
**Information technology — Abstract
Syntax Notation One (ASN.1):
Parameterization of ASN.1 specifications**

*Technologies de l'information — Notation de syntaxe abstraite
numéro un (ASN.1): Paramétrage des spécifications ASN.1*

STANDARDSISO.COM : Click to view the full PDF of ISO/IEC 8824-4:2002

PDF disclaimer

This PDF file may contain embedded typefaces. In accordance with Adobe's licensing policy, this file may be printed or viewed but shall not be edited unless the typefaces which are embedded are licensed to and installed on the computer performing the editing. In downloading this file, parties accept therein the responsibility of not infringing Adobe's licensing policy. The ISO Central Secretariat accepts no liability in this area.

Adobe is a trademark of Adobe Systems Incorporated.

Details of the software products used to create this PDF file can be found in the General Info relative to the file; the PDF-creation parameters were optimized for printing. Every care has been taken to ensure that the file is suitable for use by ISO member bodies. In the unlikely event that a problem relating to it is found, please inform the Central Secretariat at the address given below.

STANDARDSISO.COM : Click to view the full PDF of ISO/IEC 8824-4:2002

© ISO/IEC 2002

All rights reserved. Unless otherwise specified, no part of this publication may be reproduced or utilized in any form or by any means, electronic or mechanical, including photocopying and microfilm, without permission in writing from either ISO at the address below or ISO's member body in the country of the requester.

ISO copyright office
Case postale 56 • CH-1211 Geneva 20
Tel. + 41 22 749 01 11
Fax + 41 22 749 09 47
E-mail copyright@iso.org
Web www.iso.org

Published by ISO in 2003

Published in Switzerland

CONTENTS

	<i>Page</i>
Introduction	v
1 Scope	1
2 Normative references	1
2.1 Identical Recommendations International Standards	1
3 Definitions	1
3.1 Specification of basic notation	1
3.2 Information object specification	1
3.3 Constraint specification	1
3.4 Additional definitions	1
4 Abbreviations	2
5 Convention	2
6 Notation	2
6.1 Assignments	2
6.2 Parameterized definitions	2
6.3 Symbols	3
7 ASN.1 lexical items	3
8 Parameterized assignments	3
9 Referencing parameterized definitions	5
10 Abstract syntax parameters	8
Annex A – Examples	9
A.1 Example of the use of a parameterized type definition	9
A.2 Example of use of parameterized definitions together with an information object class	9
A.3 Example of parameterized type definition that is finite	10
A.4 Example of a parameterized value definition	11
A.5 Example of a parameterized value set definition	11
A.6 Example of a parameterized class definition	11
A.7 Example of a parameterized object set definition	12
A.8 Example of a parameterized object set definition	12
Annex B – Summary of the notation	13

Foreword

ISO (the International Organization for Standardization) and IEC (the International Electrotechnical Commission) form the specialized system for worldwide standardization. National bodies that are members of ISO or IEC participate in the development of International Standards through technical committees established by the respective organization to deal with particular fields of technical activity. ISO and IEC technical committees collaborate in fields of mutual interest. Other international organizations, governmental and non-governmental, in liaison with ISO and IEC, also take part in the work. In the field of information technology, ISO and IEC have established a joint technical committee, ISO/IEC JTC 1.

International Standards are drafted in accordance with the rules given in the ISO/IEC Directives, Part 2.

The main task of the joint technical committee is to prepare International Standards. Draft International Standards adopted by the joint technical committee are circulated to national bodies for voting. Publication as an International Standard requires approval by at least 75 % of the national bodies casting a vote.

Attention is drawn to the possibility that some of the elements of this document may be the subject of patent rights. ISO and IEC shall not be held responsible for identifying any or all such patent rights.

ISO/IEC 8824-4 was prepared by Joint Technical Committee ISO/IEC JTC 1, *Information technology*, Subcommittee SC 6, *Telecommunications and information exchange between systems*, in collaboration with ITU-T. The identical text is published as ITU-T Rec. X.683.

This third edition cancels and replaces the second edition (ISO/IEC 8824-4:1998), which has been technically revised. It also incorporates the Amendment ISO/IEC 8824-4:1998/Amd.1:2000.

ISO/IEC 8824 consists of the following parts, under the general title *Information technology — Abstract Syntax Notation One (ASN.1)*:

- *Part 1: Specification of basic notation*
- *Part 2: Information object specification*
- *Part 3: Constraint specification*
- *Part 4: Parameterization of ASN.1 specifications*

Introduction

Application designers need to write specifications in which certain aspects are left undefined. Those aspects will later be defined by one or more other groups (each in its own way), to produce a fully defined specification for use in the definition of an abstract syntax (one for each group).

