
**Information technology — Security
techniques — Cryptographic techniques
based on elliptic curves —**

**Part 5:
Elliptic curve generation**

*Technologies de l'information — Techniques de sécurité — Techniques
cryptographiques fondées sur les courbes elliptiques —*

Partie 5: Génération de courbes elliptiques

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Published in Switzerland

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Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.

International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 2.

The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

Attention is drawn to the possibility that some of the elements of this document may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights.

ISO/IEC 15946-5 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 27, *IT Security techniques*.

ISO/IEC 15946 consists of the following parts, under the general title *Information technology — Security techniques — Cryptographic techniques based on elliptic curves*:

- *Part 1: General*
- *Part 5: Elliptic curve generation*

Introduction

Some of the most interesting alternatives to the RSA and $F(p)$ based systems are cryptosystems based on elliptic curves defined over finite fields. The concept of an elliptic curve based public-key cryptosystem is rather simple:

- Every elliptic curve over a finite field is endowed with an addition operation “+”, under which it forms a finite abelian group.
- The group law on elliptic curves extends in a natural way to a “discrete exponentiation” on the point group of the elliptic curve.
- Based on the discrete exponentiation on an elliptic curve, one can easily derive elliptic curve analogues of the well-known public-key schemes of Diffie-Hellman and ElGamal type.

The security of such a public-key system depends on the difficulty of determining discrete logarithms in the group of points of an elliptic curve. This problem is — with current knowledge — much harder than the factorization of integers or the computation of discrete logarithms in a finite field. Indeed, since Miller and Koblitz in 1985 independently suggested the use of elliptic curves for public-key cryptographic systems, the elliptic curve discrete logarithm problem has only been shown to be solvable in certain specific, and easily recognizable, cases. There has been no substantial progress in finding an efficient method for solving the elliptic curve discrete logarithm problem on arbitrary elliptic curves. Thus, it is possible for elliptic curve based public-key systems to use much shorter parameters than the RSA system or the classical discrete logarithm based systems that make use of the multiplicative group of a finite field. This yields significantly shorter digital signatures and system parameters.

This part of ISO/IEC 15946 describes elliptic curve generation techniques useful for implementing the elliptic curve based mechanisms defined in ISO/IEC 9796-3, ISO/IEC 11770-3, ISO/IEC 14888-3, and ISO/IEC 18033-2.

It is the purpose of this part of ISO/IEC 15946 to meet the increasing interest in elliptic curve based public-key technology by describing elliptic curve generation methods to support key-exchange, key-transport and digital signatures based on an elliptic curve.

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Information technology — Security techniques — Cryptographic techniques based on elliptic curves —

Part 5: Elliptic curve generation

1 Scope

ISO/IEC 15946 specifies public-key cryptographic techniques based on elliptic curves.

This part of ISO/IEC 15946 defines elliptic curve generation techniques useful for implementing the elliptic curve based mechanisms defined in ISO/IEC 9796-3, ISO/IEC 11770-3, ISO/IEC 14888-3 and ISO/IEC 18033-2.

The scope of this part of ISO/IEC 15946 is restricted to cryptographic techniques based on elliptic curves defined over finite fields of prime power order (including the special cases of prime order and characteristic two). The representation of elements of the underlying finite field (i.e. which basis is used) is outside the scope of this part of ISO/IEC 15946.

ISO/IEC 15946 does not specify the implementation of the techniques it defines. Interoperability of products complying with ISO/IEC 15946 will not be guaranteed.

2 Normative reference(s)

The following referenced documents are indispensable for the application of this document. For dated references, only the edition cited applies. For undated references, the latest edition of the referenced document (including any amendments) applies.

ISO/IEC 15946-1, *Information technology — Security techniques — Cryptographic techniques based on elliptic curves — Part 1: General*

3 Terms and definitions

For the purposes of this document, the following terms and definitions apply.

3.1

definition field of an elliptic curve

field that includes all the coefficients of the equation describing an elliptic curve

3.2

elliptic curve

cubic curve without a singular point

NOTE 1 A definition of a cubic curve is given in [29].

NOTE 2 The set of points of E under a certain addition law forms an abelian group. In this part of ISO/IEC 15946, we only deal with finite fields F as the definition field. When we describe the definition field F of an elliptic curve E explicitly, we denote the curve as E/F .

NOTE 3 A detailed definition of an elliptic curve is given in Clause 4.

[ISO/IEC 15946-1:2008]

3.3

finite field

field containing a finite number of elements

NOTE 1 A definition of field is given in [29].

NOTE 2 For any positive integer m and a prime p , there exists a finite field containing exactly p^m elements. This field is unique up to isomorphism and is denoted by $F(p^m)$, where p is called the characteristic of $F(p^m)$.

[ISO/IEC 15946-1:2008]

3.4

hash-function

function which maps strings of bits to fixed-length strings of bits, satisfying the following two properties:

- for a given output, it is computationally infeasible to find an input which maps to this output;
- for a given input, it is computationally infeasible to find a second input which maps to the same output.

[ISO/IEC 10118-1]

NOTE 1 Computational feasibility depends on the specific security requirements and environment.

NOTE 2 For the purposes of this document, the recommended hash-functions are those defined in ISO/IEC 10118-2 and ISO/IEC 10118-3.

3.5

nearly prime number

positive integer $n = m \cdot r$, where m is a large prime number and r is a small smooth integer

NOTE The meaning of the terms large and small prime numbers is dependent on the application, and is based on bounds determined by the designer.

3.6

order of an elliptic curve $E(F)$

number of points on an elliptic curve E defined over a finite field F

3.7

smooth integer

integer r whose prime factors are all small (i.e. less than some defined bound)

4 Notation and conversion functions

4.1 Notation

In this part of ISO/IEC 15946, the following notation is used to describe public-key systems based on elliptic curve technology.

B An embedding degree, the smallest B such that $\#E(F(q)) \mid q^B - 1$.

E An elliptic curve, given by an equation of the form $Y^2 = X^3 + aX + b$ over the field $F(p^m)$ for $p > 3$, by an equation of the form $Y^2 + XY = X^3 + aX^2 + b$ over the field $F(2^m)$, or by an equation of the form $Y^2 = X^3 + aX^2 + b$ over the field $F(3^m)$, together with an extra point O_E referred to as the point at infinity. The elliptic curve is denoted by $E/F(p^m)$, $E/F(2^m)$, or $E/F(3^m)$, respectively.

NOTE 1 In applications not based on a pairing, $E/F(p)$ or $E/F(2^m)$ is preferable from an efficiency point of view. In applications that use a pairing, $E/F(p)$ or $E/F(3^m)$ is preferable from an efficiency point of view.

$\#E(F(q))$ The order (or cardinality) of $E(F(q))$.

n A prime divisor of $\#E(F(q))$.

