS-user, is less than $V(M)$ and equal to or greater than $V(A)$. $V(A)$ is set to the serial number plus one. $V(M)$, $V(R)$ and Vsc remain unchanged.

11.4.4 Major synchronization point

When an S-SYNC-MAJOR request is issued, the associated synchronization point serial number, which is indicated to the SS-user, is equal to $V(M)$. $V(R)$ remains unchanged. $V(A)$ is set to $V(M)$ if Vsc is true, otherwise it remains unchanged. $V(M)$ is then incremented by one and Vsc is set to false.

When an S-SYNC-MAJOR indication is received, the associated synchronization point serial number, which is indicated to the SS-user, is equal to $V(M)$. $V(R)$ and Vsc remain unchanged. $V(A)$ is set to $V(M)$ if Vsc is false, otherwise it remains unchanged. $V(M)$ is then incremented by one.

When an S-SYNC-MAJOR response is issued, the associated synchronization point serial number is equal to $V(M)$ minus one. No synchronization point serial number is passed with this primitive. $V(A)$ and $V(R)$ are set to $V(M)$. $V(M)$ and Vsc remain unchanged.

When an S-SYNC-MAJOR confirm is received the associated synchronization point serial number is equal to $V(M)$ minus one. No synchronization point serial number is passed with this primitive. $V(A)$ and $V(R)$ are set to $V(M)$. $V(M)$ and Vsc remain unchanged.

11.4.5 Resynchronization

When an S-RESYNCHRONIZE request is issued

- a) if the option is "abandon", there is no associated synchronization point serial number;
- b) if the option is "restart", the associated synchronization point serial number, which is supplied by the SS-user, shall be greater than or equal to V(R) and less than or equal to V(M);
- c) if the option is "set", the associated synchronization point serial number, which is supplied by the SS-user, may have any valid value.

For all options, V(A), V(M), V(R) and Vsc remain unchanged.

When an S-RESYNCHRONIZE indication is received

- d) if the option is "abandon", the associated synchronization point serial number, which is indicated to the SS-user, is greater than or equal to V(M). V(M) is set to the serial number contained in the indication;
- e) if the option is "restart", the associated synchronization point serial number, which is indicated to the SS-user, is greater than or equal to V(R). If the synchronization point serial number is greater than V(M) (see the note), the SS-user either responds to the S-RESYNCHRONIZE indication [see g)] or generates a collision (see clause 16);

NOTE — This situation can arise if the extended control QOS is provided and the S-RESYNCHRONIZE request caused an earlier S-SYNC-MINOR request to be discarded by the SS-provider.

- f) if the option is "set", the associated synchronization point serial number, which is indicated to the SS-user, may have any valid value.

For all options, V(A), V(R) and Vsc remain unchanged. For the "restart" and "set" options, V(M) remains unchanged.

When an S-RESYNCHRONIZE response is issued

- g) if the option is "abandon" or "restart", the associated synchronization point serial number, which is supplied by the SS-user, shall be equal to the value received in the S-RESYNCHRONIZE indication;
- h) if the option is "set", the associated synchronization point serial number, which is supplied by the SS-user, may have any valid value.

V(A) and V(M) are set to the synchronization point serial number and Vsc remains unchanged. V(R) is set to zero for the options "abandon" and "set"; it remains unchanged for the "restart" option.

When an S-RESYNCHRONIZE confirm is received

- i) if the option is "abandon", the associated synchronization point serial number, which is indicated to the SS-user, is greater than or equal to V(M);
- j) if the option is "restart", the associated synchronization point serial number, which is indicated to the SS-user, is equal to the synchronization point serial number in the corresponding request;

- k) if the option is "set", the associated synchronization point serial number, which is indicated to the SS-user, may have any valid value.

V(A) and V(M) are set to the synchronization point serial number and Vsc remains unchanged. V(R) is set to zero for the options "abandon" and "set"; it remains unchanged for the "restart" option.

11.4.6 Activity management

When an S-ACTIVITY-START request is issued, or when an S-ACTIVITY-START indication is received, V(A), V(M) and V(R) are set to one and Vsc remains unchanged.

When an S-ACTIVITY-RESUME request is issued, or when an S-ACTIVITY-RESUME indication is received, V(A) and V(M) are set to the synchronization point serial number supplied by the SS-user plus one; V(R) is set to one and Vsc remains unchanged.

The management of V(A), V(M), V(R) and Vsc for S-ACTIVITY-END request, indication, response and confirm is identical to that for S-SYNC-MAJOR request, indication, response and confirm respectively.

The use of S-ACTIVITY-DISCARD and S-ACTIVITY-INTERRUPT primitives has no implication on V(A), V(M), V(R) and Vsc.

12 Session connection establishment phase

12.1 Session connection service

12.1.1 Function

The session connection service enables two SS-users to establish a session connection between themselves.

Simultaneous attempts by both SS-users to establish a session connection between themselves may result in two session connections. An SS-user may always reject an unwanted connection. No architectural restrictions are placed on the number of concurrent session connections between two SS-users.

This service allows the SS-users to exchange the values of session connection parameters. By the end of the session connection establishment phase the SS-users have agreed on a set of parameter values concerning the session connection.

12.1.2 Types of primitives and their parameters

Table 9 specifies the types of session service primitives and parameters needed for session connection establishment.

Table 9 — Session connection establishment primitives and parameters

Parameter	Primitive	S-CONNECT			
		Request	Indication	Response	Confirm
Session Connection Identifier		U	C(=)	U	C(=)
Calling Session Address		M	M		
Called Session Address		M	M		
Responding Session Address				M	M
Result				M	M(=)
Quality of Service	M		M	M	M
Session Requirements	M		M(=)	M	M(=)
Initial Synchronization Point Serial Number	C		C(=)	C	C(=)
Initial Assignment of Tokens	C		C(=)	C	C(=)
SS-user data	U		C(=)	U	C(=)

Key:
M: presence of the parameter is mandatory
C: presence of the parameter is conditional
U: presence of the parameter is a user option
Blank: the parameter is absent
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

12.1.2.1 Session Connection Identifier is a parameter which is provided by the SS-users to enable them to identify the session connection. The Session Connection Identifier is transparent to the SS-provider. This parameter consists of

- a) Calling SS-user Reference (request and indication only) with a maximum of 64 octets;
- b) Called SS-user Reference (response and confirm only) with a maximum of 64 octets;
- c) Common Reference with a maximum of 64 octets;
- d) Additional Reference Information with a maximum of 4 octets.

12.1.2.2 Calling Session Address is the session address of the calling entity (see ISO 7498/Add.3).

12.1.2.3 Called Session Address is the session address of the called entity (see ISO 7498/Add.3).

12.1.2.4 Responding Session Address is the session address of the responding entity (see ISO 7498/Add.3).

12.1.2.5 Result is a parameter indicating the success or failure of the connection establishment request. Its value can be one of

- a) accept;
- b) reject by called SS-user, where the reason for failure in the result parameter is one of
 - 1) reason not specified;
 - 2) rejection by called SS-user due to temporary congestion;
 - 3) rejection by called SS-user. The user data field may be used to provide further information;
- c) reject by SS-provider where the reason of failure in the result parameter is one of
 - 1) reason not specified;
 - 2) SS-provider congestion;
 - 3) Called Session Address unknown;
 - 4) called SS-user not attached to SSAP;

Reasons 3) and 4) may be regarded as persistent.

Only value a) or b) can be present in a response. Any of the values may be present in a confirm.

12.1.2.6 Quality of Service is a list of parameters which are defined and negotiated as described in clause 10.

12.1.2.7 Session Requirements is a list of functional units subject to the restrictions defined in 9.2 and are chosen from

- a) the half-duplex functional unit;
- b) the duplex functional unit;
- c) the exceptions functional unit;
- d) the typed data functional unit;
- e) the negotiated release functional unit;
- f) the minor synchronize functional unit;
- g) the major synchronize functional unit;
- h) the resynchronize functional unit;
- i) the expedited data functional unit;
- j) the activity management functional unit;
- k) the capability data exchange functional unit.

The session requirements specified in the response indicate the called SS-user session requirements to the requestor. The acceptor may not propose both the half-duplex and the duplex functional units in the response. If only one of the half-duplex or duplex functional units was proposed in the indication, then the acceptor proposes the same functional unit in the response or refuses the connection. If the capability data exchange functional unit is proposed, the activity management functional unit is also proposed. If the exceptions functional unit is proposed, the half-duplex functional unit is also proposed. With these exceptions, additional SS-user session requirements which were not included in the indication may be included in the response. SS-user session requirements that are proposed in both the indication and the response are the ones selected for use on the session connection.

12.1.2.8 Initial Synchronization Point Serial Number identifies the initial synchronization point. The conditions for its presence and rules for its negotiation are defined in 7.6.3. Its value is in the range 0 to 999 999.