In some cases, aspects of the specification (for example, bounds) may be left undefined even at the time of abstract syntax definition, being completed by the specification of International Standardized Profiles or functional profiles from some other body.

NOTE 1 – It is a requirement imposed by this Recommendation | International Standard that any aspect that is not solely concerned with the application of constraints has to be completed prior to the definition of an abstract syntax.

In the extreme case, some aspects of the specification may be left for the implementor to complete, and would then be specified as part of the Protocol Implementation Conformance Statement.

While the provisions of ITU-T Rec. X.681 | ISO/IEC 8824-2 and ITU-T Rec. X.682 | ISO/IEC 8824-3 provide a framework for the later completion of parts of a specification, they do not of themselves solve the above requirements.

Additionally, a single designer sometimes requires to define many types, or many information object classes, or many information object sets, or many information objects, or many values, which have the same outer level structure, but differ in the types, or information object classes, or information object sets, or information objects, or values, that are used at an inner level. Instead of writing out the outer level structure for every such occurrence, it is useful to be able to write it out once, with parts left to be defined later, then to refer to it and provide the additional information.

All these requirements are met by the provision for parameterized reference names and parameterized assignments by this Recommendation | International Standard.

The syntactic form of a parameterized reference name is the same as that of the corresponding normal reference name, but the following additional considerations apply:

- When it is assigned in a parameterized assignment statement, it is followed by a list of dummy reference names in braces, each possibly accompanied by a governor; these reference names have a scope which is the right-hand side of the assignment statement, and the parameter list itself.
NOTE 2 – This is what causes it to be recognized as a parameterized reference name.
- When it is exported or imported, it is followed by a pair of empty braces to distinguish it as a parameterized reference name.
- When it is used in any construct, it is followed by a list of syntactic constructions, one for each dummy reference name, that provide an assignment to the dummy reference name for the purposes of that use only.

Dummy reference names have the same syntactic form as the corresponding normal reference name, and can be used anywhere on the right-hand side of the assignment statement that the corresponding normal reference name could be used. All such usages are required to be consistent.

INTERNATIONAL STANDARD

ITU-T RECOMMENDATION

**Information technology –
Abstract Syntax Notation One (ASN.1):
Parameterization of ASN.1 specifications**

1 Scope

This Recommendation | International Standard is part of Abstract Syntax Notation One (ASN.1) and defines notation for parameterization of ASN.1 specifications.

2 Normative references

The following Recommendations and International Standards contain provisions which, through reference in this text, constitute provisions of this Recommendation | International Standard. At the time of publication, the editions indicated were valid. All Recommendations and Standards are subject to revision, and parties to agreements based on this Recommendation | International Standard are encouraged to investigate the possibility of applying the most recent edition of the Recommendations and Standards listed below. Members of IEC and ISO maintain registers of currently valid International Standards. The Telecommunication Standardization Bureau of the ITU maintains a list of currently valid ITU-T Recommendations.

2.1 Identical Recommendations | International Standards

- ITU-T Recommendation X.680 (2002) | ISO/IEC 8824-1:2002, *Information technology – Abstract Syntax Notation One (ASN.1): Specification of basic notation*.
- ITU-T Recommendation X.681 (2002) | ISO/IEC 8824-2:2002, *Information technology – Abstract Syntax Notation One (ASN.1): Information object specification*.
- ITU-T Recommendation X.682 (2002) | ISO/IEC 8824-3:2002, *Information technology – Abstract Syntax Notation One (ASN.1): Constraint specification*.

3 Definitions

For the purposes of this Recommendation | International Standard, the following definitions apply.

3.1 Specification of basic notation

This Recommendation | International Standard uses the terms defined in ITU-T Rec. X.680 | ISO/IEC 8824-1.

3.2 Information object specification

This Recommendation | International Standard uses the terms defined in ITU-T Rec. X.681 | ISO/IEC 8824-2.

3.3 Constraint specification

This Recommendation | International Standard uses the terms defined in ITU-T Rec. X.682 | ISO/IEC 8824-3.

3.4 Additional definitions

3.4.1 normal reference name: A reference name defined, without parameters, by means of an "Assignment" other than a "ParameterizedAssignment". Such a name references a complete definition and is not supplied with actual parameters when used.