N The number of points on an elliptic curve E over $F(q)$, $\#E(F(q))$.

r The cofactor, that is $\#E(F(q)) = rn$.

4.2 Conversion functions

For the purposes of this part of ISO/IEC 15946, the following conversion functions, defined in ISO/IEC 15946-1:2008, are used.

BS2IP The bit string to integer conversion primitive

BS2OSP The bit string to octet string conversion primitive

EC2OSP _{E} The elliptic curve point to octet string conversion primitive

FE2IP _{F} The finite field element to integer conversion primitive

FE2OSP _{F} The finite field element to octet string conversion primitive

I2BSP The integer to bit string conversion primitive

I2OSP The integer to octet string conversion primitive

I2ECP The integer to elliptic curve conversion primitive

OS2BSP The octet string to bit string conversion primitive

OS2FEP _{F} The octet string to finite field element conversion primitive

OS2ECP _{E} The octet string to elliptic curve point conversion primitive

OS2IP The octet string to integer conversion primitive

5 Framework for elliptic curve generation

5.1 Types of trusted elliptic curve

There are a number of ways in which a user can obtain trust in the provenance of an elliptic curve, including the following.

- The curve may be obtained from an impartial trusted source (e.g. an international or national standard).
- The curve may be generated and/or verified by a trusted third party.

- The curve may be generated and/or verified by the user.

5.2 Overview of elliptic curve generation

There are three main ways to generate elliptic curves.

- Generate an elliptic curve by applying the order counting algorithms to a (pseudo-)randomly chosen elliptic curve. Such a technique is specified in Clause 6.
- Generate an elliptic curve by applying the complex multiplication method. Such a technique is specified in Clause 7.
- Generate an elliptic curve by lifting an elliptic curve over a small finite field to that over a reasonably large field. Such a technique is specified in Clause 8.

6 Verifiably Pseudo-Random Elliptic curve generation

6.1 Constructing Verifiably Pseudo-Random Elliptic Curves (prime case)

6.1.1 Construction algorithm

The following algorithm produces a set of elliptic curve parameters over a field $F(p)$ selected (pseudo-)randomly from the curves of appropriate order, along with sufficient information for others to verify that the curve was indeed chosen pseudo-randomly.

NOTE 1 The algorithm is consistent with [16].

It is assumed that the following quantities have been chosen:

- lower bound n_{min} for the order of the base point.
- a cryptographic hash function H with output length L_{Hash} bits.
- the bit length L of inputs to H , satisfying $L \geq L_{Hash}$.

The following notation is adopted below:

- $v = \lceil \log_2 p \rceil$,
- $s = \lfloor (v - 1)/L_{Hash} \rfloor$,
- $w = v - sL_{Hash} - 1$.

Input: a prime number p ; lower bound n_{min} for n ; a trial division bound l_{max} .

Output: a bit string X ; EC parameters a , b , n , and G .

- a) Choose an arbitrary bit string X of bit length L .
- b) Compute $h = H(X)$.
- c) Let W_0 be the bit string obtained by taking the w rightmost bits of h .
- d) Convert X to an integer $Z = \text{BS2IP}(X)$.

- e) For i from 1 to s do:
- 1) Convert the integer $(Z + i) \bmod 2^L$ to a length- L bit string X_i using I2BSP.
 - 2) Compute $W_i = H(X_i)$.
- f) Let W be the bit string obtained by the concatenation of $W_0, W_1, \dots, W_s$ as follows:

$$W = W_0 \parallel W_1 \parallel \dots \parallel W_s.$$

- g) Convert W to a finite field element $c = \text{OS2FEP}(\text{BS2OSP}(W))$.

- h) If $c = 0_F$ or $4c + 27 = 0_F$, then go to Step a).

- i) Choose finite field elements $a, b \in F(p)$ such that $b \neq 0_F$ and $cb^2 - a^3 = 0_F$.

NOTE 2 The simplest choice is $a = c$ and $b = c$. However, an implementer may want to choose differently for performance reasons.

- j) Compute the order $\#E(F(p))$ of the elliptic curve E over $F(p)$ given by $y^2 = x^3 + ax + b$.
- k) Test whether $\#E(F(p))$ is a nearly prime number using the algorithm specified in 6.1.2. If so, the output of the algorithm specified in 6.1.2 consists of the integers r, n . If not, then go to Step a).
- l) Check $E(F(p))$ satisfies the MOV-condition specified in B.2.3, that is the smallest integer B such that n divides $q^B - 1$ ensures the desirable security level. If not, then go to Step a).
- m) Test whether $\#E(F(p)) \neq p$ in order to be secure against the attack specified in B.2.2. If not, then go to Step a).
- n) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then go to Step a).
- o) Generate a point G on E of order n using the algorithm specified in 6.1.3.
- p) Output X, a, b, n, G .

NOTE 3 The necessity of near primality is described in B.2.2.

NOTE 4 Methods to compute the order $\#E(F(p))$ are described in [5], [26] and [29].

6.1.2 Test for Near Primality

Given a lower bound $n_{\min}$ and a trial division bound $l_{\max}$, the following procedures test $N = \#E(F(p))$ for near primality.

Input: positive integers $N, l_{\max}$, and $n_{\min}$.

Output: if N is nearly prime, output a prime n with $n_{\min} \leq n$ and a smooth integer r such that $N = rn$. If N is not nearly prime, output the message "not nearly prime".

- a) Set $n = N, r = 1$.
- b) For l from 2 to $l_{\max}$ do
 - 1) If l is composite then go to Step 3).
 - 2) While (l divides n)

- Set $n = n/l$ and $r = rl$.
- If $n < n_{min}$ then output “not nearly prime” and stop.

3) Next l .

- c) Test n for primality.
- d) If n is prime then output r and n and stop.
- e) Output “not nearly prime”.

NOTE Methods to test for primality are described in [3] and [4].

6.1.3 Finding a Point of Large Prime Order

If the order $\#E(F(q))$ of an elliptic curve E is nearly prime, the following algorithm efficiently produces a random point in $E(F(q))$ whose order is the large prime factor n of $\#E(F(q)) = rn$.

Input: an elliptic curve E over the field $F(q)$, a prime n , and a positive integer r not divisible by n .

Output: if $\#E(F(q)) = rn$, a point G on E of order n ; if not, the message “wrong order.”

- a) Generate a random point P (not O_E) on E .
- b) Set $G = rP$.
- c) If $G = O_E$ then go to Step a).
- d) Set $Q = nG$.
- e) If $Q \neq O_E$ then output “wrong order” and stop.
- f) Output G .

6.1.4 Verification of Elliptic Curve Pseudo-Randomness

The following algorithm determines whether an elliptic curve over $F(p)$ was generated using the method of 6.1.1. The quantities L_{Hash} , L , γ , s , and w , and the hash function H , are as in 6.1.1.