12.1.2.9 Initial Assignment of Tokens is a list of the initial sides to which the available tokens are assigned. The parameter is only required if the corresponding tokens are available. For each available token, the value in a request/indication may be one of

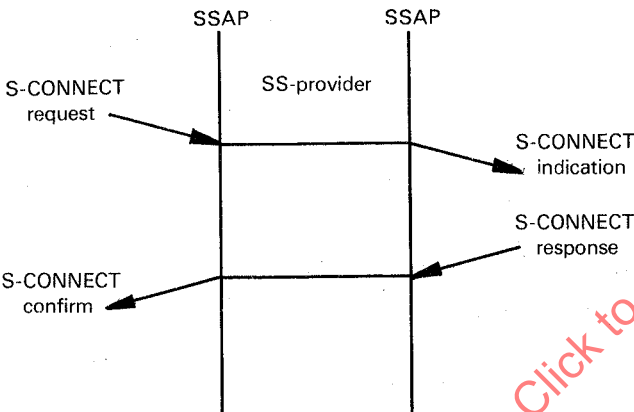
- a) requestor side;
- b) acceptor side;
- c) acceptor chooses.

The parameter in a response/confirm is absent, unless the value in the request/indication is c), in which case the acceptor replies with a) or b).

12.1.2.10 SS-user data is a parameter containing 1 to 512 octets of user information.

12.1.3 Sequence of primitives

The sequence of primitives for session connection establishment, whether accepted or rejected, is defined by the following time sequence diagram.



13 Data transfer phase

13.1 Normal data transfer service

13.1.1 Function

The normal data transfer service allows both SS-users to transfer NSSDUs over the session connection. The SS-provider should deliver each NSSDU to the SS-user as soon as possible. This service is always available on every session connection.

Use of this service is subject to the token restrictions specified in table 8.

13.1.2 Types of primitives and their parameters

Table 10 specifies the types of session service primitives and parameters needed for normal data transfer.

Table 10 — Normal data transfer primitives and parameters

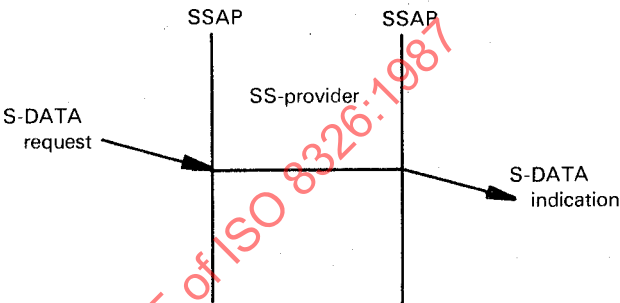
Primitive Parameter	S-DATA	
	Request	Indication
SS-user data	M	M(=)

Key:
M: presence of the parameter is mandatory
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

SS-user data parameter is an NSSDU. The size of an NSSDU is an integral number of octets greater than zero and unlimited in length.

13.1.3 Sequence of primitives

The sequence of primitives in a successful normal data transfer is defined by the following time sequence diagram.



13.2 Expedited data transfer service

13.2.1 Function

The expedited data transfer service allows SS-users to transfer XSSDUs over the session connection. The transfer of an XSSDU is free from the token and flow control constraints of the normal data transfer service, typed data transfer service and the capability data exchange service.

The SS-provider guarantees that an XSSDU will not be delivered after any subsequently submitted NSSDU or TSSDU on that session connection. The size of an XSSDU is limited.

13.2.2 Types of primitives and their parameters

Table 11 specifies the types of session service primitives and parameters needed for expedited data transfer.

Table 11 — Expedited data transfer primitives and parameters

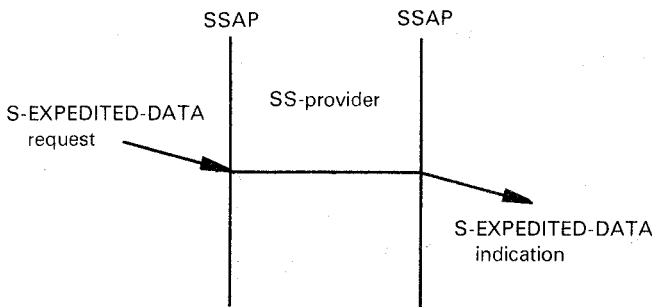
Primitive Parameter	S-EXPEDITED-DATA	
	Request	Indication
SS-user data	M	M(=)

Key:
M: presence of the parameter is mandatory
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

SS-user data parameter is an XSSDU. The size of an XSSDU is 1 to 14 octets.

13.2.3 Sequence of primitives

The sequence of primitives in a successful expedited data transfer is defined by the following time sequence diagram.



13.3 Typed data transfer service

13.3.1 Function

The typed data transfer service permits the SS-users to transfer TSSDUs over the session connection. Typed data transfers are subject to the same service restrictions as normal data transfers, except that typed data transfers are not subject to token restrictions.

13.3.2 Types of primitives and their parameters

Table 12 specifies the types of session service primitives and parameters needed for the typed data transfer service.

Table 12 — Typed data primitives and parameters

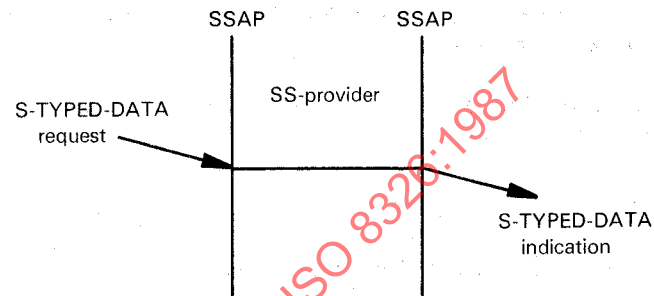
Primitive Parameter	S-TYPED-DATA	
	Request	Indication
SS-user data	M	M(=)

Key:
M: presence of the parameter is mandatory
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

SS-user data parameter is a TSSDU. The size of an integral number of octets greater than zero and unlimited in length.

13.3.3 Sequence of primitives

The sequence of primitives in a successful typed data transfer is defined by the following time sequence diagram.



13.4 Capability data exchange service

13.4.1 Function

The capability data exchange service allows SS-users to exchange a limited amount of user data while not within an activity. The service can only be initiated if activity services are available but no activity is in progress. Use of this service is subject to the token restrictions specified in table 8.

13.4.2 Types of primitives and their parameters

Table 13 specifies the types of session service primitives and parameters needed for the capability data exchange service.

SS-user data is a parameter containing 1 to 512 octets of user information.

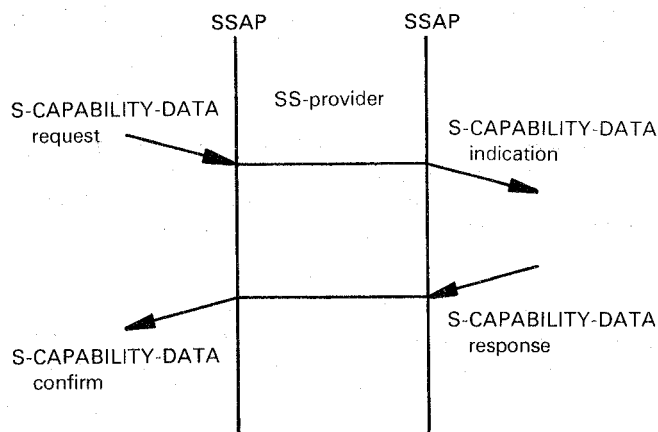
Table 13 — Capability data exchange primitives and parameters

Primitive Parameter	S-CAPABILITY-DATA			
	Request	Indication	Response	Confirm
SS-user data	U	C(=)	U	C(=)

Key:
C: presence of the parameter is conditional
U: presence of the parameter is a user option
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.4.3 Sequence of primitives

The sequence of primitives in a successful capability data exchange is defined by the following time sequence diagram.



13.5 Give tokens service

13.5.1 Function

The give tokens service allows an SS-user to surrender one or more tokens to the other SS-user, subject to the token restrictions specified in table 8.

The initial assignment of the tokens is established when the session connection is established (see 7.6.2).

13.5.2 Types of primitives and their parameters

Table 14 specifies the types of session service primitives and parameters needed for the give tokens service.

Table 14 — Give tokens primitives and parameters

Primitive Parameter	S-TOKEN-GIVE	
	Request	Indication
Tokens	M	M(=)

Key:

M: presence of the parameter is mandatory

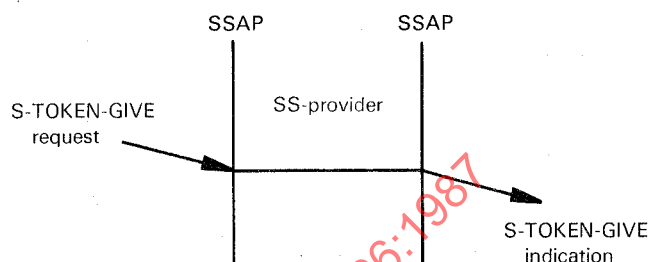
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

Tokens is a list of tokens assigned to this SS-user to be transferred to the other user. The value is any combination of

- data token;
- synchronize-minor token;
- major/activity token;
- release token.