3.4.2 parameterized reference name: A reference name defined using a parameterized assignment, which references an incomplete definition and which, therefore, must be supplied with actual parameters when used.

3.4.3 parameterized type: A type defined using a parameterized type assignment and thus whose components are incomplete definitions which must be supplied with actual parameters when the type is used.

3.4.4 parameterized value: A value defined using a parameterized value assignment and thus whose value is incompletely specified and must be supplied with actual parameters when used.

3.4.5 parameterized value set: A value set defined using a parameterized value set assignment and thus whose values are incompletely specified and must be supplied with actual parameters when used.

3.4.6 parameterized object class: An information object class defined using a parameterized object class assignment and thus whose field specifications are incompletely specified and must be supplied with actual parameters when used.

3.4.7 parameterized object: An information object defined using a parameterized object assignment and thus whose components are incompletely specified and must be supplied with actual parameters when used.

3.4.8 parameterized object set: An information object set defined using a parameterized object set assignment and thus whose objects are incompletely specified and must be supplied with actual parameters when used.

3.4.9 variable constraint: A constraint employed in specifying a parameterized abstract syntax, and which depends on some parameter of the abstract syntax.

4 Abbreviations

For the purposes of this Recommendation | International Standard, the following abbreviation applies:

ASN.1 Abstract Syntax Notation One

5 Convention

This Recommendation | International Standard employs the notational convention defined in ITU-T Rec. X.680 | ISO/IEC 8824-1, clause 5.

6 Notation

This clause summarizes the notation defined in this Recommendation | International Standard.

6.1 Assignments

The following notation which can be used as an alternative for "Assignment" (see ITU-T Rec. X.680 | ISO/IEC 8824-1, clause 12) is defined in this Recommendation | International Standard:

- ParameterizedAssignment (see 8.1).

6.2 Parameterized definitions

6.2.1 The following notation which can be used as an alternative for "DefinedType" (see ITU-T Rec. X.680 | ISO/IEC 8824-1, 13.1) is defined in this Recommendation | International Standard:

- ParameterizedType (see 9.2).

6.2.2 The following notation which can be used as an alternative for "DefinedValue" (see ITU-T Rec. X.680 | ISO/IEC 8824-1, 13.1) is defined in this Recommendation | International Standard:

- ParameterizedValue (see 9.2).

6.2.3 The following notation which can be used as an alternative for "DefinedType" (see ITU-T Rec. X.680 | ISO/IEC 8824-1, 13.1) is defined in this Recommendation | International Standard:

- ParameterizedValueSetType (see 9.2).

6.2.4 The following notation which can be used as an alternative for "ObjectClass" (see ITU-T Rec. X.681 | ISO/IEC 8824-2, 9.2) is defined in this Recommendation | International Standard:

- ParameterizedObjectClass (see 9.2).

6.2.5 The following notation which can be used as an alternative for "Object" (see ITU-T Rec. X.681 | ISO/IEC 8824-2, 11.3) is defined in this Recommendation | International Standard:

- ParameterizedObject (see 9.2).

6.2.6 The following notation which can be used as an alternative for "ObjectSet" (see ITU-T Rec. X.681 | ISO/IEC 8824-2, 12.3) is defined in this Recommendation | International Standard:

- ParameterizedObjectSet (see 9.2).

6.3 Symbols

The following notation which can be used as an alternative for "Symbol" (see ITU-T Rec. X.680 | ISO/IEC 8824-1, 12.1) is defined in this Recommendation | International Standard:

- ParameterizedReference (see 9.1).

7 ASN.1 lexical items

This Recommendation | International Standard makes use of the lexical items specified in ITU-T Rec. X.680 | ISO/IEC 8824-1, clause 11.

8 Parameterized assignments

8.1 There are parameterized assignment statements corresponding to each of the assignment statements specified in ITU-T Rec. X.680 | ISO/IEC 8824-1 and ITU-T Rec. X.681 | ISO/IEC 8824-2. The "ParameterizedAssignment" construct is:

```
ParameterizedAssignment ::=
    ParameterizedTypeAssignment
    | ParameterizedValueAssignment
    | ParameterizedValueSetTypeAssignment
    | ParameterizedObjectClassAssignment
    | ParameterizedObjectAssignment
    | ParameterizedObjectSetAssignment
```

8.2 Each "Parameterized<X>Assignment" has the same syntax as "<X>Assignment" except that following the initial lexical item there is a "ParameterList". The initial item thereby becomes a parameterized reference name (see 3.4.2):