Input: a bit string X of length L , EC parameters $q = p, a, b, n$, and $G = (x_G, y_G)$, and a positive integer n_{min} .

Output: “True” or “False”.

- a) Compute $h = H(X)$.
- b) Let W_0 be the bit string obtained by taking the w rightmost bits of h .
- c) Convert X to an integer $Z = \text{BS2IP}(X)$.
- d) For i from 1 to s do:
 - 1) Compute $Z = Z + i \bmod 2^L$.
 - 2) Convert $Z \bmod (2^L)$ to a bit string $X_i = \text{I2BSP}(Z)$.
 - 3) Compute $W_i = H(X_i)$.

e) Let W be the bit string obtained by the concatenation of $W_0, W_1, \dots, W_s$ as follows:

$$W = W_0 \parallel W_1 \parallel \dots \parallel W_s.$$

f) Convert W to a finite field element $c = \text{OS2FEP}(\text{BS2OSP}(W))$.

g) Verify the following conditions.

- $n \geq n_{min}$
- n is a prime.
- $c \neq 0_F$
- $4c + 27 \neq 0_F$.
- $b \neq 0_F$
- $cb^2 - a^3 \neq 0_F$.
- $G \neq O_E$
- $y_G^2 = x_G^3 + ax_G + b$.
- $nG = O_E$

h) If all the conditions in Step g) hold, then output “True”, otherwise output “False”.

6.2 Constructing Verifiably Pseudo-Random Elliptic Curves (binary case)

6.2.1 Construction algorithm

The following algorithm produces a set of elliptic curve parameters for a pseudo-random curve over a field $F(2^m)$, along with sufficient information for others to verify that the curve was indeed chosen pseudo-randomly.

NOTE 1 The algorithm is consistent with [16].

It is assumed that the following quantities have been chosen:

- a field $F(2^m)$
- a lower bound n_{min} for the order of the base point
- a cryptographic hash function H with output length L_{Hash} bits
- the bit length L of inputs to H , satisfying $L \geq L_{Hash}$.

The following notation is adopted below:

- $s = \lfloor (m - 1) / L_{Hash} \rfloor$,
- $w = m - sL_{Hash}$.

Input: a field $F(2^m)$; a lower bound n_{min} for n ; a trial division bound l_{max} .

Output: a bit string X ; EC parameters a, b, n , and G .

- a) Choose an arbitrary bit string X of bit length L .
- b) Compute $h = H(X)$.
- c) Let W_0 be the bit string obtained by taking the w rightmost bits of h .
- d) Convert the length- L bit string X to an integer z using BS2IP.
- e) For i from 1 to s do:
 - Convert the integer $(z + i) \bmod (2^L)$ to a length- L bit string X_i using I2BSP.
 - Compute $W_i = H(X_i)$.
- f) Let W be the bit string obtained by the concatenation of $W_0, W_1, \dots, W_s$ as follows:

$$W = W_0 \parallel W_1 \parallel \dots \parallel W_s.$$
- g) Convert the length- m bit string W to a field element b using BS2OSP and OS2FEP.
- h) If $b = 0_F$, then go to Step a).
- i) Let a be an arbitrary element in $F(2^m)$. (The simplest choice is $a = 0_F$. However, one may want to choose differently for performance reasons.)
- j) Compute the order $\#E(F(2^m))$ of the elliptic curve E over $F(2^m)$ given by $y^2 + xy = x^3 + ax^2 + b$.
- k) Test whether $\#E(F(2^m))$ is a nearly prime number using the algorithm specified in 6.1.2. If so, the output of the algorithm specified in 6.1.2 consists of the integers r, n . If not, then go to Step a).
- l) Check $E(F(2^m))$ satisfies the MOV-condition specified in B.2.3. If not, then go to Step a).
- m) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then go to Step a).
- n) Generate a point G on E of order n using the algorithm specified in 6.1.3.
- o) Output X, a, b, n, G .

NOTE 2 The necessity of near primality is described in B.2.2.

NOTE 3 Methods of computing the order $\#E(F(2^m))$ are described in [5], [26] and [29].

6.2.2 Verification of Elliptic Curve Pseudo-Randomness

The following algorithm verifies the validity of a set of elliptic curve parameters. In addition, it determines whether an elliptic curve over $F(2^m)$ was generated using the method of 6.2.1.

The quantities L_{Hash}, L, s , and w , and the hash function H , are as in 6.2.1.

Input: a bit string X of length L , EC parameters $q = 2^m, a, b, n$, and $G = (x_G, y_G)$, and a positive integer $n_{\min}$.

Output: "True" or "False".

- a) Compute $h = H(X)$.
- b) Let W_0 be the bit string obtained by taking the w rightmost bits of h .
- c) Convert the bit string X to an integer z via BS2IP.

- d) For i from 1 to s do:
- 1) Convert the integer $(z + i) \bmod (2^L)$ to a length- L bit string X_i via I2BSP.
 - 2) Compute $W_i = H(X_i)$.
- e) Let W be the bit string obtained by the concatenation of $W_0, W_1, \dots, W_s$ as follows:

$$W = W_0 \parallel W_1 \parallel \dots \parallel W_s$$

- f) Convert the length- m bit string W to the field element b' via BS2OSP and OS2FEP.
- g) Verify the following conditions.

- $n \geq n_{min}$
- n is a prime.
- $b \neq 0_F$
- $b = b'$
- $G \neq O_E$
- $y_G^2 + x_G y_G = x_G^3 + ax_G^2 + b$
- $nG = O_E$

- h) If all the conditions in Step g) hold, then output "True"; otherwise output "False".

7 Constructing Elliptic Curves by Complex Multiplication

7.1 General Construction (prime case)

The following algorithm produces an elliptic curve E over $F(p)$ with the given number of rational points N .

NOTE 1 The algorithm is based on [11], which is applied to primality proving [4].

Input: The definition field $F(p)$ and the number of points $N = rn$, where n is the largest prime divisor of N and r is a cofactor.

Output: curve parameters of elliptic curve E with $\#E(F(p)) = N$ and base point G

- a) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then execute a new input.
- b) Set $t = p + 1 - N$.
- c) Choose a pair of integers (D, V) such that $4p - t^2 = DV^2$.
- d) Construct the Hilbert class polynomial $P_D(X)$.
- e) Find a solution j_0 in $F(p)$ of $P_D(X) = 0$ modulo p .
- f) Choose $c \in F(p)^*$ and construct an elliptic curve over $F(p)$ with the j -invariant j_0 .
 - $E_{D,j_0,c}: y^2 = x^3 + (3c^2j_0 / (1728 - j_0))x + 2c^3j_0 / (1728 - j_0)$ (if $j_0 \neq 0_F, 1728$).