13.5.3 Sequence of primitives

The sequence of primitives in a successful transfer of tokens is defined by the following time sequence diagram.



13.6 Please tokens service

13.6.1 Function

The please tokens service allows an SS-user to request specific tokens, subject to the token restrictions specified in table 8.

13.6.2 Types of primitives and their parameters

Table 15 specifies the types of session service primitives and parameters needed for the please tokens service.

Table 15 — Please tokens primitives and parameters

Primitive Parameter	S-TOKEN-PLEASE	
	Request	Indication
Tokens	M	M(=)
SS-user data	U	C(=)

Key:

M: presence of the parameter is mandatory

C: presence of the parameter is conditional

U: presence of the parameter is a user option

(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

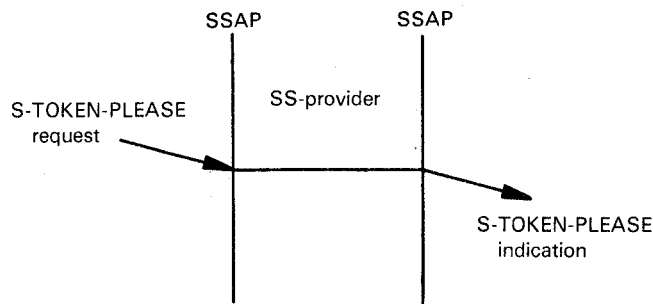
13.6.2.1 Tokens is a list of available tokens not assigned to but requested by the SS-user. The value is any combination of

- data token;
- synchronize-minor token;
- major/activity token;
- release token.

13.6.2.2 SS-user data is a parameter containing 1 to 512 octets of user information.

13.6.3 Sequence of primitives

The sequence of primitives in a successful request for tokens is defined by the following time sequence diagram.



13.7 Give control service

13.7.1 Function

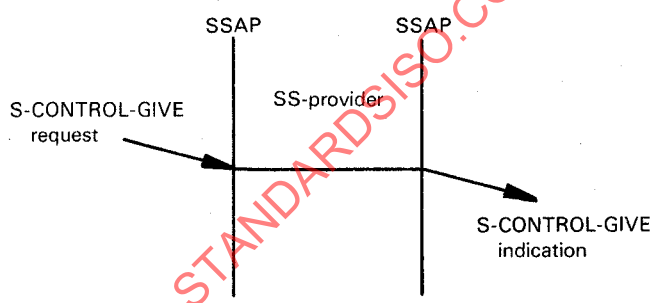
The give control service allows an SS-user to surrender the entire set of available tokens. This service is an integral part of the activity management concept. This service can only be requested when the activity management functional unit has been selected, but no activity is in progress.

13.7.2 Types of primitives and their parameters

There are no parameters associated with this service.

13.7.3 Sequence of primitives

The sequence of primitives in a successful transfer of tokens is defined by the following time sequence diagram.



13.8 Minor synchronization point service

13.8.1 Function

The minor synchronization point service allows SS-users to define minor synchronization points in the flow of NSSDUs and TSSDUs. If the activity management functional unit has been selected, this service can only be initiated within an activity. Use of this service is subject to the token restrictions specified in table 8.

The requestor may request explicit confirmation of a minor synchronization point request through the use of the Type parameter. However, the SS-provider does not require that an explicit confirmation be issued. The acceptor may issue a confirmation even if explicit confirmation is not requested.

Responses are issued in the order in which the corresponding indications were received. A further minor synchronization point request may be made while previous minor synchronization points are unconfirmed.

The confirmation of a minor or major synchronization point confirms all previously unconfirmed minor synchronization points. The number of unconfirmed minor synchronization points is not limited by the SS-provider.

Any semantics associated with request and confirmation of a minor synchronization point have no connotations to the SS-provider.

NOTE — When the duplex functional unit is selected, additional arrangements between SS-users may be required to correlate minor synchronization point requests and confirms with the flow of data from the SS-user without the synchronize-minor token.

13.8.2 Types of primitives and their parameters

Table 16 specifies the types of session service primitives and parameters needed for the minor synchronization point service.

13.8.2.1 Type is a parameter which indicates whether or not explicit confirmation is requested by the SS-user and is transparent to the SS-provider. Its value is one of

- a) explicit;
- b) optional.

13.8.2.2 Synchronization Point Serial Number is defined in 11.4.3. It is in the range 0 to 999 998.

13.8.2.3 SS-user data is a parameter containing 1 to 512 octets of user information.

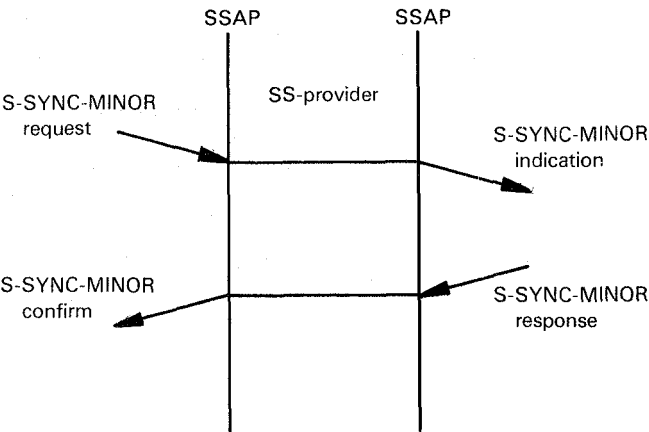
Table 16 — Minor synchronization point primitives and parameters

Parameter	Primitive	S-SYNC-MINOR			
		Request	Indication	Response	Confirm
Type		M	M(=)		
Synchronization Point Serial Number		M	M(=)	M	M(=)
SS-user data		U	C(=)	U	C(=)

Key:
M: presence of the parameter is mandatory
C: presence of the parameter is conditional
U: presence of the parameter is a user option
Blank: the parameter is absent
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.8.3 Sequence of primitives

The sequence of primitives for confirmation of a minor synchronization point is defined by the following time sequence diagram.



The response and confirm may be absent even if the Type parameter is set to explicit in the indication.

The successful confirmation of the minor synchronization point may also be achieved by issuing (instead of the S-SYNC-MINOR response to the synchronization point specified in the S-SYNC-MINOR indication):

- a) an S-SYNC-MINOR response to a subsequent S-SYNC-MINOR indication;
- b) an S-SYNC-MAJOR response to a subsequent S-SYNC-MAJOR indication;
- c) an S-SYNC-MINOR request for a subsequent minor synchronization point (provided that the synchronize-minor token has been passed from the other SS-user);
- d) an S-SYNC-MAJOR request for a subsequent major synchronization point (provided that the synchronize-minor token and, if necessary, the major/activity token have been passed from the other SS-user).

13.9 Major synchronization point service

13.9.1 Function

The major synchronization point service allows the requestor to define major synchronization points in the flow of NSSDUs, TSSDUs and XSSDUs, to completely separate the flow before and after the major synchronization point. If the activity management functional unit has been selected, this service may only be initiated within an activity. Use of this service is subject to the token restrictions specified in table 8.

After making the S-SYNC-MAJOR request, the requestor is not able to initiate any services, except for S-TOKEN-GIVE request, S-ACTIVITY-INTERRUPT request, S-ACTIVITY-DISCARD request, S-U-ABORT request or S-RESYNCHRONIZE request until the S-SYNC-MAJOR confirm is received.

After receiving the S-SYNC-MAJOR indication, in addition to any existing restrictions, the acceptor is not able to initiate S-SYNC-MAJOR request, S-SYNC-MINOR request, S-ACTIVITY-INTERRUPT request, S-ACTIVITY-DISCARD request, S-ACTIVITY-END request or S-RELEASE request until an S-SYNC-MAJOR response is issued.

Expedited data transfer services initiated by the acceptor after issuing an S-SYNC-MAJOR response are not indicated before the S-SYNC-MAJOR confirm.

13.9.2 Types of primitives and their parameters

Table 17 specifies the types of session service primitives and parameters needed for the major synchronization point service.