NOTE – ITU-T Rec. X.680 | ISO/IEC 8824-1 imposes the requirement that all reference names assigned within a module, whether parameterized or not, must be distinct.

```
ParameterizedTypeAssignment ::=
    typereference
    ParameterList
    ":" ::=
    Type
```

```
ParameterizedValueAssignment ::=
    valureference
    ParameterList
    Type
    ":" ::=
    Value
```

```
ParameterizedValueSetTypeAssignment ::=
    typereference
    ParameterList
    Type
    ":" ::=
    ValueSet
```

```
ParameterizedObjectClassAssignment ::=
    objectclassreference
    ParameterList
```

" : :="
ObjectClass

ParameterizedObjectAssignment ::=
objectreference
ParameterList
DefinedObjectClass
" : :="
Object

ParameterizedObjectSetAssignment ::=
objectsetreference
ParameterList
DefinedObjectClass
" : :="
ObjectSet

8.3 A "ParameterList" is a list of "Parameter"s between braces:

ParameterList ::= "{" Parameter "," + "}"

Each "Parameter" consists of a "DummyReference" and possibly a "ParamGovernor":

Parameter ::= ParamGovernor " : " DummyReference | DummyReference

ParamGovernor ::= Governor | DummyGovernor

Governor ::= Type | DefinedObjectClass

DummyGovernor ::= DummyReference

DummyReference ::= Reference

A "DummyReference" in "Parameter" may stand for:

- a "Type" or "DefinedObjectClass", in which case there shall be no "ParamGovernor";
- a "Value" or "ValueSet", in which case the "ParamGovernor" shall be present, and in case "ParamGovernor" is a "Governor" it shall be a "Type", and in case "ParamGovernor" is a "DummyGovernor" the actual parameter for the "ParamGovernor" shall be a "Type";
- an "Object" or "ObjectSet", in which case the "ParamGovernor" shall be present, and in case "ParamGovernor" is a "Governor" it shall be a "DefinedObjectClass", and in case "ParamGovernor" is a "DummyGovernor" the actual parameter for the "ParamGovernor" shall be a "DefinedObjectClass".

A "DummyGovernor" shall be a "DummyReference" that has no "Governor".

8.4 The scope of a "DummyReference" appearing in a "ParameterList" is the "ParameterList" itself, together with that part of the "ParameterizedAssignment" which follows the " : :=". The "DummyReference" hides any other "Reference" with the same name in that scope.

NOTE – This subclause does not apply to "identifier"s defined in "NamedNumberList"s, "Enumeration"s and "NamedBitList"s, since they are not "Reference"s. The "DummyReference" does not hide these "identifier"s (see ITU-T Rec. X.680 | ISO/IEC 8824-1, 18.11 and 19.10).

8.5 The usage of a "DummyReference" within its scope shall be consistent with its syntactic form, and, where applicable, governor, and all usages of the same "DummyReference" shall be consistent with one another.

NOTE – Where the syntactic form of a dummy reference name is ambiguous (for example, between whether it is an "objectclassreference" or "typereference"), the ambiguity can normally be resolved on the first use of the dummy reference name on the right-hand side of the assignment statement. Thereafter, the nature of the dummy reference name is known. The nature of the dummy reference is, however, not determined solely by the right-hand side of the assignment statement when it is in turn used only as an actual parameter in a parameterized reference; in this case, the nature of the dummy reference must be determined by examining the definition of this parameterized reference. Users of the notation are warned that such a practice can make ASN.1 specifications less clear, and it is suggested that adequate comments are provided to explain this for human readers.

Example

Consider the following parameterized object class assignment:

```
PARAMETERIZED-OBJECT-CLASS { TypeParam, INTEGER:valueParam, INTEGER:ValueSetParam }  
::=  
CLASS {  
  &valueField1          TypeParam,
```

```

&valueField2      INTEGER DEFAULT valueParam,
&valueField3      INTEGER (ValueSetParam),
&ValueSetField    INTEGER DEFAULT { ValueSetParam }
}

```

For the purpose of determining proper usage of the "DummyReference"s in the scope of the "Parameterized Assignment", and for that purpose only, the "DummyReference"s can be regarded to be defined as follows:

TypeParam ::= UnspecifiedType

valueParam INTEGER ::= unspecifiedIntegerValue

ValueSetParam INTEGER ::= { UnspecifiedIntegerValueSet }

where:

- TypeParam** is a "DummyReference" which stands for a "Type". Therefore **TypeParam** can be used wherever a "typereference" can be used, e.g. as a "Type" for the fixed-type value field **valueField1**.
- valueParam** is a "DummyReference" which stands for a value of an integer type. Therefore **valueParam** can be used wherever a "valuereference" of an integer value can be used, e.g. as a default value for the fixed-type value field **valueField2**.
- ValueSetParam** is a "DummyReference" which stands for a value set of an integer type. Therefore **ValueSetParam** can be used wherever a "typereference" of an integer value can be used, e.g. as a "Type" in the "ContainedSubtype" notation for **valueField3** and **ValueSetField**.