- $E_{D,j_0,c}: y^2 = x^3 + c$ (if $j_0 = 0_F$).
- $E_{D,j_0,c}: y^2 = x^3 + cx$ (if $j_0 = 1728$).

- g) Construct a random point G on $E_{D,j_0,c}(F(p))$ such that $G \neq O_E$ and $r \cdot G \neq O_E$.
- h) Set $G = r \cdot G$.
- i) If $n \cdot G = O_E$, output curve parameters of $E_{D,j_0,c}$ and the base point G . If $n \cdot G \neq O_E$, go to Step f) to choose another c .

NOTE 2 Any pair of integers (D, V) such that $4p - t^2 = DV^2$ can be used in Step c).

NOTE 3 The definition of the Diophantine equation used in Step c) is given in A.5.

NOTE 4 The definition of the Hilbert class polynomial $P_D(X)$ is given in A.2.

7.2 MNT curve (Miyaji-Nakabayashi-Takano curve)

The following algorithm produces an elliptic curve E over $F(p)$ with the embedding degree $B = 6$, which is useful for cryptosystems based on a bilinear pairing. The pairing and the embedding degree are described in A.3 and B.2.2, respectively.

NOTE 1 Detailed information is given in [20].

Input: lower and upper bound (odd integer) p_{min} and p_{max} for the definition field (in bits) and upper bound D_{max} for size of D .

Output: prime p , curve parameters of elliptic curve $E/F(p)$, the order $n = \#E(F(p))$, and basepoint G .

- a) Choose a small positive integer $D < D_{max}$ such that $D \equiv 3 \pmod{8}$ and go to Step c).
- b) If such D does not exist, then stop and output "fail".
- c) Find a pair of integers (T, U) with the smallest $U > 0$ that satisfies $T^2 - 3DU^2 = \pm 1$ using the continued fraction algorithm.
- d) Find a pair of integers (x, y) that satisfies $x^2 - 3Dy^2 = -8$ and

$$0 \leq x < 2U\sqrt{2D}, 2\sqrt{2/D} \leq y < 2T\sqrt{2/D},$$
 using the algorithm of Lagrange. If not, go to Step a).
- e) $i = 0$.
- f) Find a pair of primes (p, n) as follows:
 - 1) Compute integers x_i and y_i such that $x_i + y_i\sqrt{3D} = (x + y\sqrt{3D})(T + U\sqrt{3D})^i$.
 - 2) If $x_i \equiv 1 \pmod{6}$, then $s = (x_i - 1)/6$ and $p = 4s^2 - 1$;
 - else if $x_i \equiv -1 \pmod{6}$, then $s = (x_i + 1)/6$ and $p = 4s^2 - 1$;
 - else, $i = i + 1$ and go to Step 1).
 - 3) If $p < p_{min}$, then $i = i + 1$ and go to Step 1).
 - 4) If $p > p_{max}$, then go to Step a).

- 5) If p is prime, then $n_1 = 4s^2 + 2s + 1$ and $n_2 = 4s^2 - 2s + 1$;
 — else, $i = i + 1$ and go to Step 1).
- 6) If n_1 or n_2 prime, then $n = n_1$ or $n = n_2$, respectively and go to Step g);
 — else, $i = i + 1$ and go to Step 1).
- g) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then go to Step a).
- h) Construct the Hilbert class polynomial $P_D(X)$.
- i) Find a solution j_0 in $F(p)$ of $P_D(X) = 0$ modulo p .
- j) Choose $c \in F(p)^*$ and construct an elliptic curve over $F(p)$ with the j -invariant j_0 .
 — $E_{D,j_0,c}: y^2 = x^3 + (3c^2j_0 / (1728 - j_0))x + 2c^3j_0 / (1728 - j_0)$ (if $j_0 \neq 0_F, 1728$).
 — $E_{D,j_0,c}: y^2 = x^3 + c$ (if $j_0 = 0_F$).
 — $E_{D,j_0,c}: y^2 = x^3 + cx$ (if $j_0 = 1728$).
- k) Construct a random point G on $E_{D,j_0,c}(F(p))$, not equal to the point at infinity O_E .
- l) If $n \cdot G = O_E$, output p, E, n , and G .
- m) Else, go to Step h) to choose another $c \in F(p)^*$.

NOTE 2 The definition of the Hilbert class polynomial $P_D(X)$ is given in A.2.

NOTE 3 The continued fraction algorithm in Step c) is given in A.3 and [23].

NOTE 4 The algorithm of Lagrange in Step d) is given in A.4, [18] and [21].

NOTE 5 A technique for speeding up a protocol based on a bilinear pairing is described in [6].

7.3 BN curve (Barreto-Naehrig curve)

The following algorithm produces an elliptic curve E over $F(p)$ with the embedding degree $B = 12$, which is useful for cryptosystems based on bilinear pairings. The embedding degree is described in B.2.2.

NOTE 1 Detailed information is given in [6].

Input: the approximate desired size m of the curve order (in bits) and upper bound (odd integer) p_{max} for the definition field

Output: prime p , curve parameters of elliptic curve $E/F(p)$, the order $n = \#E(F(p))$, and basepoint G .

- a) Let $P(u) = 36u^4 + 36u^3 + 24u^2 + 6u + 1$.
- b) Compute the smallest $u \approx 2^{m/4}$ such that $\lceil \log_2 P(-u) \rceil = m$.
- c) While $p \leq p_{max}$
- 1) $t = 6u^2 + 1$.
 - 2) $p = P(-u)$ and $n = p + 1 - t$.

- 3) If p and n are prime then go to Step e).
- 4) $p = P(u)$ and $n = p + 1 - t$.
- 5) If p and n are prime, then go to Step e).
- 6) $u = u + 1$ and go to Step 1).
- d) Stop and output "fail".
- e) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then go to Step a).
- f) $b = 0$.
- g) If $b + 1$ is not represented by $b + 1 = y_0^2$ modulo p for an integer y_0 , then $b = b + 1$ and go to Step g).
- h) Set an elliptic curve $E: y^2 = x^3 + b$.
- i) Compute a square root $y_0 = \sqrt{(b + 1)}$ modulo p .
- j) Set the basepoint $G = (1, y_0) \in E$.
- k) If $n \cdot G \neq O_E$, then set $b = b + 1$ and go to Step g).
- l) Output p , E , n , and G .

NOTE 2 A technique for speeding up a protocol based on a bilinear pairing is described in [6].

7.4 F curve (Freeman curve)

The following algorithm produces an elliptic curve E over $F(p)$ with embedding degree $B = 10$, which is useful for cryptosystems based on a bilinear pairing. The embedding degree is described in B.2.2.

NOTE 1 Detailed information is given in [13].