13.9.2.1 Synchronization Point Serial Number is defined in 11.4.4. It is in the range 0 to 999 998.

13.9.2.2 SS-user data is a parameter containing 1 to 512 octets of user information.

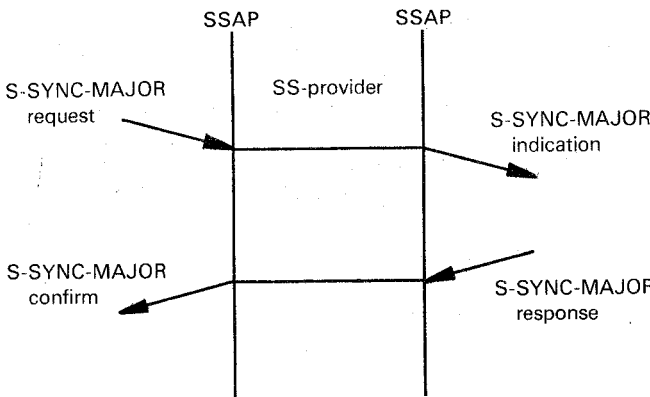
Table 17 — Major synchronization point primitives and parameters

Parameter	Primitive	S-SYNC-MAJOR			
		Request	Indication	Response	Confirm
Synchronization Point Serial Number		M	M(=)		
SS-user data		U	C(=)	U	C(=)

- Key:
- M: presence of the parameter is mandatory
 - C: presence of the parameter is conditional
 - U: presence of the parameter is a user option
 - Blank: the parameter is absent
 - (=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.9.3 Sequence of primitives

The sequence of primitives in the successful definition of a major synchronization point is defined by the following time sequence diagram.



13.10 Resynchronize service

13.10.1 Function

The resynchronize service is provided to assist orderly re-establishment of communication within the current session connection, typically following an error or lack of response by either of the SS-users or the SS-provider, or disagreements between SS-users. Requesting the service sets the session connection to an agreed defined state, including the positions of the available tokens and the value of the synchronization point serial number, which will be the next synchronization point serial number to be used.

The service may be initiated by either SS-user and has the following characteristics:

- a) after issuing the S-RESYNCHRONIZE request, the requestor is not able to initiate any services except S-U-ABORT request, until the S-RESYNCHRONIZE confirm is received;
- b) after having received an S-RESYNCHRONIZE indication, the acceptor may only issue
 - 1) S-RESYNCHRONIZE response; or
 - 2) S-RESYNCHRONIZE request (see the note); or
 - 3) S-ACTIVITY-DISCARD request (see the note); or
 - 4) S-ACTIVITY-INTERRUPT request (see the note); or
 - 5) S-U-ABORT request.

NOTE — These requests cause a collision of resynchronize requests and therefore the SS-user can only issue the request if he is going to be the collision winner (see clause 16).

- c) all undelivered data are purged;
- d) means are provided for the requesting SS-user either to set or to let the acceptor set a new assignment of each available token;
- e) means are provided to assign a new value for the synchronization point serial number;
- f) when there is an unacknowledged major synchronization point at the time of the S-RESYNCHRONIZE indication, this point remains unacknowledged. In any case, no confirmations should be issued until the resynchronization is complete and until new indications for synchronization points have been received;
- g) collision of resynchronize requests is resolved, so that only one of the colliding requests is confirmed (see clause 16).

The Resynchronize Type parameter is used to indicate the resynchronize option:

- h) “abandon” is used to request the SS-provider to resynchronize the session connection to a synchronization point which is greater than or equal to V(M). The new synchronization point serial number will be greater than any previous value used on this session connection. Where there are unacknowledged minor synchronization points at the time of the S-RESYNCHRONIZE request/indication, they remain unacknowledged;
- i) “restart” is used to return to an agreed point which is identified by a past acknowledged or unacknowledged synchronization point serial number. This point cannot be earlier than the last confirmed major synchronization point. The necessary securing of state information associated with the point is the responsibility of the SS-users;
- j) “set” is used to synchronize to any valid synchronization point serial number specified by the SS-users. When there are unacknowledged minor synchronization points at the time of the S-RESYNCHRONIZE request/indication, they remain unacknowledged.

13.10.2 Types of primitives and their parameters

Table 18 specifies the types of session service primitives and parameters needed for the resynchronize service.

Table 18 — Resynchronize primitives and parameters

Primitive	S-RESYNCHRONIZE			
	Request	Indication	Response	Confirm
Resynchronize Type	M	M(=)		M(=)
Synchronization Point Serial Number	C	M	M	M(=)
Tokens	C	C(=)	C	C(=)
SS-user data	U	C(=)	U	C(=)

Key:
M: presence of the parameter is mandatory
C: presence of the parameter is conditional
U: presence of the parameter is a user option
Blank: the parameter is absent
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.10.2.1 Resynchronize Type is a parameter which specifies one of the resynchronize options. Its value is one of

- a) abandon;
- b) restart;
- c) set.

13.10.2.2 Synchronization Point Serial Number depends on the resynchronize option and is defined in 11.4 and 11.4.5.

13.10.2.3 Tokens is a list of the available tokens for the session connection with values for their assignment following the resynchronization. For each available token, the value in a request/indication is one of

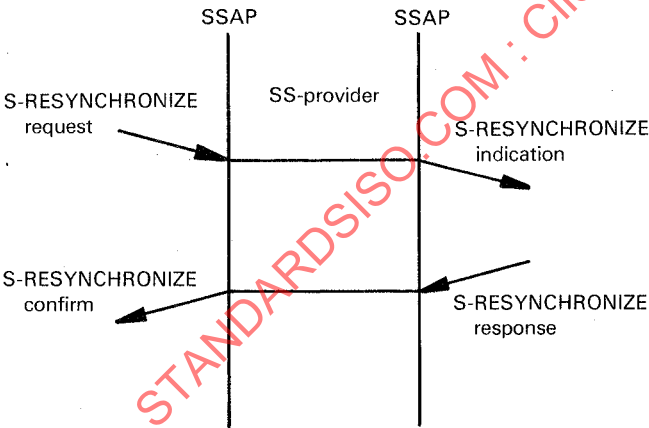
- a) requestor side;
- b) acceptor side;
- c) acceptor chooses.

The value for a response/confirm is the same as in the request/indication unless that value is c), in which case the acceptor chooses a) or b).

13.10.2.4 SS-user data is a parameter containing 1 to 512 octets of user information.

13.10.3 Sequence of primitives

The sequence of primitives in a successful resynchronization without collision is defined by the following time sequence diagram. Collision cases are defined in clause 16.



13.11 P-exception reporting service

13.11.1 Function

The P-exception reporting service permits SS-users to be notified of unanticipated situations not covered by other services. If a service cannot be completed due to SS-provider protocol errors or malfunctions, the P-exception reporting service is used to indicate this to both SS-users.

If used with the activity management service, the P-exception reporting service is only permitted while an activity is in progress or waiting for S-CAPABILITY-DATA confirm.

Following an S-P-EXCEPTION-REPORT indication, and until the error condition is cleared

- a) NSSDUs, TSSDUs and XSSDUs will be discarded by the SS-provider;
- b) synchronization point indications will not be given to the SS-users.

On receipt of an S-P-EXCEPTION-REPORT indication, either SS-user initiates one of the following services to clear the error:

- c) resynchronize;
- d) abort;
- e) activity interrupt or activity discard;
- f) give the data token (see notes).

The SS-users are not permitted to initiate any other services until the error is cleared.

NOTES

- 1 It is not recommended that the error condition be cleared by passing the data token when the resynchronize and/or activity management functional units have been selected.
- 2 If the error condition is cleared by passing the data token, data and synchronization point serial numbers may be lost. However, the SS-provider will keep track of the serial numbers of the synchronization points which have been discarded. Therefore, the synchronization point serial number indicated to the SS-user in a synchronization point request/indication made after the error condition has been cleared will reflect the fact that synchronization points have been discarded during the error condition.
- 3 XSSDUs sent after the S-TOKEN-GIVE request will be discarded if they overtake the request.
- 4 Tokens other than the data token may be transferred at the same time.

13.11.2 Types of primitives and their parameters

Table 19 specifies the types of session service primitives and parameters needed for the P-exception reporting service.

Table 19 — P-exception reporting primitives and parameters

Primitive	S-P-EXCEPTION-REPORT
Parameter	Indication
Reason	M

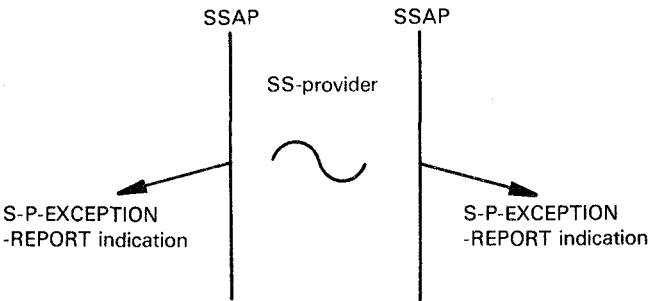
Key:
M: presence of the parameter is mandatory

Reason is a parameter specifying the reason for the exception report. Its value is one of

- a) protocol error;
- b) non-specific error.