8.6 Each "DummyReference" shall be employed at least once within its scope.

NOTE – If the "DummyReference" did not so appear, then the corresponding "ActualParameter" would have no effect on the definition, and would simply be "discarded", while to the user it might seem that some specification was taking place.

"ParameterizedValueAssignment"s, "ParameterizedValueSetTypeAssignment"s, "ParameterizedObjectAssignment"s and "ParameterizedObjectSetAssignment"s that contain either a direct or indirect reference to themselves are illegal.

8.7 In the definition of a "ParameterizedType", "ParameterizedValueSet", or "ParameterizedObjectClass", a "DummyReference" shall not be passed as a tagged type (as an actual parameter) to a recursive reference to that "ParameterizedType", "ParameterizedValueSet", or "ParameterizedObjectClass" (see A.3).

8.8 In the definition of a "ParameterizedType", "ParameterizedValueSet", or "ParameterizedObjectClass", a circular reference to the item being defined shall not be made, unless such reference is directly or indirectly marked **OPTIONAL** or, in the case of "ParameterizedType" and "ParameterizedValueSet", made through a reference to a choice type, at least one of whose alternatives is non-circular in definition.

8.9 The governor of a "DummyReference" shall not include a reference to another "DummyReference" if that other "DummyReference" also has a governor.

8.10 In a parameterized assignment the right side of the ": :=" shall not consist solely of a "DummyReference".

8.11 The governor of a "DummyReference" shall not require knowledge of either the "DummyReference" or of the parameterized reference name being defined.

8.12 When a value or value set is supplied to a parameterized type as an actual parameter, the type of the actual parameter is required to be compatible with the governor of the corresponding dummy parameter. (See ITU-T Rec. X.680 | ISO/IEC 8824-1, B.6.2 and B.6.3 for details.)

8.13 In defining a parameterized type with a value or a value set dummy parameter, the type used to govern that dummy parameter shall be a type, all of whose values are valid for use in all places to the right of the assignment where the dummy parameter is used. (See ITU-T Rec. X.680 | ISO/IEC 8824-1, B.6.5 for details.)

9 Referencing parameterized definitions

9.1 Within a "SymbolList" (in "Exports" or "Imports") a parameterized definition shall be referenced by a "ParameterizedReference":

ParameterizedReference ::= Reference | Reference "{" "}"

where "Reference" is the first lexical item in the "ParameterizedAssignment", as specified in 8.2 above.

NOTE – The first alternative of "ParameterizedReference" is provided solely as an aid to human understanding. Both alternatives have the same meaning.

9.2 Other than in "Exports" or "Imports", a parameterized definition shall be referenced by a "Parameterized<X>" construct, which can be used as an alternative for the corresponding "<X>":

```

ParameterizedType ::=
    SimpleDefinedType
    ActualParameterList

SimpleDefinedType ::=
    ExternalTypeReference |
    typereference

ParameterizedValue ::=
    SimpleDefinedValue
    ActualParameterList

SimpleDefinedValue ::=
    ExternalValueReference |
    valuereference

ParameterizedValueSetType ::=
    SimpleDefinedType
    ActualParameterList

ParameterizedObjectClass ::=
    DefinedObjectClass
    ActualParameterList

ParameterizedObjectSet ::=
    DefinedObjectSet
    ActualParameterList

ParameterizedObject ::=
    DefinedObject
    ActualParameterList

```

9.3 The reference name in the "Defined<X>" shall be a reference name to which an assignment is made in a "ParameterizedAssignment".

9.4 The restrictions on the "Defined<X>" alternative to be used, which are specified in ITU-T Rec. X.680 | ISO/IEC 8824-1 and ITU-T Rec. X.681 | ISO/IEC 8824-2 as normal reference names, apply equally to the corresponding parameterized reference names.

NOTE – In essence, the restrictions are as follows: each "Defined<X>" has two alternatives, "<x>reference" and "External<x>Reference". The former is used within the module of definition or if the definition has been imported and there is no name conflict; the latter is used where there is no imports listed (deprecated), or if there is a conflict between the imported name and a local definition (also deprecated) or between imports.

