
**Information technology —
Security techniques — Encryption
algorithms —**

**Part 3:
Block ciphers**

AMENDMENT 1: SM4

*Technologies de l'information — Techniques de sécurité —
Algorithmes de chiffrement —*

Partie 3: Chiffrement par blocs

AMENDEMENT 1: SM4





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Information technology — Security techniques — Encryption algorithms —

Part 3: Block ciphers

AMENDMENT 1: SM4

Clause 1

In the first paragraph, replace "seven different block ciphers" with "eight different block ciphers".

Replace Table 1 with the following:

Block length	Algorithm name (see #)	Key length
64 bits	TDEA (4.2)	128 or 192 bits
	MISTY (4.3)	128 bits
	CAST-128 (4.4)	
	HIGHT (4.5)	
128 bits	AES (5.2)	128, 192 or 256 bits
	Camellia (5.3)	128 bits
	SEED (5.4)	
	SM4 (5.5)	

5.1

Replace the sentence with the following:

In this clause, four 128-bit block ciphers are specified: AES in 5.2, Camellia in 5.3, SEED in 5.4, and SM4 in 5.5.

5.5

Add new subclause 5.5 as follows:

5.5 SM4

5.5.1 The SM4 algorithm

The SM4 algorithm is a symmetric block cipher that can process data blocks of 128 bits, using a cipher key with length of 128 bits under 32 rounds.

5.5.2 SM4 encryption

A 128-bit block P is transformed into a 128-bit block C using the following procedure, where for $i = 0, 1, 2, 3$ the X_i are 32-bit variables, and for $i = 0, 1, \dots, 31$ the rk_i are 32-bit subkeys:

$$(1) P = X_0 \parallel X_1 \parallel X_2 \parallel X_3$$

(2) for $i = 0$ to 31:

$$X_{i+4} = F(X_i, X_{i+1}, X_{i+2}, X_{i+3}, rk_i)$$

$$(3) C = X_{35} \parallel X_{34} \parallel X_{33} \parallel X_{32}$$

5.5.3 SM4 decryption

The decryption operation is identical to the encryption operation, except that the rounds (and therefore the subkeys) are used in reverse order:

$$(1) C = X_{35} \parallel X_{34} \parallel X_{33} \parallel X_{32}$$

(2) for $i = 31$ to 0:

$$X_i = F(X_{i+4}, X_{i+1}, X_{i+2}, X_{i+3}, rk_i)$$

$$(3) P = X_0 \parallel X_1 \parallel X_2 \parallel X_3$$

5.5.4 SM4 functions

5.5.4.1 Function F

The function F is used for both encryption and decryption. The function F is defined as follows:

$$F(X_0, X_1, X_2, X_3, rk) = X_0 \oplus T(X_1 \oplus X_2 \oplus X_3 \oplus rk)$$

where X_i ($i = 0, 1, 2, 3$) and rk are bit strings of length 32, T is a permutation defined in 5.5.4.2.

5.5.4.2 Permutation T and T'

5.5.4.2.1 General

The permutation T is used both for encryption and decryption. T is a composition of a nonlinear transformation τ and a linear transformation L, that is $T(\cdot) = L(\tau(\cdot))$. The permutation T' is used for the key schedule. T' is a composition of the nonlinear transformation τ and a linear transformation L', that is $T'(\cdot) = L'(\tau(\cdot))$. T, T', L, L' and τ are all transformations on 32-bit strings.

5.5.4.2.2 Nonlinear transformation τ

The nonlinear transformation τ is defined as follows, where for $i = 0, 1, 2, 3$ the a_i are bytes and S is an S-box defined in 5.5.4.2.4:

$$\tau(a_0 \parallel a_1 \parallel a_2 \parallel a_3) = S(a_0) \parallel S(a_1) \parallel S(a_2) \parallel S(a_3).$$

5.5.4.2.3 Linear transformation L and L'

The linear transformation L is defined as follows (B is a 32-bit variable):

$$L(B) = B \oplus (B \lll 2) \oplus (B \lll 10) \oplus (B \lll 18) \oplus (B \lll 24).$$

The linear transformation L' is defined as follows (B is a 32-bit variable):

$$L'(B) = B \oplus (B \lll 13) \oplus (B \lll 23).$$

5.5.4.2.4 S-box S

The S-box S used in the transformation τ is presented in hexadecimal form in Table 17.

Table 17 — SM4 S-box

	0	1	2	3	4	5	6	7	8	9	a	b	c	d	e	f
0	d6	90	e9	fe	cc	e1	3d	b7	16	b6	14	c2	28	fb	2c	05
1	2b	67	9a	76	2a	be	04	c3	aa	44	13	26	49	86	06	99
2	9c	42	50	f4	91	ef	98	7a	33	54	0b	43	ed	cf	ac	62
3	e4	b3	1c	a9	c9	08	e8	95	80	df	94	fa	75	8f	3f	a6
4	47	07	a7	fc	f3	73	17	ba	83	59	3c	19	e6	85	4f	a8
5	68	6b	81	b2	71	64	da	8b	f8	eb	0f	4b	70	56	9d	35
6	1e	24	0e	5e	63	58	d1	a2	25	22	7c	3b	01	21	78	87
7	d4	00	46	57	9f	d3	27	52	4c	36	02	e7	a0	c4	c8	9e
8	ea	bf	8a	d2	40	c7	38	b5	a3	f7	f2	ce	f9	61	15	a1
9	e0	ae	5d	a4	9b	34	1a	55	ad	93	32	30	f5	8c	b1	e3
a	1d	f6	e2	2e	82	66	ca	60	c0	29	23	ab	0d	53	4e	6f
b	d5	db	37	45	de	fd	8e	2f	03	ff	6a	72	6d	6c	5b	51
c	8d	1b	af	92	bb	dd	bc	7f	11	d9	5c	41	1f	10	5a	d8
d	0a	c1	31	88	a5	cd	7b	bd	2d	74	d0	12	b8	e5	b4	b0
e	89	69	97	4a	0c	96	77	7e	65	b9	f1	09	c5	6e	c6	84
f	18	f0	7d	ec	3a	dc	4d	20	79	ee	5f	3e	d7	cb	39	48

5.5.5 SM4 key schedule

The key scheduling part accepts a 128-bit master key $MK = MK_0 \parallel MK_1 \parallel MK_2 \parallel MK_3$, and yields 32 subkeys, as shown below.

$$(1) K_0 \parallel K_1 \parallel K_2 \parallel K_3 = (MK_0 \oplus FK_0) \parallel (MK_1 \oplus FK_1) \parallel (MK_2 \oplus FK_2) \parallel (MK_3 \oplus FK_3)$$

(2) for $i = 0$ to 31:

$$rk_i = K_{i+4} = K_i \oplus T'(K_{i+1} \oplus K_{i+2} \oplus K_{i+3} \oplus CK_i)$$

The constants FK_i ($i = 0, 1, 2, 3$) are as follows (in hexadecimal form):

$$FK_0 = a3b1bac6, FK_1 = 56aa3350, FK_2 = 677d9197, FK_3 = b27022dc.$$

The constants CK_i ($i = 0, 1, \dots, 31$) are defined as follows. Suppose $CK_i = ck_{i,0} \parallel ck_{i,1} \parallel ck_{i,2} \parallel ck_{i,3}$, where $ck_{i,j}$ are bytes, and $ck_{i,j} = (4i+j) \times 7 \pmod{256}$ ($i = 0, 1, \dots, 31, j = 0, 1, 2, 3$).

Thus, the values of CK_i ($i = 0, 1, \dots, 31$) are (in hexadecimal form):

00070e15, 1c232a31, 383f464d, 545b6269,
 70777e85, 8c939aa1, a8afb6bd, c4cbd2d9,
 e0e7eef5, fc030a11, 181f262d, 343b4249,
 50575e65, 6c737a81, 888f969d, a4abb2b9,
 c0c7ced5, dce3eaf1, f8ff060d, 141b2229,
 30373e45, 4c535a61, 686f767d, 848b9299,
 a0a7aeb5, bcc3cad1, d8dfe6ed, f4fb0209,
 10171e25, 2c333a41, 484f565d, 646b7279.

Annex B

Insert the following line after `id-bc128-seed`:

```
id-bc128-sm4 OID ::= {id-bc128 sm4(4)}
```

Replace { OID id-bc128-seed PARMS KeyLength } , with the following:

```
{ OID id-bc128-seed PARMS KeyLength } |  
{ OID id-bc128-sm4 PARMS KeyLength },
```

Annex D

Change the title to "Numerical examples".

Replace "test vectors" with "numerical examples".

D.1

Change the heading to "General".

Replace the text with the following :

This annex provides numerical examples for TDEA, MISTY1, CAST-128, HIGHT, AES, Camellia, SEED, and SM4 ciphers. In these examples, all data are expressed in hexadecimal.

D.9

Add new clause D.9 as follows:

D.9 SM4 numerical examples

D.9.1 SM4 encryption

Given inputs (plaintext and key), output (ciphertext and subkeys) and intermediate values are described.

Input plaintext: 01 23 45 67 89 ab cd ef fe dc ba 98 76 54 32 10.

Input key: 01 23 45 67 89 ab cd ef fe dc ba 98 76 54 32 10.

The subkeys and the values of the output of each round:

$rk_0 = f12186f9$ $X_4 = 27fad345$

$rk_1 = 41662b61$ $X_5 = a18b4cb2$

$rk_2 = 5a6ab19a$ $X_6 = 11c1e22a$

$rk_3 = 7ba92077$ $X_7 = cc13e2ee$

$rk_4 = 367360f4$ $X_8 = f87c5bd5$

$rk_5 = 776a0c61$ $X_9 = 33220757$

$rk_6 = \text{b6bb89b3}$	$X_{10} = \text{77f4c297}$
$rk_7 = \text{24763151}$	$X_{11} = \text{7a96f2eb}$
$rk_8 = \text{a520307c}$	$X_{12} = \text{27dac07f}$
$rk_9 = \text{b7584dbd}$	$X_{13} = \text{42dd0f19}$
$rk_{10} = \text{c30753ed}$	$X_{14} = \text{b8a5da02}$
$rk_{11} = \text{7ee55b57}$	$X_{15} = \text{907127fa}$
$rk_{12} = \text{6988608c}$	$X_{16} = \text{8b952b83}$
$rk_{13} = \text{30d895b7}$	$X_{17} = \text{d42b7c59}$
$rk_{14} = \text{44ba14af}$	$X_{18} = \text{2ffc5831}$
$rk_{15} = \text{104495a1}$	$X_{19} = \text{f69e6888}$
$rk_{16} = \text{d120b428}$	$X_{20} = \text{af2432c4}$
$rk_{17} = \text{73b55fa3}$	$X_{21} = \text{ed1ec85e}$
$rk_{18} = \text{cc874966}$	$X_{22} = \text{55a3ba22}$
$rk_{19} = \text{92244439}$	$X_{23} = \text{124b18aa}$
$rk_{20} = \text{e89e641f}$	$X_{24} = \text{6ae7725f}$
$rk_{21} = \text{98ca015a}$	$X_{25} = \text{f4cba1f9}$
$rk_{22} = \text{c7159060}$	$X_{26} = \text{1dcdfa10}$
$rk_{23} = \text{99e1fd2e}$	$X_{27} = \text{2ff60603}$
$rk_{24} = \text{b79bd80c}$	$X_{28} = \text{eff24fdc}$
$rk_{25} = \text{1d2115b0}$	$X_{29} = \text{6fe46b75}$
$rk_{26} = \text{0e228aeb}$	$X_{30} = \text{893450ad}$
$rk_{27} = \text{f1780c81}$	$X_{31} = \text{7b938f4c}$
$rk_{28} = \text{428d3654}$	$X_{32} = \text{536e4246}$
$rk_{29} = \text{62293496}$	$X_{33} = \text{86b3e94f}$
$rk_{30} = \text{01cf72e5}$	$X_{34} = \text{d206965e}$
$rk_{31} = \text{9124a012}$	$X_{35} = \text{681edf34}$

The output ciphertext: 68 1e df 34 d2 06 96 5e 86 b3 e9 4f 53 6e 42 46.

D.9.2 SM4 encryption 1 000 000 times

Given inputs (plaintext and key), output (ciphertext) after encryption iteratively 1 000 000 times is described.

Input plaintext: 01 23 45 67 89 ab cd ef fe dc ba 98 76 54 32 10.