13.11.3 Sequence of primitives

The sequence of primitives in a successful P-exception report is defined by the following time sequence diagram.



13.12 U-exception-reporting service

13.12.1 Function

The U-exception reporting service permits an SS-user to report an exception condition subject to the token restrictions specified in table 8.

If used with the activity management service, the U-exception reporting service is only permitted while an activity is in progress.

Following an S-U-EXCEPTION-REPORT request, and until the error condition is cleared

- a) NSSDUs, TSSDUs and XSSDUs will be discarded by the SS-provider;
- b) synchronization point indications will not be given to the requestor of the S-U-EXCEPTION-REPORT;
- c) the requestor is only permitted to issue S-U-ABORT request.

On receipt of an S-U-EXCEPTION-REPORT indication, the acceptor initiates one of the following services to clear the error:

- d) resynchronize;
- e) abort;
- f) activity interrupt or activity discard;
- g) give the data token (see notes 1 to 4).

The acceptor is not permitted to initiate any other services until the error is cleared.

NOTES

- 1 It is not recommended that the error condition be cleared by passing the data token when the resynchronize and/or activity management functional units have been selected.
- 2 If the error condition is cleared by passing the data token, data and synchronization point serial numbers may be lost. However, the SS-provider will keep track of the serial numbers of the synchronization points which have been discarded. Therefore, the synchronization

point serial number indicated to the SS-user in a synchronization point request/indication made after the error condition has been cleared will reflect the fact that synchronization points have been discarded during the error condition.

- 3 XSSDUs sent after the S-TOKEN-GIVE request will be discarded if they overtake the request.
- 4 Tokens other than the data token may be transferred at the same time.

13.12.2 Types of primitives and their parameters

Table 20 specifies the types of session service primitives and parameters needed for the U-exception reporting service.

Table 20 – U-exception reporting primitives and parameters

Primitive Parameter	S-U-EXCEPTION-REPORT	
	Request	Indication
Reason	M	M(=)
SS-user data	U	C(=)

Key:
M: presence of the parameter is mandatory
C: presence of the parameter is conditional
U: presence of the parameter is a user option
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

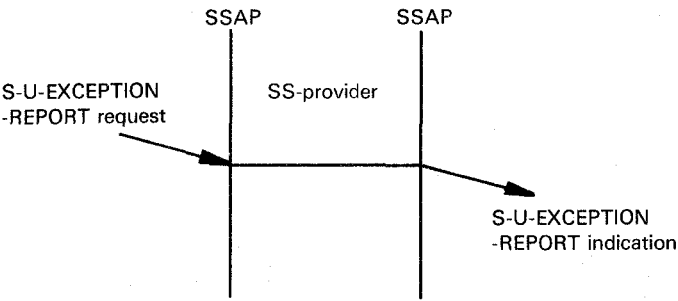
13.12.2.1 Reason is a parameter specifying the reason for the exception report and is transparent to the SS-provider. Its value is one of

- a) SS-user receiving ability jeopardized (i.e. data received may not be handled correctly);
- b) local SS-user error;
- c) sequence error;
- d) demand data token;
- e) unrecoverable procedural error;
- f) non-specific error.

13.12.2.2 SS-user data is a parameter containing 1 to 512 octets of user information.

13.12.3 Sequence of primitives

The sequence of primitives in a successful U-exception report is defined by the following time sequence diagram.



13.13 Activity start service

13.13.1 Function

The activity start service allows an SS-user to indicate that a new activity is entered. The value of the next synchronization point serial number to be used is set to one (see 11.4.6). The service can only be initiated if no activity is in progress and subject to the token restrictions specified in table 8.

13.13.2 Types of primitives and their parameters

Table 21 specifies the types of session service primitives and parameters needed for the activity start service.

Table 21 — Activity start primitives and parameters

Parameter \ Primitive	S-ACTIVITY-START	
	Request	Indication
Activity Identifier	M	M(=)
SS-user data	U	C(=)

Key:

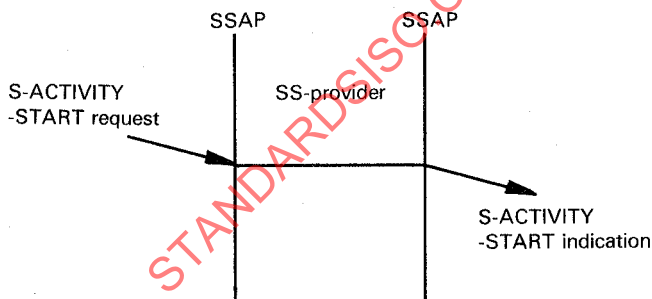
- M: presence of the parameter is mandatory
- C: presence of the parameter is conditional
- U: presence of the parameter is a user option
- (=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.13.2.1 Activity Identifier is a parameter which is provided by the SS-users to enable them to identify the new activity and is transparent to the SS-provider. This parameter has a maximum of 6 octets.

13.13.2.2 SS-user data is a parameter containing 1 to 512 octets of user information.

13.13.3 Sequence of primitives

The sequence of primitives in a successful activity start is defined by the following time sequence diagram.



13.14 Activity resume service

13.14.1 Function

The activity resume service allows an SS-user to indicate that a previously interrupted activity is resumed. A new activity identifier is provided by the SS-user together with the identifier of the activity being resumed and the next synchronization point serial number to be used minus one. In the case when the resumed activity was originally started on another session connection, the session connection identifier of that session connection is also provided by the SS-user.

The service can only be initiated if no activity is in progress and subject to the token restrictions specified in table 8.

13.14.2 Types of primitives and their parameters

Table 22 specifies the types of session service primitives and parameters needed for the activity resume.

Table 22 — Activity resume primitives and parameters

Parameter \ Primitive	S-ACTIVITY-RESUME	
	Request	Indication
Activity Identifier	M	M(=)
Old Activity Identifier	M	M(=)
Synchronization Point Serial Number	M	M(=)
Old Session Connection Identifier	U	C(=)
SS-user data	U	C(=)

Key:

- M: presence of the parameter is mandatory
- C: presence of the parameter is conditional
- U: presence of the parameter is a user option
- (=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.14.2.1 Activity Identifier is a parameter which is provided by the SS-users to enable them to give a new identifier to the activity being resumed and is transparent to the SS-provider. This parameter has a maximum of 6 octets.

13.14.2.2 Old Activity Identifier is the original identifier of the activity being resumed and is transparent to the SS-provider.

13.14.2.3 Synchronization Point Serial Number is provided by the SS-user and is defined in 11.4.6.

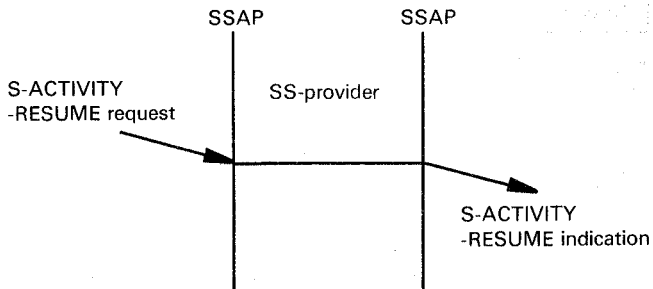
13.14.2.4 Old Session Connection Identifier is the session connection identifier of the session connection in which the activity being resumed was originally started and is transparent to the SS-provider. It consists of

- a) Calling SS-user Reference with a maximum of 64 octets;
- b) Called SS-user Reference with a maximum of 64 octets;
- c) Common Reference with a maximum of 64 octets;
- d) Additional Reference Information with a maximum of 4 octets.

13.14.2.5 SS-user data is a parameter containing 1 to 512 octets of user information.

13.14.3 Sequence of primitives

The sequence of primitives in a successful activity resume is defined by the following time sequence diagram.



13.15 Activity interrupt service

13.15.1 Function

The activity interrupt service allows an SS-user to abnormally terminate the current activity so that work achieved before the interruption is not cancelled, and may be resumed later.

The service can only be initiated if an activity is in progress and subject to the token restrictions specified in table 8. After receipt of the confirm, all available tokens are assigned to the SS-user which issued the request.

After issuing an S-ACTIVITY-INTERRUPT request, the requestor is not able to initiate any services, except S-U-ABORT request, until the S-ACTIVITY-INTERRUPT confirm is received.

After receiving an S-ACTIVITY-INTERRUPT indication, the acceptor is not able to initiate any services, except S-U-ABORT request, until the S-ACTIVITY-INTERRUPT response is issued.

Use of this service may cause loss of data which has not yet been delivered to the SS-user.