9.5 The "ActualParameterList" is:

```

ActualParameterList ::=
    "{" ActualParameter "," + "}"

ActualParameter ::=
    Type
    Value
    ValueSet
    DefinedObjectClass
    Object
    ObjectSet

```

9.6 There shall be exactly one "ActualParameter" for each "Parameter" in the corresponding "ParameterizedAssignment" and they shall appear in the same order. The particular choice of "ActualParameter", and the governor (if any) shall be determined by examination of the syntactic form of the "Parameter" and the environment in which it occurs in the "ParameterizedAssignment". The form of the "ActualParameter" shall be the form required to replace the "DummyReference" everywhere in its scope (see 8.4).

Example

The parameterized object class definition of the previous example (see 8.5) can be referenced, for instance, as follows:

MY-OBJECT-CLASS ::= PARAMETERIZED-OBJECT-CLASS { BIT STRING, 123, {4 | 5 | 6} }

9.7 The actual parameter takes the place of the dummy reference name in determining the actual type, value, value set, object class, object, or object set that is being referenced by this instance of use of the parameterized reference name.

9.8 The meaning of any references which appear in the "ActualParameter", and the tag default applicable to any tags which so appear, are determined according to the tagging environment of the "ActualParameter" rather than that of the corresponding "DummyReference".

NOTE – Thus, parameterization, like referencing, selection types, and **COMPONENTS OF**, among others, is not exactly textual substitution.

Example

Consider the following modules:

```

M1 DEFINITIONS AUTOMATIC TAGS ::= BEGIN
    EXPORTS T1;

    T1 ::= SET {
        f1    INTEGER,
        f2    BOOLEAN
    }

END

M2 DEFINITIONS EXPLICIT TAGS ::= BEGIN
    IMPORTS T1 FROM M1;

    T3 ::= T2{T1}

    T2{X} ::= SEQUENCE {
        a    INTEGER,
        b    X
    }

END

```

Application of 9.8 implies that the tag for the component **f1** of **T3** (i.e. @T3.b.f1) will be implicitly tagged because the tagging environment of the dummy parameter **X**, namely explicit tagging, does not affect the tagging of the components of the actual parameter **T1**.

Consider the module **M3**:

```

M3 DEFINITIONS AUTOMATIC TAGS ::= BEGIN
    IMPORTS T1 FROM M1;

    T5 ::= T4{T1}

    T4{Y} ::= SEQUENCE {
        a    INTEGER,
        b    Y
    }

END

```

Application of ITU-T Rec. X.680 | ISO/IEC 8824-1, 30.6, implies that the tag for the component **b** of **T5** (i.e. @T5.b) will be explicitly tagged because the dummy parameter **Y** is always explicitly tagged, hence **T5** is equivalent to:

```

T5 ::= SEQUENCE {
    a    [0] IMPLICIT INTEGER,
    b    [1] EXPLICIT SET {
        f1    [0] INTEGER,
        f2    [1] BOOLEAN
    }
}

```

while **T3** is equivalent to:

```

T3 ::= SEQUENCE {
    a    INTEGER,
    b    SET {
        f1    [0] IMPLICIT INTEGER,
        f2    [1] IMPLICIT BOOLEAN
    }
}

```

10 Abstract syntax parameters

10.1 ITU-T Rec. X.681 | ISO/IEC 8824-2, Annex B, provides the **ABSTRACT-SYNTAX** information object class and recommends its use to define abstract syntaxes, using as an example an abstract syntax defined as the set of values of a single ASN.1 type which was not parameterized at the outer level.

10.2 Where the ASN.1 type used to define the abstract syntax *is* parameterized, some parameters may be supplied as actual parameters when the abstract syntax is defined, while others may be left as parameters of the abstract syntax itself.

Example

If a parameterized type has been defined called **YYY-PDU** with two dummy references (the first an object set of some defined object class, and the second an integer value for a bound, say), then:

```
yyy-Abstract-Syntax { INTEGER:bound } ABSTRACT-SYNTAX ::=  
{ YYY-PDU { {ValidObjects} , bound } IDENTIFIED BY {yyy 5} }
```

defines a parameterized abstract syntax in which the object set has been resolved, but **bound** remains as a parameter of the abstract syntax.

An abstract syntax parameter shall be used:

- a) directly or indirectly in the context of a constraint;
- b) directly or indirectly as actual parameters that eventually are used in the context of a constraint.