Input: lower and upper bound p_{min} and p_{max} for the size of the definition field (in bits) and upper bound D_{max} for size of D .

Output: prime p , curve parameters of elliptic curve $E/F(p)$, the order $n = \#E(F(p))$, and basepoint G .

- a) Choose a small positive integer $D < D_{max}$ such that $D \equiv 43$ or $67 \pmod{120}$ and $15D$ is square-free and go to Step c).
- b) If such D does not exist, then stop and output "fail".
- c) Find a pair of integers (T, U) with the smallest $U > 0$ that satisfies $T^2 - 15DU^2 = \pm 1$ using the continued fraction algorithm.
- d) Let $d = T - U\sqrt{15D}$.
- e) Let $g = d^2$ if $T^2 - 15DU^2 = -1$, and let $g = d$ otherwise.
- f) Find a pair of integers (x, y) that satisfies $x^2 - 15Dy^2 = -20$ and $0 \leq x < 10U\sqrt{3D}$, $2\sqrt{1/(3D)} \leq y < 2T\sqrt{1/(3D)}$ using the algorithm of Lagrange.
- g) For the solution (x, y) in Step f).

- 1) If $x = \pm 5 \pmod{15}$, then:
 - Let $s = (-5 \pm x)/15$.
 - Let $p = 25s^4 + 25s^3 + 25s^2 + 10s + 3$.
 - Let $n = 25s^4 + 25s^3 + 15s^2 + 5s + 1$.
- 2) Else, go to Step f) to the next (x, y) .
- 3) If $p > p_{max}$, go to Step f) to use the next (x, y) .
- 4) Else if $p < p_{min}$, then go to Step 6).
- 5) If p and n are primes, go to Step h).
- 6) Find a pair of integers (x', y') such that $x' + y'\sqrt{15D} = (x + y\sqrt{15D}) \cdot g$.
- 7) Let $x = x'$ and $y = y'$ and return to Step 1).
- h) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then go to Step a).
- i) Construct the Hilbert class polynomial $P_D(X)$.
- j) Find a solution j_0 in $F(p)$ of $P_D(X) = 0$ modulo p .
- k) Choose $c \in F(p)^*$ and construct an elliptic curve E over $F(p)$ with j -invariant j_0 :
 - $E_{D,j_0,c} : y^2 = x^3 + (3c^2j_0 / (1728 - j_0))x + 2c^3j_0 / (1728 - j_0)$ (if $j_0 \neq 0_F, 1728$).
 - $E_{D,j_0,c} : y^2 = x^3 + c$ (if $j_0 = 0_F$).
 - $E_{D,j_0,c} : y^2 = x^3 + cx$ (if $j_0 = 1728$).
- l) Construct a random point G on $E_{D,j_0,c}(F(p))$, not equal to the point at infinity O_E .
- m) If $n \cdot G = O_E$, output p, E, n and G .
- n) Else, go to Step k) to choose another $c \in F(p)^*$.

NOTE 2 The definition of the Hilbert class polynomial $P_D(X)$ in Step i) is given in A.2.

NOTE 3 The continued fraction algorithm in Step c) is given in A.3 and [23].

NOTE 4 The algorithm of Lagrange in Step f) is given in A.4, [18] and [21].

NOTE 5 A technique to speed up a protocol based on a bilinear pairing is described in [6].

7.5 CP curve (Cocks-Pinch curve)

The following algorithm produces an elliptic curve E over $F(p)$ with arbitrary embedding degree B , which is useful for cryptosystems based on a bilinear pairing. The embedding degree is described in B.2.2.

NOTE 1 Detailed information is given in [7].

Input: a positive integer B and a set R of prime numbers n ($n-1$ is divisible by B).

Output: prime p , curve parameters of elliptic curve $E/F(p)$, the order $n \cdot r = \#E(F(p))$, and basepoint G .

- a) Choose a small square-free positive integer D and n in R such that $-D$ is a square modulo n .
 - 1) Find a B -th root of unity z in $F(n) \setminus \{0, 1\}$.
 - 2) $t' = z + 1$
 - 3) $y' = (t'-2) / \sqrt[3]{(-D)} \pmod{n}$
 - 4) Let t be an integer such that t is equal to t' modulo n , and let y be an integer such that y is equal to y' modulo n .

NOTE 2 $t = t'$ and $y = y'$ are permitted.

- 5) $p = (t^2 + Dy^2) / 4$
- 6) If p is not prime, then go to Step a).
- b) Test whether the prime divisor n satisfies the condition described in B.2.4 for cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5. If not, then go to Step a).
- c) Construct the Hilbert class polynomial $P_D(X)$.
- d) Find a solution j_0 in $F(p)$ of $P_D(X) = 0$ modulo p .
- e) Choose $c \in F(p)^*$ and construct an elliptic curve over $F(p)$ with the j -invariant j_0 .
 - $E_{D,j_0,c}: y^2 = x^3 + (3c^2j_0 / (1728 - j_0))x + 2c^3j_0 / (1728 - j_0)$ (if $j_0 \neq 0_F, 1728$).
 - $E_{D,j_0,c}: y^2 = x^3 + c$ (if $j_0 = 0_F$).
 - $E_{D,j_0,c}: y^2 = x^3 + cx$ (if $j_0 = 1728$).
- f) Set a cofactor $r = (p + 1 - t) / n$.
- g) Construct a random point G on $E_{D,j_0,c}(F(p))$ such that $G \neq O_E$ and $r \cdot G \neq O_E$.
- h) Set $G = r \cdot G$.
- i) If $n \cdot G = O_E$, output n , G , and the elliptic curve E .
- j) Else, go to Step e) to choose another $c \in F(p)^*$.

NOTE 3 The definition of the Hilbert class polynomial $P_D(X)$ in Step c) is given in A.2.

NOTE 4 A technique for speeding up a protocol based on a bilinear pairing is described in [6].

8 Constructing Elliptic Curves by Lifting

The following algorithm produces an elliptic curve E over $F(p^m)$ by lifting an elliptic curve E over $F(p)$.

NOTE The algorithm is based on [29].

Input: small finite field $F(p)$, elliptic curve E over $F(p)$, lower and upper bound N_{min} and N_{max} for the order of elliptic curve (in bits).

Output: extension degree m , order $N_m = \#E(F(p^m))$, basepoint G , and order n of G .

- a) Count the order of $N = \#E(F(p))$, which is easily executed since $F(p)$ is small.
- b) Set $t = p + 1 - N$ and compute algebraic integers α and β that satisfy $t = \alpha + \beta$ and $p = \alpha\beta$.
- c) Set $m = 1$.
- d) Find a triple of (m, N_m, n) as follows:
 - 1) Compute $N_m = p^m + 1 - (\alpha^m + \beta^m)$ and $q = p^m$, which is an integer.
 - 2) If $N_m < N_{min}$, then $m = m + 1$ and go to Step 1).
 - 3) If $N_m > N_{max}$, then stop and output "fail".
 - 4) Test whether N_m is a nearly prime number using the algorithm specified in 6.1.2. If so, the output of 6.1.2 consists of the integers r and n . If not, then $m = m + 1$ and go to Step 1).
 - 5) Check whether $E(F(q))$ satisfies the MOV-condition specified in B.2.3, that is the smallest integer B such that n divides $q^B - 1$ ensures the desirable security level. If not, then $m = m + 1$ and go to Step 1).
- e) Generate a point G on $E(F(q))$ of order n using the algorithm specified in 6.1.3.
- f) Output an extension degree m , the order $N_m = \#E(F(q))$, a basepoint G and the order n .

Annex A (informative)

Background information on elliptic curves

A.1 j -invariant

Let $F(q)$ be a finite field with $q = p^m$, where prime $p > 3$. Let E be an elliptic curve over $F(q)$ given by the short Weierstrass equation,

$$Y^2 = X^3 + aX + b \text{ with } a, b \in F(q),$$

where the inequality $4a^3 + 27b^2 \neq 0_F$ holds in $F(q)$. Then the j -invariant is defined as

$$j = 1728 \cdot (4a^3) / (4a^3 + 27b^2).$$

Let $F(2^m)$, for some $m \geq 1$, be a finite field. Let E be an elliptic curve over $F(2^m)$ given by the equation

$$Y^2 + XY = X^3 + aX + b \text{ with } a, b \in F(2^m)$$

where $b \neq 0_F$. Then the j -invariant is defined as

$$j = 1/b.$$

Let $F(3^m)$, for some $m \geq 1$, be a finite field. Let E be an elliptic curve over $F(3^m)$ given by the equation

$$Y^2 = X^3 + aX^2 + b \text{ with } a, b \in F(3^m)$$

such that $a, b \neq 0_F$. Then the j -invariant is defined as

$$j = -a^3/b.$$

A.2 Hilbert class polynomial

The construction of elliptic curves by complex multiplication uses the theory of imaginary quadratic fields $Q(\sqrt{-D})$. In the case of the imaginary quadratic field $Q(\sqrt{-D})$, the Hilbert class field K is the extension field of $Q(\sqrt{-D})$, which is the unramified abelian extension of $Q(\sqrt{-D})$. The Hilbert class polynomial $P_D(X)$ is defined by the minimum polynomial of K over $Q(\sqrt{-D})$. In the construction of elliptic curves by complex multiplication, the fact that the j -invariants of elliptic curves $E/F(p)$ are given as a solution of a Hilbert class polynomial $P_D(X)$ modulo p is used.

NOTE 1 These facts are described in [7] and [10].

NOTE 2 Online databases of Hilbert class polynomials are available in [17].

A.3 Cryptographic pairing

A cryptographic pairing e_n satisfies the conditions of non-degeneracy, bilinearity, and computability. A pairing e_n is defined over $\langle G_1 \rangle \times \langle G_2 \rangle$ as follows,

$$e_n : \langle G_1 \rangle \times \langle G_2 \rangle \rightarrow \mu_n,$$

where $\langle G_1 \rangle$ and $\langle G_2 \rangle$ are the cyclic groups of order n and μ_n is the cyclic group of the n -th roots of unity. A pairing e_n is realized by restricting the domain of the Weil or Tate pairings.

A.4 Pell Equation

The Pell equation is of the form

$$T^2 - dU^2 = \pm 1,$$

where d is a fixed integer. In the construction of elliptic curves by complex multiplication, the Pell equation with a positive integer d that is not a perfect square is used. Then all positive integer solutions of (T, U) are given by using the least positive solution (T_0, U_0) with the smallest $U_0 > 0$ as follows:

$$T + U\sqrt{d} = (T_0 + U_0\sqrt{d})^k$$

for $k = 1, 2, \dots$

NOTE These facts are described in [24].

A.5 The Diophantine equation $x^2 - dy^2 = n$

In the construction of elliptic curves by complex multiplication, the Diophantine equation $x^2 - dy^2 = n$ is used. Here n is an integer and d is a positive integer that is not a perfect square. The number of integer solutions of this equation is zero or infinite. All positive solutions (x, y) are given by using the least positive solution (T_0, U_0) with the smallest $U_0 > 0$ of the related Pell equation of $T^2 - dU^2 = 1$.

NOTE Details are described in [24].

Annex B (informative)

Background information on elliptic curve cryptosystems

B.1 Definition of cryptographic problems

B.1.1 The elliptic curve discrete logarithm problem (ECDLP)

For an elliptic curve $E/F(q)$, the base point $G \in E(F(q))$ with order n , and a point $P \in E(F(q))$, the elliptic curve discrete logarithm problem (with respect to the base point G) is to find the integer $x \in [0, n-1]$ such that $P = xG$ if such an x exists.

The security of elliptic curve cryptosystems is based on the believed hardness of the elliptic curve discrete logarithm problem.

B.1.2 The computational elliptic curve Diffie-Hellman problem (ECDHP)

For an elliptic curve $E/F(q)$, the base point $G \in E(F(q))$ with order n , and points $aG, bG \in E(F(q))$, the computational elliptic curve Diffie-Hellman problem is to compute abG .

The security of some elliptic curve cryptosystems is based on the believed hardness of the computational elliptic curve Diffie-Hellman problem.

B.1.3 The decisional elliptic curve Diffie-Hellman problem (ECDDHP)

For an elliptic curve $E/F(q)$, the base point $G \in E(F(q))$ with order n , and points $aG, bG, Y \in E(F(q))$, the decisional elliptic curve Diffie-Hellman problem is to decide whether $Y = abG$ or not.

The security of some elliptic curve cryptosystems is based on the believed hardness of the decisional elliptic curve Diffie-Hellman problem.

B.1.4 The bilinear Diffie-Hellman problem (BDHP)

The bilinear Diffie-Hellman problems are described in two ways according to the corresponding cryptographic bilinear maps.

- For two groups $\langle G_1 \rangle$ and $\langle G_2 \rangle$ with order n , a cryptographic bilinear map $e_n : \langle G_1 \rangle \times \langle G_2 \rangle \rightarrow \mu_n$, $aG_1, bG_1 \in \langle G_1 \rangle$, and $aG_2, cG_2 \in \langle G_2 \rangle$, the bilinear Diffie-Hellman problem is to compute $e_n(G_1, G_2)^{abc}$.
- For a group $\langle G_1 \rangle$ with order n , a cryptographic bilinear map $e_n : \langle G_1 \rangle \times \langle G_1 \rangle \rightarrow \mu_n$, and $aG_1, bG_1, cG_1 \in \langle G_1 \rangle$, the bilinear Diffie-Hellman problem is to compute $e_n(G_1, G_1)^{abc}$.