13.15.2 Types of primitives and their parameters

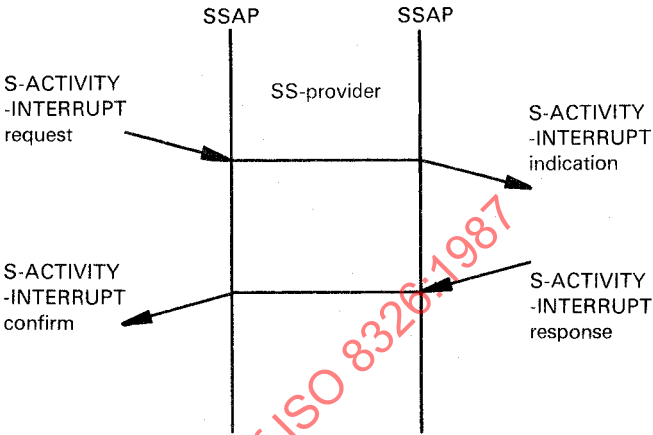
Table 23 specifies the types of session service primitives and parameters needed for the activity interrupt service.

Reason is a parameter specifying the reason for the activity interrupt and is transparent to the SS-provider. Its value is one of

- a) SS-user receiving ability jeopardized (i.e. data received may not be handled correctly);
- b) local SS-user error;
- c) sequence error;
- d) demand data token;
- e) unrecoverable procedural error;
- f) non-specific error.

13.15.3 Sequence of primitives

The sequence of primitives in a successful activity interrupt is defined by the following time sequence diagram.



13.16 Activity discard service

13.16.1 Function

The activity discard service allows an SS-user to abnormally terminate the current activity. There is an implied meaning to the SS-user that the previous content of this activity is cancelled, but this is not controlled by the SS-provider.

The service can only be initiated if an activity is in progress and subject to the token restrictions specified in table 8. After receipt of the confirm, all available tokens are assigned to the SS-user which issued the request.

After issuing an S-ACTIVITY-DISCARD request, the requestor is not able to initiate any services, except S-U-ABORT request, until the S-ACTIVITY-DISCARD confirm is received.

After receiving an S-ACTIVITY-DISCARD indication, the acceptor is not able to initiate any services, except S-U-ABORT request, until the S-ACTIVITY-DISCARD response is issued.

Use of this service may cause loss of data which has not yet been delivered to the SS-user.

Table 23 — Activity interrupt primitives and parameters

Parameter	Primitive	S-ACTIVITY-INTERRUPT			
		Request	Indication	Response	Confirm
Reason		U	C(=)		

Key:

M: presence of the parameter is mandatory

C: presence of the parameter is conditional

U: presence of the parameter is a user option

Blank: the parameter is absent

(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.16.2 Types of primitives and their parameters

Table 24 specifies the types of session service primitives and parameters needed for the activity discard service.

Reason is a parameter specifying the reason for the activity discard and is transparent to the SS-provider. Its value is one of

- a) SS-user receiving ability jeopardized (i.e. data received may not be handled correctly);
- b) local SS-user error;
- c) sequence error;
- d) demand data token;
- e) unrecoverable procedural error;
- f) non-specific error.

13.16.3 Sequence of primitives

The sequence of primitives in a successful activity discard is defined by the following time sequence diagram.

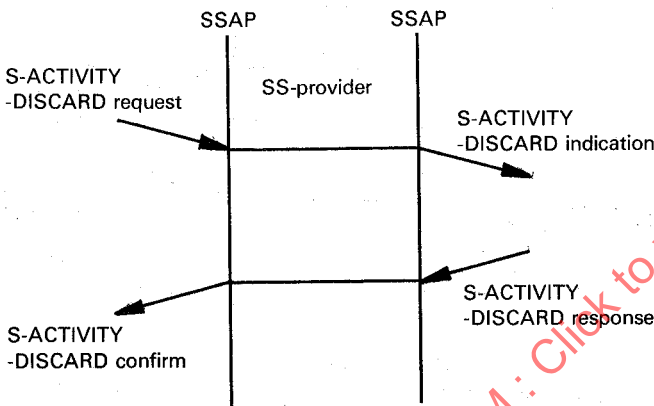


Table 24 — Activity discard primitives and parameters

Parameter \ Primitive	S-ACTIVITY-DISCARD			
	Request	Indication	Response	Confirm
Reason	U	C(=)		

- Key:
- C: presence of the parameter is conditional
 - U: presence of the parameter is a user option
 - Blank: the parameter is absent
 - (=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

Table 25 — Activity end primitives and parameters

Parameter \ Primitive	S-ACTIVITY-END			
	Request	Indication	Response	Confirm
Synchronization Point Serial Number	M	M(=)		
SS-user data	U	C(=)	U	C(=)

- Key:
- M: presence of the parameter is mandatory
 - C: presence of the parameter is conditional
 - U: presence of the parameter is a user option
 - Blank: the parameter is absent
 - (=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

13.17 Activity end service

13.17.1 Function

The activity end service allows an SS-user to indicate the end of an activity, and has the effect of setting a major synchronization point. This service can only be invoked if an activity is in progress and subject to the token restrictions specified in table 8.

After issuing the S-ACTIVITY-END request, in addition to any existing restrictions, the requestor is not able to initiate any services, except for S-U-ABORT request, S-ACTIVITY-INTERRUPT request, S-ACTIVITY-DISCARD request or S-TOKEN-GIVE request until the S-ACTIVITY-END confirm is received.

After receiving the S-ACTIVITY-END indication, in addition to any existing restrictions, the acceptor is not able to initiate S-SYNC-MAJOR request, S-SYNC-MINOR request, S-ACTIVITY-INTERRUPT request, S-ACTIVITY-DISCARD request, S-ACTIVITY-END request or S-RELEASE request until the S-ACTIVITY-END response is issued.

If the activity management functional unit has been selected, the SS-user is not allowed to initiate any services, except activity start, activity resume, token management, capability data, expedited data, typed data, normal data, release or abort, until an activity is started or resumed.

13.17.2 Types of primitives and their parameters

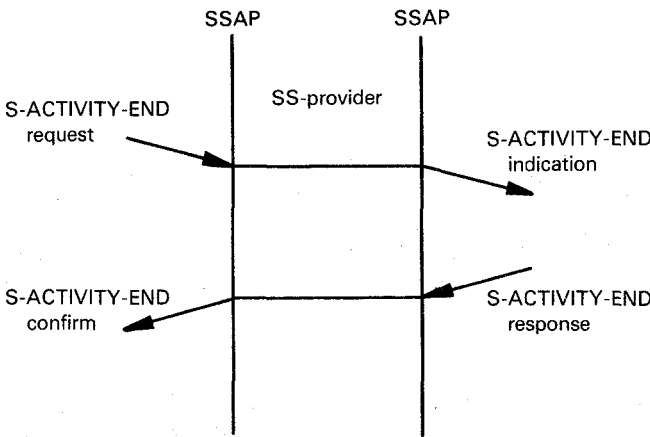
Table 25 specifies the types of session service primitives and parameters needed for the activity end service.

13.17.2.1 Synchronization Point Serial Number is defined in 11.4.6.

13.17.2.2 SS-user data is a parameter containing 1 to 512 octets of user information.

13.17.3 Sequence of primitives

The sequence of primitives in a successful normal termination of an activity is defined by the following time sequence diagram.



14 Session connection release phase

14.1 Orderly release service

14.1.1 Function

The orderly release service is always provided and allows either SS-user to release the session connection in an orderly manner. This is done co-operatively between the two SS-users without the loss of data after all in-transit data have been delivered and accepted by both SS-users.

Use of this service is subject to the token restrictions specified in table 8. If the release token is available the acceptor may refuse the release and continue the session connection without loss of data. If the release token is not available, the acceptor cannot refuse the release.

14.1.2 Types of primitives and their parameters

Table 26 specifies the types of session service primitives and parameters needed for the orderly release service.

14.1.2.1 Result is a parameter indicating whether or not the session release is granted. Its value may be one of

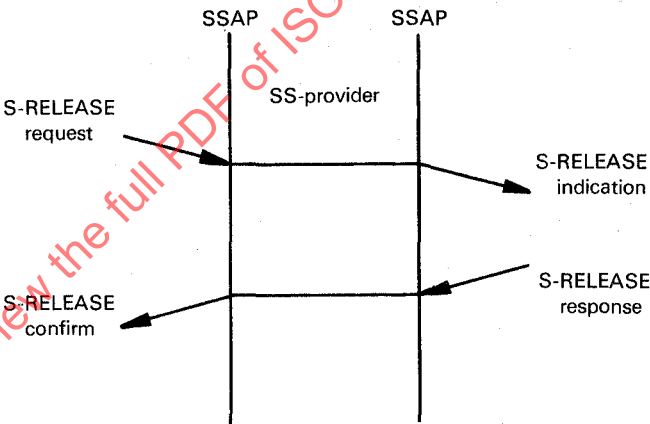
- a) affirmative;
- b) negative.