NOTE – See the example in A.2, and the example in ITU-T Rec. X.680 | ISO/IEC 8824-1, F.5.

10.3 A constraint whose value set depends on one or more parameters of the abstract syntax is a variable constraint. Such constraints are determined after the definition of the abstract syntax (perhaps by International Standardized Profiles or in Protocol Implementation Conformance Statements).

NOTE – If somewhere in the chain of definitions involved in the specification of the constraint values a parameter of the abstract syntax appears, the constraint is a variable constraint. It is a variable constraint even if the value set of the resulting constraint is independent of the actual value of the parameter of the abstract syntax.

Example

The value of **((1..3) EXCEPT a) UNION (1 .. 3)** is always 1..3 no matter what the value of **a** is, nonetheless it is still a variable constraint if **a** is a parameter of the abstract syntax.

10.4 Formally, a variable constraint does not constrain the set of values in the abstract syntax.

NOTE – It is strongly recommended that constraints that are expected to remain as variable constraints in an abstract syntax have an exception specification using the notation provided by ITU-T Rec. X.680 | ISO/IEC 8824-1, 49.4.

Annex A

Examples

(This annex does not form an integral part of this Recommendation | International Standard)

A.1 Example of the use of a parameterized type definition

Suppose that a protocol designer frequently needs to carry an authenticator with one or more of the fields of the protocol. This will be carried as a **BIT STRING**, alongside the field. Without parameterization, **Authenticator** would need to be defined as a **BIT STRING**, then **authenticator** would need to be added wherever it was to appear, with text to identify what it applied to. Alternatively, the designer could adopt the discipline of turning any field that has an authenticator into a **SEQUENCE** of that field and **authenticator**. The parameterization mechanism provides a convenient short-hand for doing this task.

First we define the parameterized type **SIGNED{}**:

```
SIGNED { ToBeSigned } ::= SEQUENCE
{
    authenticated-data      ToBeSigned,
    authenticator           BIT STRING
}
```

then, in the body of the protocol, the notation (for example):

```
SIGNED { OrderInformation }
```

is a type notation standing for:

```
SEQUENCE
{
    authenticated-data      OrderInformation,
    authenticator           BIT STRING
}
```

Suppose further that for some fields, the sender is to have the option of adding the authenticator or not. This could be achieved by making the **BIT STRING** optional, but a more elegant solution (less bits on the line) would be to define another parameterized type:

```
OPTIONALLY-SIGNED {ToBeSigned} ::= CHOICE
{
    unsigned-data [0]      ToBeSigned,
    signed-data   [1]      SIGNED { ToBeSigned }
}
```

NOTE – The tagging in the **CHOICE** is not necessary if the writer ensures that none of the uses of the parameterized type produce an actual argument which is a **BIT STRING** (the type of **SIGNED**), but is useful in preventing errors in other parts of the specification.

A.2 Example of use of parameterized definitions together with an information object class

Use information object classes to collect all the parameters for an abstract syntax. In that way the number of parameters for an abstract syntax can be reduced to one which is an instance of the collection class. The "InformationFromObject" production can be used to extract information from the parameter object.

Example

```
-- An instance of this class contains all the parameters for the abstract
-- syntax, Message-PDU.
```

```
MESSAGE-PARAMETERS ::= CLASS {
    &maximum-priority-level      INTEGER,
    &maximum-message-buffer-size INTEGER,
    &maximum-reference-buffer-size  INTEGER
}
WITH SYNTAX {
    THE MAXIMUM PRIORITY LEVEL IS      &maximum-priority-level
    THE MAXIMUM MESSAGE BUFFER SIZE IS  &maximum-message-buffer-
size
    THE MAXIMUM REFERENCE BUFFER SIZE IS  &maximum-reference-buffer-
size
}
```

```

    }
    -- The "ValueFromObject" production is used to extract values
    -- from the abstract syntax parameter, "param". The values can be
    -- used only in constraints. In addition the parameter is passed
    -- through to another parameterized type.

    Message-PDU { MESSAGE-PARAMETERS : param } ::= SEQUENCE {
        priority-level    INTEGER (0..param.&maximum-priority-level),
        message           BMPString (SIZE (0..param.&maximum-message-buffer-size)),
        reference         Reference { param }
    }

    Reference { MESSAGE-PARAMETERS : param } ::=
        SEQUENCE OF IA5String (SIZE (0..param.&maximum-reference-buffer-size))
    -- Definition of a parameterized abstract syntax information object.
    -- The abstract syntax parameter is used only in constraints.

    message-Abstract-Syntax { MESSAGE-PARAMETERS : param }
        ABSTRACT-SYNTAX ::=
        {
            Message-PDU { param }
            IDENTIFIED BY { joint-iso-ccitt asn1(1) examples(123) 0 }
        }