The security of some elliptic curve cryptosystems is based on the believed hardness of the elliptic curve bilinear Diffie-Hellman problem.

B.1.5 The elliptic curve discrete logarithm problem with auxiliary inputs (ECDLP with auxiliary inputs)

The security of some cryptosystems is based on the elliptic curve discrete logarithm problem with auxiliary inputs. Some of them are as follows (the notation follows from the original definitions of the problems in B.1.1, B.1.2, and B.1.4).

- ECDLP with additional inputs $x^2G, x^3G, \dots, x^kG$
- ECDHP with additional inputs $a^2G, a^3G, \dots, a^kG$
- BDHP with additional inputs $a^2G_1, a^3G_1, \dots, a^kG_1$

B.2 Algorithms to determine discrete logarithms on elliptic curves

B.2.1 Security of ECDLP

The security of ECDLP depends on the selection of elliptic curves $E/F(q)$ and the size n of the order of the base point G . The size of n should be 160 bits or more to achieve the desired level of security.

This section gives the overviews of algorithms to solve ECDLP. The elliptic curve $E/F(q)$ shall be chosen to meet the defined security objectives against the following algorithms to solve ECDLP. The size of n shall be set to meet the defined security objectives against the baby-step-giant-step algorithm and various variants of the Pollard ρ algorithm.

B.2.2 Overview of algorithms

The following techniques are available to determine discrete logarithms on an elliptic curve:

- The Pohlig-Silver-Hellman algorithm. This is a “divide-and-conquer” method which reduces the discrete logarithm problem for an elliptic curve E defined over $F(q)$ to the discrete logarithm in the cyclic subgroups of prime order dividing $\#E(F(q))$.
- The baby-step-giant-step algorithm and various variants of the Pollard- ρ algorithm.

NOTE 1 The various variants of the Pollard- ρ algorithm are described in [29].

- The algorithm of Frey-Rück and the Menezes-Okamoto-Vanstone algorithm which both transform the discrete logarithm problem in a cyclic subgroup of E with prime order n to the smallest extension field $F(q^B)$ of $F(q)$ such that n divides $(q^B - 1)$, where B is called the embedding degree. The Frey-Rück algorithm runs under weaker conditions than the algorithm published by Menezes-Okamoto-Vanstone.
- The algorithm of Araki-Satoh, Smart and Semaev which solves the discrete logarithm problem for an elliptic curve E defined over $F(p^m)$ in the case $\#E(F(p^m)) = p^m$.

Unlike the situation of the discrete logarithm in the multiplicative group of some finite field there is no known “index-calculus” available in the case of elliptic curves. As for attacks using covering for special type of covers, e.g. the Weil descent attack, the GHS attack, etc. See Chapter 22 of [5].

NOTE 2 The Pohlig-Silver-Hellman and baby-step-giant-step algorithms work generally on all kinds of elliptic curves while the Frey-Rück, the Menezes-Okamoto-Vanstone, Araki-Satoh, Smart, and Semaev algorithms work only on curves with special properties.

B.2.3 The MOV-condition

Let n be as defined in the set of elliptic curve domain parameters, where n is a prime divisor of $\#E(F(q))$ and q is a power of a prime p . A value B is given as the smallest integer such that n divides $p^B - 1$. As mentioned above, the Frey-Rück and Menezes-Okamoto-Vanstone algorithms reduce the discrete logarithm problem in an elliptic curve over $F(q)$ to the discrete logarithm in the finite field $F(p^B)$ for some $B \geq 1$. By using the attack, the difficulty of the discrete logarithm problem in an elliptic curve $E/F(q)$ is related to the discrete logarithm problem in a finite field $F(p^B)$. The *subfield-adjusted MOV-condition* describes the degree B that ensures the security level of the discrete logarithm problem in an elliptic curve by the discrete logarithm problem in finite field. For some applications based on the Weil and Tate pairing, a reasonably small value of B such as 6 or more is preferable.

NOTE Information on the degree B is described in [15].

B.2.4 The condition of prime divisor n

For some cryptosystems based on ECDLP, ECDHP, or BDHP with auxiliary inputs as in B.1.5, the prime divisor n should satisfy the following conditions: there is no divisor d of $n - 1$ such that $(\log n)^2 < d < n^{1/2}$ and there is no divisor e of $n + 1$ such that $(\log n)^2 < e < n^{1/2}$. The divisors d and e are possibly composite.

NOTE Detailed information on d and e is given in [8].

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Annex C (informative)

Numerical examples

C.1 Numerical examples of Verifiably Pseudo-Random Elliptic Curves

C.1.1 Introduction

Refer to [12] for this Clause. The parameters are chosen from a seed using SHA-1.

C.1.2 Elliptic curve over a prime field (192 bits)

p ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff
 $2^{192}-2^{64}-1$

Equation of E $y^2 = x^3 + ax + b$

a ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff

b 64210519 e59c80e7 0fa7e9ab 72243049 feb8deec c146b9b1

(seed) X 3045ae6f c8422f64 ed579528 d38120ea e12196d5

(compressed) G 03 188da80e b03090f6 7cbf20eb 43a18800 f4ff0afd 82ff1012

(uncompressed) G 04 188da80e b03090f6 7cbf20eb 43a18800 f4ff0afd 82ff1012
 07192b95 ffc8da78 631011ed 6b24cdd5 73f977a1 1e794811

n ffffffff ffffffff ffffffff 99def836 146bc9b1 b4d22831

(cofactor) r 1

C.1.3 Elliptic curve over a prime field (224 bits)

p ffffffff ffffffff ffffffff 00000000 00000000 00000001
 $2^{224}-2^{96}+1$

a ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff

b b4050a85
 0c04b3ab f5413256 5044b0b7 d7bfd8ba 270b3943 2355ffb4

(seed) X db713447 99d5c7fc dc45cb59f a3b9ab8f 6a948bc5

(compressed) G 02 b70e0cbd
 6bb4bf7f 321390b9 4a03c1d3 56c21122 343280d6 115c1d21

(uncompressed) G 04 b70e0cbd 6bb4bf7f
 321390b9 4a03c1d3 56c21122 343280d6 115c1d21 bd376388
 b5f723fb 4c22dfe6 cd4375a0 5a074764 44d58199 85007e34

N ffffffff
 ffffffff ffffffff ffff16a2 e0b8f03e 13dd2945 5c5c2a3d

(cofactor) r 1

C.1.4 Elliptic curve over a prime field (256 bits)

P ffffffff 00000001
 00000000 00000000 00000000 ffffffff ffffffff ffffffff
 $2^{224}(2^{32}-1)+2^{192}+2^{96}-1$

A ffffffff 00000001
 00000000 00000000 00000000 ffffffff ffffffff ffffffff

B 5ac635d8 aa3a93e7
 b3ebbd55 769886bc 651d06b0 cc53b0f6 3bce3c3e 27d2604b

(seed) X c49d3608 86e70493 6a6678e1 139d26b7 819f7e90

(compressed) G 03 6b17d1f2 e12c4247
 f8bce6e5 63a440f2 77037d81 2deb33a0 f4a13945 d898c296

(uncompressed) G 04 6b17d1f2 e12c4247 f8bce6e5 63a440f2
 77037d81 2deb33a0 f4a13945 d898c296 4fe342e2 fe1a7f9b
 8ee7eb4a 7c0f9e16 2bce3357 6b315ece cbb64068 37bf51f5

N ffffffff 00000000
 ffffffff ffffffff bce6faad a7179e84 f3b9cac2 fc632551

(cofactor) r 1

C.1.5 Elliptic curve over a prime field (384 bits)

P ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff
 ffffffff ffffffff ffffffff 00000000 00000000 ffffffff
 $2^{384}-2^{128}-2^{96}+2^{32}-1$

A ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff
 ffffffff ffffffff ffffffff 00000000 00000000 ffffffff

B b3312fa7 e23ee7e4 988e056b e3f82d19 181d9c6e fe814112
 0314088f 5013875a c656398d 8a2ed19d 2a85c8ed d3ec2aef

(seed) X a335926a a319a27a 1d00896a 6773a482 7acdac73

(compressed) G	03 aa87ca22 be8b0537 8eb1c71e f320ad74 6e1d3b62 8ba79b98 59f741e0 82542a38 5502f25d bf55296c 3a545e38 72760ab7
(uncompressed) G	04 aa87ca22 be8b0537 8eb1c71e f320ad74 6e1d3b62 8ba79b98 59f741e0 82542a38 5502f25d bf55296c 3a545e38 72760ab7 3617de4a 96262c6f 5d9e98bf 9292dc29 f8f41dbd 289a147c e9da3113 b5f0b8c0 0a60b1ce 1d7e819d 7a431d7c 90ea0e5f
n	ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff c7634d81 f4372ddf 581a0db2 48b0a77a ecec196a ccc52973
(cofactor) r	1

C.1.6 Elliptic curve over a prime field (521 bits)

p	01ff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff 2 ⁵²¹ -1
a	01ff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff
b	0051 953eb961 8e1c9a1f 929a21a0 b68540ee a2da725b 99b315f3 b8b48991 8ef109e1 56193951 ec7e937b 1652c0bd 3bb1bf07 3573df88 3d2c34f1 ef451fd4 6b503f00
(seed) X	d09e8800 291cb853 96cc6717 393284aa a0da564ba
(compressed) G	0200c6 858e06b7 0404e9cd 9e3ecb66 2395b442 9c648139 053fb521 f828af60 6b4d3dba a14b5e77 efe75928 fe1dc127 a2ffa8de 3348b3c1 856a429b f97e7e31 c2e5bd66
(uncompressed) G	04 00c6858e 06b70404 e9cd9e3e cb662395 b4429c64 8139053f b521f828 af606b4d 3dbaa14b 5e77efe7 5928fe1d c127a2ff a8de3348 b3c1856a 429bf97e 7e31c2e5 bd660118 39296a78 9a3bc004 5c8a5fb4 2c7d1bd9 98f54449 579b4468 17afb1d7 273e662c 97ee7299 5ef42640 c550b901 3fad0761 353c7086 a272c240 88be9476 9fd16650
n	01ff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff ffffffff 51868783 bf2f966b 7fcc0148 f709a5d0 3bb5c9b8 899c47ae bb6fb71e 91386409
(cofactor) r	1

C.2 Numerical examples of MNT curve

C.2.1 Introduction

The examples of Miyaji-Nakabayashi-Takano (MNT) curve in this Clause are given in [22].

C.2.2 Elliptic curve over a prime field (160 bits)

p	8c72d321 e48aa141 9b22f914 cb43c112 b76d7ae5
a	8c72d321 e48aa141 9b22f914 cb43c112 b76d7ae2
b	299ce219 b7b01348 fc2b5007 b6ab1ee1 005676f7
(compressed) G	03 00000000 00000000 00000000 00000000 00000002
(uncompressed) G	04 00000000 00000000 00000000 00000000 00000002 0be8f0d3 623edada ce4c2fac a541679b 002f1d07
n	8c72d321 e48aa141 9b23b6b2 e4a85a07 3822640f
(cofactor) r	1

C.2.3 Elliptic curve over a prime field (256 bits)

p	b1d5054e 2f7b68aa ee7ef918 74dd140c 6919af9b 719ed905 f6529c2a 424a6332
a	b1d5054e 2f7b68aa ee7ef918 74dd140c 6919af9b 719ed902 f6529c2a 424a6332
b	ae3dd5d1 f97c497c 1d5452d1 b074a6c0 6a25d4e5 819ccd1c 6e974d68 ef44f266
(compressed) G	02 00000000 00000000 00000000 00000000 00000000 00000003
(uncompressed) G	04 00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000000 00000003 693d7af8 c4a29f8d e56e477f f569661c 4dcd2227 aac17b09 e4b4b0b7 03b978ce
n	b1d5054e 2f7b68ab e99c585a 8419ae9f b45c620e 5ef666c3 f6529c2a 424a6332
(cofactor) r	1

C.3 Numerical examples of BN curve**C.3.1 Summary**

All of the following examples are chosen so that p is the largest prime satisfying $p \equiv 3 \pmod{4}$ and $p \equiv 4 \pmod{9}$ for the largest parameter u with minimum Hamming weight, allowing the extension field $F(p^2)$ to be represented as $F(p)[i]/(i^2 + 1)$ and the extension field $F(p^{2m})$ to be represented as $F(p^2)[z]/(z^m - v)$ for $m = 2, 3, 6$ and $v = 1 + i$. Computation of square (or cube) roots needed for point and/or pairing compression is also simplified in both $F(p)$ and $F(p^2)$. Furthermore, the curve equation has the form $E: y^2 = x^3 + 3$ with the obvious basepoint $G = (1, 2)$, and the sextic twist $E'/F(p^2)$ of the form $E': y'^2 = x'^3 + 3v$ contains a subgroup of order n and cofactor $h = 2p - n$, with basepoint $G' = hG_0'$ where G_0' is a point with x -coordinate $x_0' = 1$. Finally, the isomorphism $\psi: E'/F(p^2) \rightarrow E/F(p^{12})$ takes the form $\psi(x', y') = (x' v^{-1} z^4, y' v^{-1} z^3)$, with $z^6 = v$. These properties