The latter value may be given only if the release token is available.

14.1.2.2 SS-user data is a parameter containing 1 to 512 octets of user information.

14.1.3 Sequence of primitives

The sequence of primitives in a successful orderly session release is defined by the following time sequence diagram.



A collision of S-RELEASE requests may occur when no tokens are available. This results in S-RELEASE indications to both SS-users. In this case, the calling SS-user should send the S-RELEASE response after receiving the S-RELEASE indication from the called SS-user. The called SS-user should not send his S-RELEASE response before receiving the S-RELEASE confirm from the calling SS-user.

14.2 U-abort service

14.2.1 Function

The U-abort service provides the means by which either SS-user can instantaneously release the session connection and have the other SS-user informed of this release. Use of this service will cause loss of undelivered data.

Table 26 — Orderly release primitives and parameters

Primitive	S-RELEASE			
	Request	Indication	Response	Confirm
Result			M	M(=)
SS-user data	U	C(=)	U	C(=)

Key:
M: presence of the parameter is mandatory
C: presence of the parameter is conditional
U: presence of the parameter is a user option
Blank: the parameter is absent
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

14.2.2 Types of primitives and their parameters

Table 27 specifies the types of session service primitives and parameters needed for the U-abort service.

Table 27 — U-abort primitives and parameters

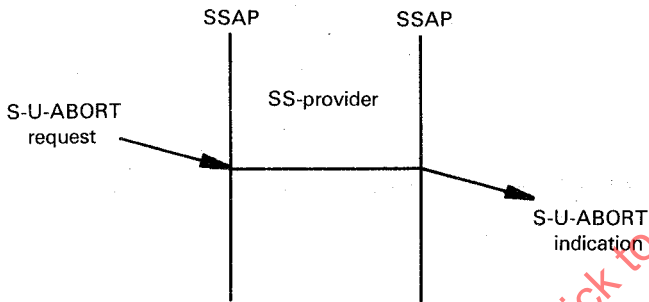
Parameter \ Primitive	S-U-ABORT	
	Request	Indication
SS-user data	U	C(=)

Key:
C: presence of the parameter is conditional
U: presence of the parameter is a user option
(=): the value of the parameter is identical to the value of the corresponding parameter of the preceding SS primitive

SS-user data is a parameter containing 1 to 9 octets of user information.

14.2.3 Sequence of primitives

The sequence of primitives in a successful U-abort is defined by the following time sequence diagram.



14.3 P-abort service

14.3.1 Function

The P-abort service provides the means by which the SS-provider may indicate the release of the session connection for reasons internal to the SS-provider. Use of this service will cause loss of undelivered data. A reason code of limited size is passed from the SS-provider to the SS-user.

14.3.2 Types of primitives and their parameters

Table 28 specifies the types of session service primitives and parameters needed for the P-abort service.

Table 28 — P-abort primitives and parameters

Parameter \ Primitive	S-P-ABORT	
	Request	Indication
Reason		M

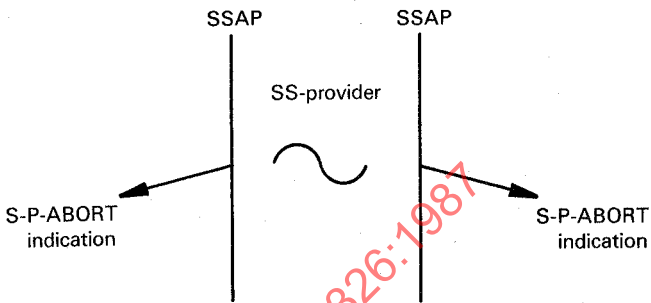
Key:
M: presence of the parameter is mandatory

Reason is a parameter indicating the reason for the abort. Its value is one of

- a) transport disconnect;
- b) protocol error;
- c) undefined.

14.3.3 Sequence of primitives

The sequence of primitives in a successful P-abort is defined by the following time sequence diagram.



15 Sequences of primitives

15.1 State tables

Annex A contains state tables which define the constraints on the sequences in which the session service primitives may occur. The constraints determine the order in which the session services occur, but do not fully specify when they may occur. Other constraints will affect the ability of an SS-user or the SS-provider to issue a primitive at any particular time.

15.2 Sequences of primitives at one session connection end-point

The possible sequences of primitives at one session connection end-point may be derived directly from the state tables in annex A.

16 Collision

16.1 Collision as viewed by the SS-user

The SS-provider resolves collisions between those requests that may destroy SS-user data. If a collision occurs, one of the SS-users will receive an unexpected indication while awaiting one of the following:

- a) S-RESYNCHRONIZE confirm;
- b) S-ACTIVITY-INTERRUPT confirm;
- c) S-ACTIVITY-DISCARD confirm;
- d) clearing the error state after issuing an S-U-EXCEPTION-REPORT request.

Table 29 defines the indications that may be received which indicate that the SS-user has lost a collision resolved by the SS-provider.

16.2 Collision resolution by the SS-provider

The SS-provider resolves colliding SS-user requests according to the following rules.

In the case of collision between two of the following types of requests, the first in the list takes precedence.

- a) S-U-ABORT request;
- b) S-ACTIVITY-DISCARD request;
- c) S-ACTIVITY-INTERRUPT request;
- d) S-RESYNCHRONIZE (abandon) request;
- e) S-RESYNCHRONIZE (set) request;

- f) S-RESYNCHRONIZE (restart) request;
- g) S-U-EXCEPTION-REPORT request.

Possible collisions of the same request are handled as follows:

- h) if two S-RESYNCHRONIZE (abandon) requests collide, the calling SS-user request takes precedence;
- i) if two S-RESYNCHRONIZE (restart) requests collide, the request with the lowest serial number takes precedence. If the serial numbers are equal, the calling SS-user request takes precedence;
- j) if two S-RESYNCHRONIZE (set) requests collide, the calling SS-user request takes precedence.

Table 29 — Indications resulting from collision resolution

SS-user receives SS-user is waiting for	ER	RR	RS	RA	AI	AD	AB
Clearing error state after S-U-EXCEPTION-REPORT request		X	X	X	X	X	X
S-RESYNCHRONIZE (restart) confirm		X	X	X	X	X	X
S-RESYNCHRONIZE (set) confirm			X	X	X	X	X
S-RESYNCHRONIZE (abandon) confirm				X	X	X	X
S-ACTIVITY-INTERRUPT confirm							X
S-ACTIVITY-DISCARD confirm							X

- Key:
- X: indication may be received
 - Blank: indication will not be received
 - AB: S-P-ABORT indication or S-U-ABORT indication
 - AD: S-ACTIVITY-DISCARD indication
 - AI: S-ACTIVITY-INTERRUPT indication
 - ER: S-U-EXCEPTION-REPORT indication or S-P-EXCEPTION-REPORT indication
 - RA: S-RESYNCHRONIZE (abandon) indication
 - RR: S-RESYNCHRONIZE (restart) indication
 - RS: S-RESYNCHRONIZE (set) indication

Annex A

State tables

(This annex forms part of the standard.)

A.1 General

This annex describes the session service in terms of state tables. The state tables show the state of an SS-user, the events that occur at the session service boundary, the actions taken by the SS-user and the resultant state.

These state tables do not constitute a formal definition of the session service; they are included to provide a more precise definition of the relationships between session service primitives defined in clauses 12, 13 and 14.

Table 30 specifies the abbreviated name and name of each incoming event generated by the SS-provider.

Table 31 specifies the abbreviated name and name of each state.

Table 32 specifies the abbreviated name and name of each outgoing event generated by the SS-user.

Table 33 summarizes the operations on the variables V(A), V(M), V(R) and Vsc.

Table 34 specifies the specific actions.

Table 35 specifies the predicates.

Tables 36 to 43 specify the state tables.

A.2 Notation for state tables

A.2.1 Incoming events, states and outgoing events are represented by their abbreviated names.

A.2.2 Specific actions are represented by the notation [n], where n is the number of the specific action in table 34.

A.2.3 Predicates are represented by the notation pn, where n is the number of the predicate in table 35.

A.2.4 Boolean operators are represented by the following notation:

&	AND
^	NOT
OR	OR

A.3 Conventions for entries in state tables

A.3.1 The intersection of each state and incoming or outgoing event which is invalid is left blank.

A.3.2 The intersection of each state and incoming or outgoing event which is valid contains entries which are either

a) an action list which

- 1) may contain specific actions;
- 2) always contains the resultant state;

or

b) one or more conditional action lists, each consisting of

- 1) a predicate expression comprising predicates and boolean operators;
- 2) an action list [(as in A.3.2a)].

NOTE — The action lists and conditional action lists use the notation in clause A.2.

A.4 Actions to be taken by the SS-user

The state tables define the action to be taken by the SS-user.

A.4.1 Invalid intersections

If the intersection of the state and an incoming or outgoing event is invalid, any action taken by the SS-user is a local matter.

A.4.2 Valid intersections

If the intersection of the state and incoming event is valid, one of the following actions shall be taken.

A.4.2.1 If the intersection contains an action list, the SS-user shall take the specific actions in the order specified in the state table.

A.4.2.2 If the intersection contains one or more conditional action lists, for each predicate expression that is true the SS-user shall take the specific actions in the order given in the action list associated with the predicate expression. If none of the predicate expressions are true, the SS-user shall take one of the actions defined in A.4.1.

A.5 Definitions of sets and variables

The following sets and variables are specified in this International Standard.

A.5.1 Functional units

The set of all functional units specified in this International Standard is defined as

$$\text{fu-dom} = \{\text{FD, HD, EXCEP, TD, NR, SY, MA, RESYN, EX, ACT, CD}\}$$

where

FD = Duplex functional unit
 HD = Half-duplex functional unit
 EXCEP = Exceptions functional unit
 TD = Typed data functional unit
 NR = Negotiated release functional unit
 SY = Minor synchronize functional unit
 MA = Major synchronize functional unit
 RESYN = Resynchronize functional unit
 EX = Expedited data functional unit
 ACT = Activity management functional unit
 CD = Capability data exchange functional unit

A boolean function FU is defined over fu-dom as follows:

for f in fu-dom

$\text{FU}(f) = \text{true}$ if and only if the functional unit f has been selected during the session connection establishment phase.

The value is set when the S-CONNECT response is issued or the S-CONNECT confirm is received.

A.5.2 Tokens

The set of all tokens specified in this International Standard is defined as

$$\text{tk-dom} = \{\text{mi, ma, tr, dk}\}$$

where

mi = synchronize-minor token
 ma = major/activity token
 tr = release token
 dk = data token

The following boolean functions are defined over tk-dom:

a) $\text{AV}(t)$, for t in tk-dom, is a function which defines the availability of the corresponding token and has the following values:

$\text{AV}(\text{mi}) = \text{FU}(\text{SY})$
 $\text{AV}(\text{dk}) = \text{FU}(\text{HD})$
 $\text{AV}(\text{tr}) = \text{FU}(\text{NR})$
 $\text{AV}(\text{ma}) = \text{FU}(\text{MA}) \text{ or } \text{FU}(\text{ACT})$

b) $\text{OWNED}(t)$, for t in tk-dom, is a function which defines the assignment of the corresponding token and is defined as

$\text{OWNED}(t) = \text{true}$: if the token is assigned to the SS-user
 $\text{OWNED}(t) = \text{false}$: if the token is not assigned to the SS-user

$\text{OWNED}(t)$ is not defined if $\text{AV}(t) = \text{false}$. $\text{OWNED}(t)$ is set when

- 1) the S-CONNECT response is issued or the S-CONNECT confirm is received; or
- 2) the S-RESYNCHRONIZE response is issued or the S-RESYNCHRONIZE confirm is received; or
- 3) the S-TOKEN-GIVE request is issued or the S-TOKEN-GIVE indication is received; or
- 4) the S-CONTROL-GIVE request is issued or the S-CONTROL-GIVE indication is received;
- 5) the S-ACTIVITY-INTERRUPT response is issued or the S-ACTIVITY-INTERRUPT confirm is received;
- 6) the S-ACTIVITY-DISCARD response is issued or the S-ACTIVITY-DISCARD confirm is received.

c) $\text{I}(t)$, for t in tk-dom, is a function which, when true, indicates that the SS-user has Initiating rights for the behaviour controlled by the token. This applies even if the corresponding token is not available:

$$\text{I}(t) = \neg \text{AV}(t) \text{ OR } \text{OWNED}(t)$$

d) $\text{A}(t)$, for t in tk-dom, is a function which, when true, indicates that the SS-user has Accepting rights for the behaviour controlled by the token. This applies even if the corresponding token is not available:

$$\text{A}(t) = \neg \text{AV}(t) \text{ OR } \neg \text{OWNED}(t)$$

e) $\text{II}(t)$, for t in tk-dom, is a function which, when true, indicates that the SS-user has Initiating rights as $\text{I}(t)$, but this applies to the case when the behaviour may only be initiated if the corresponding token is available and owned:

$$\text{II}(t) = \text{AV}(t) \text{ AND } \text{OWNED}(t)$$

f) $\text{AA}(t)$, for t in tk-dom, is a function which, when true, indicates that the SS-user has Accepting rights as $\text{A}(t)$, but only if the corresponding token is available, but not owned:

$$\text{AA}(t) = \text{AV}(t) \text{ AND } \neg \text{OWNED}(t)$$

A.5.3 SET of tokens

The following subsets of tk-dom are defined:

$\text{RT} = \{\text{tokens requested in the input event}\}$

$\text{GT} = \{\text{tokens given in the input event}\}$

For the purpose of the following function definitions, two further sets are defined:

$F = \{\text{AV, OWNED, I, A, II, AA}\}$
 (the set of functions defined in A.5.2)

$S = \text{the set of subsets of tk-dom}$

The following functions are defined over F and S:

a) $\text{ALL}(f, s)$, for f in F and s in S:

$\text{ALL}(f, s) = \text{true}$: all of the $f(t)$ for t in s are true or s is empty;

For example:

$ALL(A, tk-dom) = \text{true}$: none of the available tokens are owned (for example on receipt of an S-RELEASE indication)

b) $ANY(f, s)$, for f in F and s in S :

$ANY(f, s) = \text{true}$: any $f(t) = \text{true}$ for t in s and s is not empty;

For example:

$ANY(II, tk-dom) = \text{true}$: at least one of the available tokens is owned.

A.5.4 Variables

A.5.4.1 Vact

Vact is a boolean variable having the following values when the activity management functional unit has been selected [$FU(Act) = \text{true}$]:

Vact = true : an activity is in progress;

Vact = false: no activity is in progress;

Vact has no defined value if $FU(Act) = \text{false}$.

Vact is set as follows:

a) Vact is set false during the connection establishment phase, if the activity management functional unit has been selected [$FU(Act) = \text{true}$]. Otherwise, Vact is not set;

b) Vact is set true when the S-ACTIVITY-START request or S-ACTIVITY-RESUME request is issued or the S-ACTIVITY-START indication or S-ACTIVITY-RESUME indication is received (only possible when $FU(Act) = \text{true}$);

c) Vact is set false when the S-ACTIVITY-DISCARD response or S-ACTIVITY-INTERRUPT response is issued or the S-ACTIVITY DISCARD confirm or S-ACTIVITY-INTERRUPT confirm is received;

d) Vact is set false when the S-ACTIVITY-END response is issued or the S-ACTIVITY-END confirm is received.

A.5.4.2 Vrsp and Vrspnb

These variables are used to resolve resynchronization collisions.

Vrsp indicates what kind of resynchronization is currently in progress:

Vrsp = no: no resynchronization in progress

Vrsp = a: resynchronize abandon

Vrsp = r: resynchronize restart

Vrsp = s: resynchronize set

Vrspnb indicates the serial number in the case of resynchronize restart.

Vrsp and, if necessary Vrspnb, are set when an S-RESYNCHRONIZE request is issued or an S-RESYNCHRONIZE indication is received. Vrsp is set to no when the SS-user goes to STA713.

A.5.4.3 Vcoll

Vcoll is a boolean variable having the following values:

Vcoll = true : a collision of S-RELEASE requests has been detected;

Vcoll = false: there has not been a collision of S-RELEASE requests.

This variable is set false during the session connection establishment phase.

A.5.4.4 V(A)

V(A) is used by the SS-user and is the lowest serial number to which a synchronization point confirmation is expected. No confirmation is expected when $V(A) = V(M)$.

A.5.4.5 V(M)

V(M) is used by the SS-user and is the next serial number to be used.

A.5.4.6 V(R)

V(R) is used by the SS-user and is the lowest serial number to which resynchronization restart is permitted.

A.5.4.7 Vsc

Vsc is a boolean variable having the following values:

Vsc = true : the SS-user has the right to issue minor synchronization point responses when $V(A)$ is less than $V(M)$;

Vsc = false: the SS-user does not have the right to issue minor synchronization point responses.

Vsc is set false during the session connection establishment phase and when an S-SYNC-MINOR request is issued. Vsc is set true when an S-SYNC-MINOR indication is received.

NOTE — Table 33 summarizes the operations on V(A), V(M), V(R) and Vsc.

A.5.4.8 Vdnr

Vdnr is a boolean variable having the following values:

Vdnr = true : an S-RELEASE confirm has been received in STA09 (following a collision of S-RELEASE requests);

Vdnr = false: no S-RELEASE confirm has been received.

This variable is set to false during the connection establishment phase.