```

The class **MESSAGE-PARAMETERS** and the parameterized abstract syntax object, **message-Abstract-Syntax**, are used as follows:

```

-- This instance of MESSAGE-PARAMETERS defines parameter values
-- for the abstract syntax.

my-message-parameters MESSAGE-PARAMETERS ::= {
    THE MAXIMUM PRIORITY LEVEL IS 10
    THE MAXIMUM MESSAGE BUFFER SIZE IS 2000
    THE MAXIMUM REFERENCE BUFFER SIZE IS 100
}

-- The abstract syntax can now be defined with all variable constraints specified.

my-message-Abstract-Syntax ABSTRACT-SYNTAX ::=
    message-Abstract-Syntax { my-message-parameters }

```

A.3 Example of parameterized type definition that is finite

When specifying a parameterized type which represents a generic list, specify the type so that the resulting ASN.1 notation is finite. For example, we may specify:

```

List1 { ElementTypeParam } ::= SEQUENCE {
    elem ElementTypeParam,
    next List1 { ElementTypeParam } OPTIONAL
}

```

which is finite, for when it is used:

```

IntegerList1 ::= List1 { INTEGER }

```

the resulting ASN.1 notation is as you would normally define it:

```

IntegerList1 ::= SEQUENCE {
    elem INTEGER,
    next IntegerList1 OPTIONAL
}

```

Contrast this to the following:

```

List2 { ElementTypeParam } ::= SEQUENCE {
    elem ElementTypeParam,
    next List2 { [0] ElementTypeParam } OPTIONAL
}

IntegerList2 ::= List2 { INTEGER }

```

where the resulting ASN.1 notation is infinite:

```

IntegerList2 ::= SEQUENCE {
    elem INTEGER,
    next SEQUENCE {
        elem [0] INTEGER,
        next SEQUENCE {
            elem [0][0] INTEGER,
            next SEQUENCE {
                elem [0][0][0] INTEGER,
                next SEQUENCE {
                    -- and so on
                } OPTIONAL
            } OPTIONAL
        } OPTIONAL
    } OPTIONAL
}

```

A.4 Example of a parameterized value definition

If a parameterized string value is defined as follows:

```

genericBirthdayGreeting { IA5String : name } IA5String ::= { "Happy birthday, ",
name, "!!" }

```

then the following two string values are the same:

```

greeting1 IA5String ::= genericBirthdayGreeting { "John" }
greeting2 IA5String ::= "Happy birthday, John!!"

```

A.5 Example of a parameterized value set definition

If two parameterized value sets are defined as follows:

```

QuestList1 {IA5String : extraQuest} IA5String ::= { "Jack" | "John" | extraQuest }
QuestList2 {IA5String : ExtraQuests} IA5String ::= { "Jack" | "John" |
ExtraQuests }

```

then the following value sets denote the same value set:

```

SetOfQuests1 IA5String ::= { QuestList1 { "Jill" } }
SetOfQuests2 IA5String ::= { QuestList2 { {"Jill"} } }
SetOfQuests3 IA5String ::= { "Jack" | "John" | "Jill" }

```

and the following value sets denote the same value set:

```

SetOfQuests4 IA5String ::= { QuestList2 { {"Jill" | "Mary"} } }
SetOfQuests5 IA5String ::= { "Jack" | "John" | "Jill" | "Mary" }

```

Notice that a value set is *always* specified within braces, even when it is a parameterized value set reference. By omitting the braces from a reference to an "identifier" that was created in a value set assignment or from a reference to a "ParameterizedValueSetType" the notation is that of a "Type", not a value set.

A.6 Example of a parameterized class definition

The following parameterized class can be used to define error classes which contain error codes of different types. Note that the **ErrorCodeType** parameter is used only as a "DummyGovernor" for the **ValidErrorCodes** parameter:

```

GENERIC-ERROR { ErrorCodeType, ErrorCodeType : ValidErrorCodes } ::= CLASS {
    &errorCode ValidErrorCodes
}
WITH SYNTAX {
    CODE &errorCode
}

```

The parameterized class definition can be used as follows to define different classes which share some characteristics like the same defined syntax: